

**CONSTRUCTION AND OPTIMALITY OF  
EXTENDED MOST BALANCED GROUP DIVISIBLE  
TREATMENT INCOMPLETE BLOCK DESIGNS**

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**DEDICATED TO:**

**MY MOTHER**

**AND**

**THE MEMORY OF MY FATHER**

ABSTRACT

CONSTRUCTION AND OPTIMALITY OF EXTENDED MOST BALANCED  
GROUP DIVISIBLE TREATMENT INCOMPLETE BLOCK DESIGNS

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The Balanced Treatment Incomplete Block (BTIB) Designs were proposed by Bechhofer and Tamhane (1981) for the problem of comparing  $v \geq 2$  test-treatments with a control in "b" blocks of size "k" each. The requirement of complete symmetry of information matrix (called M matrix), however, imposes severe restriction on the class of such designs and, as a consequence, there are a number of design parameters for which BTIB designs do not exist. A new class of designs, henceforth called Extended Most Balanced Group Divisible Treatment Incomplete Block ( E MB GD TIB ) Designs, is therefore introduced here which, upto some

extent, relaxes the condition of complete symmetry of the information matrix and hence, leads to the study of a larger class, which includes BTIB designs as a sub-class. The symmetry of the information matrix is, however, not affected much and still the elements of the leading diagonal remain equal. Thus conventional idea of A-optimality (minimizing the trace of  $M^{-1}$  matrix) remain intact and A-optimal designs can, therefore, be obtained easily.

The minimal complete class of generator designs ( MCCGD ), (a finite set of elementary designs which are used in the construction of all other admissible and non-equivalent designs in the class) are constructed for  $k = 2,3,4$  and  $v = 4,6,8,10$ . The A-optimal designs for all these  $k$  and  $v$  values are also obtained.

In addition, very powerful and practical computer algorithms, programs/subroutines are designed, which can run very efficiently to provide all the designs as well as the A-optimal designs for any given  $(k,v)$ . Also, programs are developed which help much in finding minimal complete class of generator designs and can provide a complete structure for any design in the class.

**KEY WORDS:** BTIB Designs, Treatment-control Comparison, Reinforced Designs, Group-divisible Designs, A-optimal Designs, Optimal Designs.

## ÖZET

# GELİŞTİRİLMİŞ EN DENGELİ GRUPLANDIRILABİLİR MUAMELELİ TAM OLMAYAN BLOK TASARIMLARININ HAZIRLANMASI VE OPTİMALLIĞI

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Dengeli muameleli tam olmayan Blok (BTIB) tasarımları, Bechoffer ve Tamhane (1981) tarafından  $v > 2$  test-muamelesinin "b" sayıdaki blokta bir kontrol ile mukayesesi problemi için önerilmişti. Bununla beraber bilgi matrisinin (M matrisi) tam simetriye sahip olması koşulu, bu sınıftaki tasarımlar üzerine katı kısıtlamalar getirmektedir ve sonuç olarak da bazı tasarım parametreleri için BTIB tasarımları bulmak olanaksızdır. Bundan böyle "Geliştirilmiş En Dengeli Gruplandırılabilir Muameleli Tam Olmayan Blok Tasarımları" (E MB GD TIB) diye adlandırılan yeni sınıf tasarımlar burada takdim olunacaktır ki, bir dereceye kadar bu tasarımlar bilgi matrisinin tam simetri koşulunu aramayacak ve dolayısıyla, BTIB tasarımlarını da bir alt sınıf olarak içeren daha geniş bir sınıftaki tasarımların incelenmesine olanak verecektir. Maamafih,

bilgi matrisinin simetrisi çok etkilenmemekte ve baştaki köşegen elemanları eşit olarak kalmaktadır. Böylece, alışlagelmiş olan A-optimal olma ( $M^{-1}$  matrisi'nin izinin minimum olması) fikri korunmuştur ve A-optimal tasarımlarını kolayca elde etmek mümkün olacaktır.

En küçük tam sınıf oluşturan jeneratör tasarımlar (MCCGD), (bunlar sonlu sayıda olup diğer bütün kabul edilebilir ve aynı sınıfta eşdeğeri bulunmayan tasarımların oluşturulmasında kullanılan elemanter tasarımlardır)  $k=2,3,4$  ve  $v=4,6,8,10$  için oluşturulmuştur. Bütün bu  $k$  ve  $v$  değerleri için A-optimal tasarımlar da elde edilmiştir.

Ayrıca, çok güçlü ve kullanışlı bilgisayar algoritmaları, programları/alt-programları hazırlanmıştır, ki bunlar verilen herhangi  $(k,v)$  için A-optimal tasarımların yanında diğer bütün tasarımlarında elde edilmesinde çok randımanlı olarak kullanılabilir. Bunlara ilaveten, en küçük tam sınıf jeneratör tasarımların bulunmasında çok faydalı olacak olan programlar da mevcuttur ve bunlar bu sınıftaki herhangi bir tasarım için bütün yapıyı verebileceklerdir.

Anahtar Kelimeler: BTIB Tasarımlar, Muamele-kontrol Mukayeseleri, Takviyeli Tasarımlar, Gruplara Ayrılabilen Tasarımlar, A-optimal Tasarımlar, Optimal Tasarımlar.

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LIST OF SYMBOLS

$b$	Number of blocks
$k_j$	Size of $j$ th block
$n$	Total experimental size
$y_{iju}$	Response of $u$ th plot in $j$ th block treated by $i$ th treatment
$\bar{y}$	Mean response
$\alpha_i$	$i$ th treatment effect
$\beta_j$	$j$ th block effect
$e_{iju}$	Error term
$v$	Number of treatments
$\phi$	Incidence matrix for treatments
$\psi$	Incidence matrix for blocks
$y$	Vector of responses
$J_n$	$n \times 1$ vector of unity
$\alpha$	Vector of treatment effects
$\beta$	Vector of block effects
$e$	Vector of errors
$r_i$	Total number of times $i$ th treatment is repeated
$r_{ij}$	Number of times $i$ th treatment is repeated in $j$ th block
$M$	Information matrix
$\hat{\alpha}$	Vector of estimates of treatment effects
$Q$	Vector of adjusted treatment total
$m_{it}$	$(i, t)$ th element of vector $M$

LIST OF SYMBOLS (CONTINUED)

$\lambda_{il}$	Number of times $i$ th and $l$ th treatment appear together in a block over the whole design
$\lambda_0$	Number of times a test-treatment appears with control in a block in a design
$\lambda_1$	Number of times a test-treatment appears with another test-treatment in a block in a design
D	Represents a design
$\lambda_2$	Number of times a test-treatment appears with another test-treatment (of second associates) in a block
$m_1$	Control appears "1" times in " $m_1$ " blocks
r	Number of times each test-treatment is repeated
$r_0$	Number of times the control is repeated
N	Number of TT-pairs
$M_j$	Number of TT-pairs in $j$ th BB combination
$b_j$	Number of blocks in $j$ th BB combination
R	Incidence matrix of the design
S	Design matrix
$\gamma$	Eigen values of "S" matrix
$\underline{m}$	Vector of allocation of the control

## CHAPTER 1.

### 1.1 INTRODUCTION:

In many areas of Science it is desirable to compare simultaneously several test-treatments (or varieties) with a standard or existing one, called the control. Also, many practical situations require the blocking of experimental units in order to cut down on bias and improve the precision of the experiment. The most suitable designs are the Randomized Block Designs which require a block size equal to the number of treatments (including the control) involved. If the total number of treatment is large, then a large block is required to accommodate all the treatments. This create practical difficulties in handling the experiment, as the size of experiment increases considerably. Incomplete block designs are therefore suggested, in which size of the block is less than the number of treatments. Balanced Incomplete Block (BIB) Designs are well known designs of this type, but they are not appropriate for comparing test-treatments with the control, as they assign equal weight for all treatment comparisons, including the control; whereas, more efficiency is required for the treatment-control comparison. One way out from this

problem is to construct BIB designs and add equal number of control to each block. Such a design is called augmented or reinforced design and is more efficient than BIB designs for comparing treatments with a control, but it is less efficient for pair wise comparisons between treatments. Such designs are studied by Das(1958), Constantine(1983) and many other authors.

A new general class of designs is proposed by Bechhofer and Tamhane (1981), that is appropriate for the above problem. They are called Balanced Treatment Incomplete Block ( BTIB ) Designs and are balanced with respect to treatment other than the control. They study the structure of these designs and give the method of their construction. A lot of work has been done in this area since then. Ture (1982) studied in detail the construction of these designs for a variety of parameter range and also discussed the A-optimality of such designs. Bechhofer and Tamhane (1983) focus their attention on the problem of optimal allocation of experimental units among the test-treatments and the control to minimize the total size of the experiment. Majumdar and Notz (1983) established optimal properties of some BTIB designs. Hedayat and Majumdar (1984) have an extensive study of A-optimal BTIB designs in blocks of size 2, through a combination of theoretical results and



numerical investigations. Hedayat and Majumdar (1985) investigated the optimality of augmented BIB designs and were successful in solving the problem analytically. Stufken (1987) extended the results obtained by Hedayat and Majumdar (1985), and discussed the A-optimality of R- (reinforced) and S- (Step) type designs.

Majumdar (1981) shows that one needs to consider only binary designs, in which each treatment appears at most once in a block, in order to search for A-optimal designs. Since treatments are equally replicated in binary BTIB designs, it follows that in many cases an exactly A-optimal design will not exist. In such a case one might try to construct a design which is as close to a BTIB design as possible. Here by being close to a BTIB design it is meant that a design is as equally replicated as possible, and each treatment-control pair in a block should occur (over the whole design) as equally as possible and the same should be true for the repetition of all treatment-treatment pairs.

This leads to the study of Most Balanced designs, the topic of research in the following pages of this dissertation.

**1.2 PRELIMINARIES:**

**1.2.1 GENERAL BLOCK DESIGNS:**

Consider "b" blocks of more or less homogeneous experimental units or plots. Suppose j<sup>th</sup> block consists of k<sub>j</sub> plots. The design is said to be an incomplete block design if the number of treatments "v", is greater than the number of plots in a block. There are  $n = \sum_{j=1}^b k_j$  experimental units. The usual model is

$$Y_{iju} = \gamma + \alpha_i + \beta_j + e_{iju}, \quad \begin{matrix} i = 1, 2, \dots, v \\ j = 1, 2, \dots, b \\ u = 1, 2, \dots, k_j \end{matrix}$$

where i<sup>th</sup> treatment has been applied to u<sup>th</sup> plot belonging to j<sup>th</sup> block.

$Y_{iju}$  is the response from such a plot

$\alpha_i$  is the i<sup>th</sup> treatment effect

$\beta_j$  is the j<sup>th</sup> block effect

$e_{iju}$  is the residual of the plot

Consider the incidence matrices for treatments and blocks,

$$\phi = (\phi_1, \phi_2, \dots, \phi_v)$$

and  $\psi = (\psi_1, \psi_2, \dots, \psi_b)$

where,  $\phi_i = (\phi_{i1}, \phi_{i2}, \dots, \phi_{in})$

and  $\psi_j = (\psi_{j1}, \psi_{j2}, \dots, \psi_{jn})$

where,  $\phi_{iu} = \begin{cases} 1 & , \text{ if } i \text{ th treatment allotted to} \\ & u \text{ th plot} \\ 0 & , \text{ otherwise} \end{cases}$

and,  $\psi_{ju} = \begin{cases} 1 & , \text{ if plot } u \text{ receiving treatment} \\ & i \text{ is in block } j \\ 0 & , \text{ otherwise} \end{cases}$

The above linear model can be written as

$$y = \gamma \cdot J_n + \phi \alpha + \beta \cdot \psi + e$$

where,  $y$  is  $n \times 1$  vector of observations

$J_n$  is  $n \times 1$  vector of unity

$\alpha$  is  $v \times 1$  vector of treatment effects

$\beta$  is  $b \times 1$  vector of block effects

$e$  is  $n \times 1$  vector of residuals

The grand total is  $G = J_n' y$

Treatment total is  $T = (T_1, T_2, \dots, T_v)' = \phi' \cdot y$

Block totals is  $B = (B_1, B_2, \dots, B_b)' = \psi' \cdot y$

The normal equations of the least-square estimation are

$$\begin{bmatrix} J_n' \\ \phi' \\ \psi' \end{bmatrix} \cdot \begin{bmatrix} J_n & \phi & \psi \end{bmatrix} \cdot \begin{bmatrix} \hat{\gamma} \\ \hat{\alpha} \\ \hat{\beta} \end{bmatrix} = \begin{bmatrix} J_n' \\ \phi' \\ \psi' \end{bmatrix} \cdot y$$

that is  $\begin{bmatrix} n & r' & k' \\ r & D_r & R \\ k & R' & D_k \end{bmatrix} \cdot \begin{bmatrix} \hat{\gamma} \\ \hat{\alpha} \\ \hat{\beta} \end{bmatrix} = \begin{bmatrix} G \\ T \\ B \end{bmatrix}$

where,  $r = (r_1, r_2, \dots, r_v)'$

$r_i$  is the total number of times  $i$  th treatment is repeated in the design.

$$k = (k_1, k_2, \dots, k_b)'$$

$D_r = \text{diag}(r_i)$  is a  $v \times v$  diagonal matrix of  $r_i$ 's and off diagonal elements zeros.

$D_k = \text{diag}(k_j)$  is a  $b \times b$  diagonal matrix of  $k_j$ 's and off diagonal elements zeros.

$R = (r_{ij})$  is a  $v \times b$  matrix of elements,  $r_{ij}$  = number of times  $i$  th treatment appears in  $j$  th block. There are  $v + b + 1$  linear equations as a whole, but they are linearly dependent. Indeed,

$$\begin{aligned} n \hat{\gamma} + r \hat{\alpha} + k \hat{\beta} &= J'_v (r \hat{\gamma} + D_r \hat{\alpha} + R \hat{\beta}) \\ &= J'_b (k \hat{\gamma} + R \hat{\alpha} + D_k \hat{\beta}) \end{aligned}$$

Thus we consider only two equations,

$$r \hat{\gamma} + D_r \hat{\alpha} + R \hat{\beta} = T \quad 1.2.1$$

$$k \hat{\gamma} + R \hat{\alpha} + D_k \hat{\beta} = B \quad 1.2.2$$

Premultiply eq.1.2.2 by  $R D_k^{-1}$  and subtract from eq.1.2.1 to obtain,

$$(D_r - R D_k^{-1} R') \hat{\alpha} = T - R D_k^{-1} B \equiv Q = (Q_i) \text{ (say)}$$

where,

$$Q_i = T_i - \sum_{j=1}^b r_{ij} \cdot \frac{B_j}{k_j}$$

is called the adjusted (adjusted for blocks) treatment totals, and  $\sum_{i=1}^v Q_i = 0$ .

Let,  $M = D_r - R D_k^{-1} R' = (m_{il})$ ;  $i, l = 1, 2, \dots, v$ .

Then,

$$M \hat{\alpha} = Q \quad 1.2.3$$

### 1.2.2 INFORMATION MATRIX:

In eq.(3), above ,  $M = ( m_{ij} )$  is called Information Matrix for estimating  $\alpha$ . Here,

$$m_{ii} = r_i - \sum_{j=1}^b r_{ij}^2 / k_j$$

and,

$$m_{it} = - \sum_{j=1}^b (r_{ij} \times r_{tj}) / k_j , \quad i \neq t; \quad i, t = 1, 2, \dots, v$$

$M^{-1}$  is the variance-covariance matrix, whose diagonal elements give the variance of the estimates and off diagonal elements give the covariances.

### 1.2.3 PROPER AND CONNECTED DESIGNS AND ESTIMABILITY:

In general all treatment effects cannot be estimated. However, it is easy to show that if a linear function  $a \alpha$  is estimable, then this must be a treatment contrast e.g.  $\alpha_i - \alpha_t$  for all  $i$  and  $t$ .

A design is called a proper design if all blocks have the same number of plots  $k$ , i.e.  $k_1 = k_2 = \dots = k_p = k$ . Bechhofer and Tamhane (1981) show that for a proper design with parameters  $b$ ,  $k$  and  $v$  (total  $v+1$  treatments including the control), the information

matrix for estimating all contrasts  $\alpha_0 - \alpha_i$  ( $i = 1, 2, \dots, v$ ) is a  $v \times v$  symmetric non-negative definite matrix  $M = (m_{it})$  such that,

$$m_{it} = \begin{cases} r_i - \frac{1}{k} \sum_{j=1}^b r_{ij}^2 & , \text{ if } i = t \\ -\frac{1}{k} \lambda_{it} & , \text{ if } i \neq t \end{cases}$$

where,

$\lambda_{it}$  = number of times  $i$ th and  $t$ th treatments appears together in a block over the whole design.

In equireplicate design each treatment is repeated equal number of times over the whole design i.e.  $r_1 = r_2 = \dots = r_v = r$

In binary design a treatment appears at most once in a block i.e.  $r_{ij} = 0$  or  $1$ . Therefore,  $r_{ij}^2 = r_{ij}$ , and the information matrix  $M = (m_{it})$  is given by,

$$m_{it} = \begin{cases} r_i - \frac{1}{k} \sum_{j=1}^b r_{ij} & , \text{ if } i = t \\ -\frac{1}{k} \lambda_{it} & , \text{ if } i \neq t \end{cases}$$

A treatment and a block are said to be associated if the treatment occurs in the block. A design is connected if all the treatments and block are associated with each other. For a connected design every

treatment contrast is estimable and

$$\text{rank} ( M ) = \text{number of treatments} - 1$$

### 1.3 BALANCED TREATMENT INCOMPLETE BLOCK (BTIB) DESIGNS:

#### 1.3.1 DEFINITION:

Consider an incomplete block design in "v" treatments and a control denoted by "0" in "b" blocks of size "k" (  $v \geq k$  ) each

The design is a BTIB design if,

$$\lambda_{0i} = \lambda_0, \quad i = 1, 2, \dots, v$$

$$\text{and} \quad \lambda_{it} = \lambda_1, \quad i, t = 1, 2, \dots, v; i \neq t$$

Hence, a BTIB design is such that each test-treatment appears with control ( in a block ) same number of times over the design, and any pair of test-treatments appears together ( in a block ) the same number of times over the design.

#### 1.3.2 INFORMATION MATRIX:

The information matrix of a BTIB design is given by ( see eq. A.1 of Bechhofer and Tamhane 1981),

$$m_{it} = \begin{cases} \frac{1}{k} \{ \lambda_0 + (v-1) \cdot \lambda_1 \}, & i = t, i = 1, 2, \dots, v \\ -\frac{1}{k} \lambda_1, & i \neq t; i, t = 1, 2, \dots, v \end{cases}$$

So the information matrix is completely symmetric. Since,  $M^{-1}$  is the variance-covariance matrix of the vector of estimates  $(\hat{\alpha}_0 - \hat{\alpha}_1, \dots, \hat{\alpha}_0 - \hat{\alpha}_v)$ , a BTIB design provides BLUE 's  $\hat{\alpha}_0 - \hat{\alpha}_i$  ( $i = 1, 2, \dots, v$ ), where,

$$\hat{\alpha}_0 - \hat{\alpha}_i = \frac{\lambda_1 Q_0 - \lambda_0 Q_i}{\lambda_0 (\lambda_0 + v \lambda_1)}, \quad i = 1, 2, \dots, v$$

where,

$$Q_i = k \cdot T_i - \sum_{j=1}^b r_{ij} \cdot B_j$$

also,

$$V(\hat{\alpha}_0 - \hat{\alpha}_i) = \frac{k (\lambda_0 + \lambda_1)}{\lambda_0 (\lambda_0 + v \lambda_1)} \sigma^2 = \tau^2 \sigma^2$$

$$\text{Corr.}\{\hat{\alpha}_0 - \hat{\alpha}_i, \hat{\alpha}_0 - \hat{\alpha}_t\} = \frac{\lambda_1}{\lambda_0 + \lambda_1} = \rho, \quad \begin{matrix} i \neq t \\ i, t = 1, 2, \dots, v \end{matrix}$$

$$\text{Cov.}\{\hat{\alpha}_0 - \hat{\alpha}_i, \hat{\alpha}_0 - \hat{\alpha}_t\} = \frac{k \lambda_1}{\lambda_0 (\lambda_0 + v \lambda_1)}, \quad \begin{matrix} i \neq t \\ i, t = 1, 2, \dots, v \end{matrix}$$

### 1.3.3 MINIMAL COMPLETE CLASS OF GENERATOR DESIGNS(MCCGD):

For given  $(v, k)$ , a generator design is a BTIB design no proper subset of whose blocks forms a BTIB design, and no block of which contains only one of  $v+1$  treatments.

Generator designs are important for the



construction of BTIB designs, as any design is either a generator design or a union of copies of generator designs. In fact, if for given  $(v,k)$  there are "n" generator designs  $D_i$  ( $i = 1, 2, \dots, n$ ) where  $D_i$  has parameters  $(b_i, \lambda_0^{(i)}, \lambda_1^{(i)})$ ,  $i = 1, 2, \dots, n$ ; then the BTIB design  $D = \bigcup_{i=1}^n f_i \cdot D_i$  obtained by taking union of  $f_i$  replications of  $D_i$  has parameters  $(b, \lambda_0, \lambda_1)$  given by,

$$b = \sum_{i=1}^n f_i b_i, \quad \lambda_0 = \sum_{i=1}^n f_i \lambda_0^{(i)}, \quad \lambda_1 = \sum_{i=1}^n f_i \lambda_1^{(i)}$$

Notz and Tamhane (1981) show that for a given  $(v,k)$  there exists only finitely many generator designs.

For given  $(v,k)$  a design  $D_2$  with parameters  $(b_2, \lambda_0^{(2)}, \lambda_1^{(2)})$  is inadmissible with respect to  $D_1$  with parameters  $(b_1, \lambda_0^{(1)}, \lambda_1^{(1)})$  iff  $b_1 \leq b_2$ ,  $\tau_1^2 \leq \tau_2^2$ ,  $\rho_1 \geq \rho_2$  with at least one inequality strict. The  $D_1$  and  $D_2$  are equivalent if  $b_1 = b_2$ ,  $\lambda_0^{(1)} = \lambda_0^{(2)}$ ,  $\lambda_1^{(1)} = \lambda_1^{(2)}$ .

A smallest set of generator designs with the property that any admissible design can either be constructed from the set or it is equivalent to a design which is constructed from the set, is called the Minimal Complete Class of Generator Designs (MCCGD).

1.4 PARTIALLY BALANCED TREATMENT  
INCOMPLETE BLOCK DESIGNS:

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1.4.1 PARTIALLY BALANCED SCHEME:

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Extensive studies have been done in the area of finding BTIB designs and catalog of these designs for values of  $v$ ,  $b$ , and  $k$  in the practical range  $2 \leq k \leq v \leq 10$  and  $v \leq b \leq 50$ , has been published by Hedayat and Majumdar (1984, 1985). However, there are good many values of  $v$ ,  $b$ , and  $k$  in the practical range where no BTIB design exists. This is due to the severe restriction of complete symmetry of information matrix (i.e. all diagonal elements are equal and all off diagonal elements are equal). In other words, if the restrictions that,

$$\lambda_{0i} = \lambda_0, \quad i = 1, 2, \dots, v \text{ and}$$
$$\lambda_{it} = \lambda_1, \quad i, t = 1, 2, \dots, v; i \neq t$$

can be relaxed in some way, there is hope of finding a more general class of designs having members existing even for those values of  $v$ ,  $b$ , and  $k$  for which BTIB design do not exist. The following idea is therefore, proposed.

Take,  $\lambda_{0i} = \lambda_0, \quad i = 1, 2, \dots, v$

$$\lambda_{it} = \begin{cases} \lambda_1, & \text{for some pairs } i \text{ \& } t, i \neq t \\ \lambda_2, & \text{for some other pairs } i \text{ \& } t, i \neq t \end{cases}$$

Note that the information matrix of such a design is still symmetrical, but not completely ( all diagonal elements are equal but off diagonal elements are different).

Before going to discuss in detail the construction and existence of such designs, in next chapter, it is necessary to revise some old and to give some new preliminary definitions.

DEFINITION 1: ( BIB designs )

An arrangement of "v" treatments in "b" blocks of "k" experimental units each is called a Balanced Incomplete Block (BIB) Design if it satisfies the following three conditions:

- (i) each block contains k different treatments.
- (ii) each treatment occurs in r blocks, and
- (iii) any two treatments occurs together in  $\lambda$  blocks

DEFINITION 2: ( PBIB Designs )

Given v treatments, a relation satisfying the following conditions is said to be an association scheme with m associate classes:

- (i) any two treatments are either first or second or,....., m th associates.
- (ii) any treatment has  $n_i$  i th associates,  $i = 1, 2, \dots, m$ .
- (iii) if two treatments are i th associates, then

the number of treatments that are  $j$  th associates of the first and  $k$  th associates of the second is  $p_{jk}^i$  ( $i, j, k = 1, 2, \dots, m$ ).

A design is said to be partially balanced incomplete block design if the following conditions are satisfied:

- (i) each block contains  $k$  different treatments.
- (ii) each treatment occurs in  $r$  blocks.
- (iii) any two treatments that are  $i$  th associates occurs together in  $\lambda_i$  blocks,  $i = 1, 2, \dots, m$ .

DEFINITION 3: ( GD PBIB Designs )

This is a special case of PBIB designs with  $m = 2$ . Let there be  $v = m.n$  treatments, divided into " $m$ " groups of " $n$ " treatments each in the following way.

		T R E A T M E N T S			
GROUP 1	1	2	3	.....	n
2	n+1	n+2	n+3	.....	2n
⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮
m	(m-1)n+1	(m-1)n+2	(m-1)n+3	.....	mn

Treatments belonging to the same group are first associates, and treatments belonging to different groups are the second associates. This association scheme is called a Group Divisible association leading to

group-divisible partially balanced incomplete block designs ( GD PBIBD ). So, any two treatments that are in the same group occur together in  $\lambda_1$  blocks and any two treatments that are in different groups occur together in  $\lambda_2$  blocks.

DEFINITION 4: ( BTIB Designs )

A design in  $v+1$  treatments  $0, 1, 2, \dots, \dots, v$ , ( 0 is the control ) in  $b$  blocks of size  $k$  each (  $v \geq k$  ) is a balanced treatment incomplete block design if,

$$\lambda_{0i} = \lambda_0, \quad i = 1, 2, \dots, v$$

$$\lambda_{it} = \lambda_1, \quad i, t = 1, 2, \dots, v; i \neq t$$

DEFINITION 5: ( PBTIB Designs )

Consider  $v$  test-treatments  $1, 2, \dots, v$  and a control 0. Apply the association scheme described in Definition 2 of partially incomplete block designs on test-treatments only. A design is partially balanced treatment incomplete block design if any two treatments that are  $i$ th associate occur together in  $\lambda_i$  blocks,  $i = 1, 2, \dots, m$ .

DEFINITION 6: ( GD TIB Designs )

Let  $v = m.n$ , and  $v+1$  treatments  $0, 1, 2, \dots, v$ , can be divided in  $m+1$  groups  $0, 1, 2, \dots, m$ , such that group 0 contains control treatment only and the remaining  $m$  groups are as given in

**Definition 3** of GD PBIB designs. Two test-treatments, if they are in the same group are first associates, and if they are in different groups are second associates. A design is group divisible treatment incomplete block design if,

$$\lambda_{0i} = \lambda_0, \quad i = 1, 2, \dots, v$$

$$\lambda_{it} = \begin{cases} \lambda_1, & i \text{ \& t are first associates, } i \neq t \\ \lambda_2, & i \text{ \& t are second associates, } i \neq t \end{cases}$$

#### 1.4.2 MOST BALANCED SCHEME:

##### DEFINITION 1: ( MB GD IB Designs )

If a design is GD PBIB design with 2 groups and  $\lambda_2 = \lambda_1 \pm 1$  it is called Most Balanced Group Divisible Incomplete Block Design ( MB GD IBD ). Furthermore, a design is binary MB GD IB design if each treatment has the same number of replications and appears in each block at most once ( Cheng 1978).

##### DEFINITION 2: ( MB GD TIB Designs )

A GD TIB design in Definition 6 of Section 1.4.1 is Most Balanced Group Divisible Treatment Incomplete Block Design ( MB GD TIBD ) if  $m = 2$  and  $\lambda_2 = \lambda_1 \pm 1$ , and it is binary if each test-treatment is

equi-replicated and occurs in each block at most once. Note that no restriction is imposed upon the occurrence of control in blocks.

**DEFINITION 3: ( E MB GD TIB Designs )**

An extended family of MB GD TIB Designs is defined as a family of designs with  $m = 2$  and  $\lambda_2 = \lambda_1 + \delta$ , where  $\delta = -1, 0, 1$ . In this way it contains all BIB designs, BTIB designs and MB GD TIB designs.

**EXAMPLE 1:** The following design is a member of extended family with  $k = 3, v = 4, b = 10$

BLOCK		1	2	3	4	5	6	7	8	9	10
PLOT	1	0	0	0	0	0	0	1	1	1	2
	2	1	1	1	2	2	3	2	2	3	3
	3	2	3	4	3	4	4	3	4	4	4

For this allocation,

$$r_0 = r_1 = r_2 = r_3 = r_4 = r = 6$$

$$\lambda_{01} = \lambda_{02} = \lambda_{03} = \lambda_{04} = \lambda_0 = 3$$

$$\lambda_{12} = \lambda_{13} = \lambda_{14} = \lambda_{23} = \lambda_{24} = \lambda_{34} = \lambda_1 = 3$$

Since,  $\lambda_0 = \lambda_1$  this is a BIB design.

**EXAMPLE 2:** Let  $k = 3, v = 4, b = 6$

BLOCK		1	2	3	4	5	6
PLOT	1	0	0	0	0	0	0
	2	1	1	1	2	2	3
	3	2	3	4	3	4	4

Here,  $r_0 = 6, r_i = r = 3, i = 1, 2, 3, 4$

$$\lambda_{0i} = \lambda_0 = 3, \quad i = 1, 2, 3, 4$$

$$\lambda_{it} = \lambda_1 = 1, \quad i < t, i, t = 1, 2, 3, 4$$

This is a BTIB design and hence, a member of the family

**EXAMPLE 3:** Let  $k = 3, v = 4, b = 8$

BLOCK		1	2	3	4	5	6	7	8
PLOT	1	0	0	0	0	0	0	1	1
	2	1	2	2	2	0	0	2	3
	3	4	3	4	4	1	3	3	4

This is a MB GD TIB design with the following allocation scheme.

$$m = 2, n = 2$$

group 0 contains control

group 1 contains treatments 1 and 2

group 2 contains treatments 3 and 4

Hence, (1,2) and (3,4) are first associates and (1,3),

(1,4), (2,3), (2,4) are second associates.

Here,  $r_0 = 8, r_i = r = 4, i = 1, 2, 3, 4$

$$\lambda_{0i} = \lambda_0 = 3, \quad i = 1, 2, 3, 4$$

$$\lambda_{it} = \begin{cases} \lambda_1 = 1, & i \neq t \text{ are first associates} \\ \lambda_2 = 2, & i \neq t \text{ are second associates} \end{cases}$$



DEFINITION 4: ( Basic Blocks )

A block that contains control is called a "basic block" (BB). A BB is identified by the number of TC-pairs in it. When  $k = 4$ , for example, there are two different types of BB 's.

Type 1:  $\begin{bmatrix} 0 \\ 0 \\ x \\ y \end{bmatrix}$  referred as 4-BB, since it contains 4 TC-pairs viz.  $(0,x), (0,y), (0,x), (0,y)$ , where  $x, y$  are test-treatments.

Type 2:  $\begin{bmatrix} 0 \\ x \\ y \\ z \end{bmatrix}$  referred as 3-BB, since it contains 3 TC-pairs viz.  $(0,x), (0,y), (0,z)$ , where  $x, y, z$  are test-treatments.

DEFINITION 5: ( Basic Block Combinations )

Given  $(v,k)$  and an E MB GD TIB design with parameters  $(b, \lambda_0, \lambda_1, \lambda_2)$ , the total number of TC-pairs are  $v \lambda_0$  which can be allocated to different possible combinations of BBs, where each such combination must have  $v \lambda_0$  such pairs.

A BB combination can be represented by an ordered pair. For example, if  $k = 4, v = 4$ , a BB combination is represented by an ordered pair  $(U, V)$ , which means that there are  $U$  of 4-BBs and  $V$  of 3-BBs in the combination.

## CHAPTER 2

### 2.1 THEORETICAL RESULTS:

#### 2.1.1 THEORETICAL FACTS:

consider a design  $D$  with  $(k, v, b)$  and

$\lambda_{0i}$  = number of times treatment  $i$  and control are paired in a block.

$\lambda_{it}$  = number of times treatment  $i$  and treatment  $t$  ( $i \neq t$ ) are paired in a block.

Let control appears " $l$ " times in " $m_1$ " blocks

#### FACT 2.1

Given  $(k, v)$  and an E MB GD TIB design with parameters  $b, \lambda_0, \lambda_1, \lambda_2$ , the control appears a total of  $v \cdot \lambda_0$  times with the test-treatments (TC-pairs) i.e.

$$\sum_{i=1}^v \lambda_{0i} = v \lambda_0$$

#### FACT 2.2

Given  $(k, v)$  and an E MB GD TIB design with parameters  $b, \lambda_0, \lambda_1, \lambda_2$ , the number of times a test-treatment is paired with other test-treatments is equal to

$$\frac{1}{2} \left[ (v - 2) \lambda_1 + v \lambda_2 \right]$$

Consider the following association scheme,

group 1 contains treatments 1,2,.....,v/2

group 2 contains treatments (v/2)+1,(v/2)+2,.....,v

So,

number of times a test-treatment is paired with other test-treatments of the same group (i.e. first associates) =  $\left( \frac{v}{2} - 1 \right) \lambda_1$

number of times a test-treatment is paired with other test-treatments of different group (i.e. second associates) =  $(v/2) \lambda_2$

The sum of these two gives the result.

FACT 2.3

Given  $(k,v)$  and an E MB GD TIB design with parameters  $b, \lambda_0, \lambda_1, \lambda_2$ , the total number of Treatment- Treatment (test) pairs (TT-pairs) is equal to

$$\frac{v}{4} \left[ (v - 2) \lambda_1 + v \lambda_2 \right]$$

Actually,

$$\begin{aligned} \sum_{\substack{i,t \\ \text{first} \\ \text{assoc.}}} \lambda_{it} &= \text{no. of TT-pairs (first associate type)} \\ &= 2 \cdot \left[ \frac{v/2}{2} \right] \lambda_1 = \frac{1}{4} v (v - 2) \lambda_1 \end{aligned}$$

$$\sum_{\substack{i,t \\ \text{second} \\ \text{assoc.}}} \lambda_{it} = \text{no. of TT-pairs (second associate type)}$$

$$= \begin{bmatrix} v/2 \\ 1 \end{bmatrix} \begin{bmatrix} v/2 \\ 1 \end{bmatrix} \lambda_2 = \frac{1}{4} v^2 \lambda_2$$

The sum of these two gives the result.

**FACT 2.4**

A binary E MB GD TIB design is equi-replicate in terms of test-treatments, with the common replicate size "r" given by,

$$r = ( b k - r_0 ) / v \quad 2.1.1$$

**FACT 2.5**

For a design with parameters k, v, b,

$$(i) \quad r_0 = \sum_{I=1}^{k-1} 1 m_1 \quad 2.1.2$$

$$(ii) \quad \sum_{i=1}^v \lambda_{0i} = r_0 ( k - 1 ) - \sum_{I=1}^{k-1} 1 ( 1 - 1 ) m_1 \quad 2.1.3$$

$$(iii) \quad \sum_{\substack{i,t=1 \\ i \neq t}}^v \lambda_{it} = \frac{1}{2} b k ( k - 1 ) - \frac{1}{2} \sum_{I=1}^{k-1} 1 ( 1 - 1 ) m_1 - \sum_{i=1}^v \lambda_{0i} \quad 2.1.4$$

For proof see Ture ( 1982 ).

**EXAMPLE:** Consider the following MB GD TIB design with the following association scheme,

group 0 contains control

group 1 contains treatments 1,2

group 2 contains treatments 3,4

BLOCK		1	2	3	4	5	6
PLOT	1	0	0	1	1	1	2
	2	1	3	2	2	3	3
	3	2	4	3	4	4	4

Here,  $k = 3, v = 4, b = 6$

$$\lambda_0 = 1, \lambda_1 = 3, \lambda_2 = 2$$

$$m_1 = 2, m_2 = 0$$

See that,

$$r_0 = 2 = \sum_{i=1}^k 1 m_1 \quad (\text{FACT 2.5(i) verified})$$

$$\sum_{i=1}^4 \lambda_{0i} = 4 \quad (\text{FACT 2.1 and FACT 2.5(ii) verified})$$

$$\sum_{\substack{i,t=1 \\ i \neq t}}^4 \lambda_{it} = 14 \quad (\text{FACT 2.2 and FACT 2.5(iii) verified})$$

FACT 2.6 ( Ture 1982 )

For  $v = 4, k = 4$  the BB combinations are given by (see Definition 5 of section 1.4.2).

(i)  $( U = \lambda_0 - 3j + 3, V = 4j - 4 ), j = 1, 2, \dots$

$\dots, [\lambda_0/3] + 1, \text{ if } \lambda_0 \text{ is even} \quad \text{with}$

$M_j = \lambda_0 + 9j - 9 \text{ and } b_j = \lambda_0 + j - 1, \quad \text{where}$

$[x]$  is the largest integer smaller than  $x$ .

(ii) (  $U = \lambda_0 - 3j$  ,  $v = 4j$  ),  $j = 1, 2, \dots, [\lambda_0/3]$   
 if  $\lambda_0 (> 1)$  is odd with  $M_j = \lambda_0 + 9j$  and  
 $b_j = \lambda_0 + j$

Where  $M_j$  is the number of TT-pairs in the  $j$  th BB combination and  $b_j$  is the number of blocks it contains.

FACT 2.7 ( Ture 1982)

Consider all possible BB combinations for a specified  $(v, k, \lambda_0, \lambda_1, \lambda_2)$ , then

$$b \geq b_{j_0}$$

where  $j_0$  is the index of the BB combination with  $b_{j_0}$  BBS for which  $|N - M_{j_0}|$  is closest to zero;  $N$  is the total number of TT-pairs in the design and  $M_{j_0}$  is the number of TT-pairs in the  $j_0$  th BB combination.

2.1.2 NON-EXISTENCE OF DESIGNS:

THEOREM: An Extended Most Balanced Group Divisible Treatment Incomplete Block Design with parameters  $k, v, b, r, \lambda_0, \lambda_1,$  and  $\lambda_2$  is non-existent if,

$$b < \begin{cases} 2 & , \text{when } r = \lambda_1 > \lambda_2 \\ v - 1, & \text{when } \lambda_1 < \lambda_2, \lambda_1 < r, r = \lambda_1 + \frac{v(\lambda_2 - \lambda_1)}{2} \\ v, & \text{when } r > \lambda_1 \geq \lambda_2, \text{ or } \\ & \lambda_1 < \lambda_2, r > \lambda_1 + \frac{v(\lambda_2 - \lambda_1)}{2} \end{cases}$$

**PROOF:**

Let  $R = (r_{ij})$  be the incidence matrix of the  $(v+1 \times b)$

design, where  $r_{ij}$  = no. of times  $i$ th treatment is replicated in  $j$ th block;  $i = 0, 1, 2, \dots, v$ ;  $j = 1, 2, \dots, b$ .

Let,  $S = R R' = (s_{it})$ , with

$$s_{it} = \begin{cases} \sum_{j=1}^b r_{0j}^2 = s_{00}, & i = t = 0 \\ \lambda_0, & i = 0, t = 1, 2, \dots, v \\ r, & i = t; i, t = 1, 2, \dots, v \\ \lambda_1, & i \text{ and } t \text{ are first associates} \\ \lambda_2, & i \text{ and } t \text{ are second associates} \end{cases}$$

Note that  $r \geq \lambda_1, \lambda_2$

Let,  $\nu_0, \nu_1, \dots, \nu_v$  are  $v+1$  eigen values of  $S$ , given by the following equation,

$$| S - \nu I | = 0$$

where,

$$[S - \nu I] = \begin{bmatrix} s_{00} - \nu & \lambda_0 & \dots & \lambda_0 & \lambda_0 & \dots & \lambda_0 \\ \lambda_0 & r - \nu & \dots & \lambda_1 & \lambda_2 & \dots & \lambda_2 \\ \vdots & \vdots & & \vdots & \vdots & & \vdots \\ \lambda_0 & \lambda_1 & \dots & r - \nu & \lambda_2 & \dots & \lambda_2 \\ \lambda_0 & \lambda_2 & \dots & \lambda_2 & r - \nu & \dots & \lambda_1 \\ \vdots & \vdots & & \vdots & \vdots & & \vdots \\ \lambda_0 & \lambda_2 & \dots & \lambda_2 & \lambda_1 & \dots & r - \nu \end{bmatrix}$$

The equation  $|S - \nu I| = 0$  reduces to the following polynomial equation,

$$(r - \nu - \lambda_1)^{\nu-2} \left[ r - \nu + \left(\frac{\nu}{2} - 1\right) \lambda_1 - \frac{\nu}{2} \lambda_2 \right] \left[ (s_{00} - \nu) \left\{ r - \nu + \left(\frac{\nu}{2} - 1\right) \lambda_1 + \frac{\nu}{2} \lambda_2 \right\} - \nu \lambda_0^2 \right] = 0$$

Solution of this equation gives the following  $\nu + 1$  roots

$$\nu_0, \nu_1 = \frac{1}{2} \left[ \left\{ r + \left(\frac{\nu}{2} - 1\right) \lambda_1 + \frac{\nu}{2} \lambda_2 + s_{00} \right\} \pm \left[ \left\{ \left( r + \left(\frac{\nu}{2} - 1\right) \lambda_1 + \frac{\nu}{2} \lambda_2 \right)^2 - s_{00}^2 \right\} + 4 \nu \lambda_0^2 \right]^{1/2} \right]$$

At least one of these two roots is positive.

$$\nu_2, \dots, \nu_{\nu-1} = r - \lambda_1 = \mu \quad (\text{say})$$

$$\nu_\nu = r + \left(\frac{\nu}{2} - 1\right) \lambda_1 - \frac{\nu}{2} \lambda_2 = \mu + \frac{\nu}{2} (\lambda_1 - \lambda_2) = \mu^* \quad (\text{say})$$

Now,  $\text{rank}(R) = \text{rank}(S) \leq b$

Whereas, the  $\text{rank}(S)$  is equal to the number of positive eigen values of  $S$ , which can be obtained by considering the following possibilities.

$$\mu \begin{cases} = 0, & \text{if } r = \lambda_1 \\ > 0, & \text{if } r > \lambda_1 \end{cases}$$

Note that,

$$\mu^* = 0, \text{ if } \mu = 0 \text{ (i.e. } r = \lambda_1) \text{ and } \lambda_1 = \lambda_2, \text{ or}$$

$$\text{if } \mu > 0 \text{ (i.e. } r > \lambda_1), \lambda_1 < \lambda_2, \text{ and}$$

$$\mu = \frac{\nu}{2} (\lambda_2 - \lambda_1).$$



$\mu^* > 0$ , if  $\mu = 0$  (i.e.  $r = \lambda_1$ ) and  $\lambda_1 > \lambda_2$ , or  
 if  $\mu > 0$  (i.e.  $r > \lambda_1$ ) and  $\lambda_1 = \lambda_2$ , or  
 if  $\mu > 0$  (i.e.  $r > \lambda_1$ ) and  $\lambda_1 > \lambda_2$ , or  
 if  $\mu > 0$  (i.e.  $r > \lambda_1$ ),  $\lambda_1 < \lambda_2$ , and  

$$\mu > \frac{v}{2} (\lambda_2 - \lambda_1)$$

The following are, therefore, true

- (i)  $\text{rank}(R) \leq 2$ , if  $\mu = 0$ ,  $\mu^* = 0$  (i.e. if  $r = \lambda_1 = \lambda_2$ ), hence  $b \geq 1$
- (ii)  $\text{rank}(R) = 2$  or  $3$ , if  $\mu = 0$ ,  $\mu^* > 0$  (i.e. if  $r = \lambda_1 > \lambda_2$ ), hence  $b \geq 2$
- (iii)  $\text{rank}(R) = v - 1$  or  $v$ , if  $\mu > 0$ ,  $\mu^* = 0$  (i.e. if  $\lambda_1 < \lambda_2$ ,  $\lambda_1 < r$ , and  $r = \lambda_1 + \frac{v}{2} (\lambda_2 - \lambda_1)$ ), hence  $b \geq v - 1$
- (iv)  $\text{rank}(R) = v$  or  $v + 1$ , if  $\mu > 0$ ,  $\mu^* > 0$  (i.e.  $r > \lambda_1 \geq \lambda_2$  or  $\lambda_1 < \lambda_2$ , and  $r > \lambda_1 + \frac{v}{2} (\lambda_2 - \lambda_1)$ ), hence  $b \geq v$

Possibility (i) is trivial. The converse of possibilities (ii), (iii), and (iv) give the conditions required for the non-existence of the design in the theorem. Hence, the theorem is proved.

**REMARK:** Although the theorem is completely valid to discard a design as non-existent, it does not provide sufficient conditions for the existence of a design i.e. a design may not be constructed even if its non-existence is not proved. However, the theorem is extremely useful as it eliminates the useless effort of trying to construct a

design which is actually non-existent. It is our experience that out of hundreds of designs studied there are very few designs which disobey the condition of non-existence but even the construction of them is not possible.

The usefulness of this theorem is explained by the following examples.

EXAMPLE 1: Consider a BTIB design with parameters

$$k = 4, v = 4, b = 4, r = 3, \lambda_0 = 3, \lambda_1 = 1, \lambda_2 = 1$$

The following relation holds,

$$r > \lambda_1 = \lambda_2$$

According to the possibility (iv),  $b$  should be greater than or equal to  $v$  for the possible existence of the design. Since, it is so, the design may exist.

EXAMPLE 2: Consider a MB GD TIB design with,

$$k = 3, v = 4, b = 8, r = 4, \lambda_0 = 3, \lambda_1 = 1, \lambda_2 = 2$$

Here,  $\lambda_1 < \lambda_2$  and  $r > \lambda_1 + \frac{v}{2} (\lambda_2 - \lambda_1) = 3$ .

This implies that  $b \geq v = 4$  for the possible existence.

The design may exist.

EXAMPLE 3: Consider a BIB design with parameters,

$$k = 4, v = 6, b = 4, r = 2, \lambda_0 = 1, \lambda_1 = 1, \lambda_2 = 1.$$

Since,  $r > \lambda_1 = \lambda_2$  implies that  $b \geq v = 6$  for the possible existence, the design cannot exist.

**EXAMPLE 4:** Let us have a MB GD TIB design with,

$$k = 3, v = 8, b = 7, r = 2, \lambda_0 = 1, \lambda_1 = 1, \lambda_2 = 0$$

The relation  $r > \lambda_1 > \lambda_2$  gives the condition  $b \geq 8$  for the possible existence. This is not true here, because  $b = 7$ . Hence, the design will not exist.

### 2.1.3 INFORMATION MATRIX AND RELATED MEASURES:

The information matrix of a GD TIB design with parameters  $k, v, b, \lambda_0, \lambda_1,$  and  $\lambda_2$  is given by,

$$M = (m_{it})_{(vxv)}$$

where,

$$m_{it} = \begin{cases} \frac{1}{k} \left( \lambda_0 + \frac{v-2}{2} \lambda_1 + \frac{v}{2} \lambda_2 \right), & i = t \\ -\frac{1}{k} \lambda_1, & i \text{ and } t \text{ are first associates} \\ -\frac{1}{k} \lambda_2, & i \text{ and } t \text{ are second associates} \end{cases}$$

Let,  $\alpha = \frac{1}{k} \left\{ \lambda_0 + \frac{v-2}{2} + \frac{v}{2} \lambda_2 \right\}, \beta = -\frac{1}{k} \lambda_1, \gamma = -\frac{1}{k} \lambda_2$

So, the information matrix  $M$  can be written as

$$M = \left[ \begin{array}{c|c} P & Q \\ \hline Q & P \end{array} \right]$$

where,

$$P = \left[ \begin{array}{cccc} \alpha & \beta & \dots & \beta \\ \beta & \alpha & \dots & \alpha \\ \vdots & \vdots & \ddots & \vdots \\ \alpha & \alpha & \dots & \beta \end{array} \right], \text{ and } Q = \gamma \cdot J = \gamma \cdot \left[ \begin{array}{cccc} 1 & 1 & \dots & 1 \\ 1 & 1 & \dots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \dots & 1 \end{array} \right]$$

The variance-covariance matrix for estimating  $(\hat{\alpha}_0 - \hat{\alpha}_1)$ ,  
 $i = 1, 2, \dots, v$  is given by

$$V = \sigma^2 M^{-1}$$

where,

$$M^{-1} = \left[ \begin{array}{c|c} P^{-1} ( I + Q S^{-1} Q P^{-1} ) & - P^{-1} Q S^{-1} \\ \hline - S^{-1} Q P^{-1} & S^{-1} \end{array} \right]$$

where,  $S = P - Q P^{-1} Q$ .

The four elements in  $M^{-1}$  are as follows,

(i)  $P^{-1} ( I + Q S^{-1} Q P^{-1} ) = ( p_{ij} )$ , where

$$p_{ij} = \begin{cases} A \{ \alpha + ( \frac{v}{2} - 2 ) \beta \} + E, & \text{if } i = j \\ - \beta A + E & , \text{if } i \neq j \end{cases}$$

(ii), (iii)  $- P^{-1} Q S^{-1} = - S^{-1} Q P^{-1} = -\gamma (\alpha - \beta) (d - e) \text{ADJ}$

(iv)  $S^{-1} = D ( q_{ij} )$ , where

$$q_{ij} = \begin{cases} d + ( \frac{v}{2} - 2 ) e & , \text{if } i = j \\ - e & , \text{if } i \neq j \end{cases}$$

Where,  $A = (\alpha - \beta)^{-1} \{ \alpha + ( \frac{v}{2} - 1 ) \beta \}^{-1}$

$D = (d - e)^{-1} \{ d + ( \frac{v}{2} - 1 ) e \}^{-1}$

$E = \frac{v}{2} \gamma^2 (\alpha - \beta)^2 (d - e) A^2 D$

$d = \alpha - \frac{v}{2} \gamma^2 (\alpha - \beta) A$

$e = \beta - \frac{v}{2} \gamma^2 (\alpha - \beta) A$

The diagonal elements of V, gives the variance of the estimates.

$$v(\hat{\alpha}_0 - \hat{\alpha}_i) = \frac{\sigma^2 k \{ (\lambda_0 + \lambda_1 + \frac{v}{2} \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_2) - \frac{v}{2} (\frac{v}{2} - 1) \lambda_2^2 \}}{\lambda_0 (\lambda_0 + v \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_1 + \frac{v}{2} \lambda_2)}$$

$$= \eta^2 \sigma^2 \quad (\text{say}) \quad , \quad \text{for } i = 1, 2, \dots, v$$

The off diagonal elements of V give the covariance terms

$$\text{cov}\{(\hat{\alpha}_0 - \hat{\alpha}_i), (\hat{\alpha}_0 - \hat{\alpha}_t)\} = \begin{cases} \frac{\sigma^2 k \{ \lambda_1 + (\lambda_0 + \frac{v}{2} \lambda_2) + \frac{v}{2} \lambda_2^2 \}}{\lambda_0 (\lambda_0 + v \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_1 + \frac{v}{2} \lambda_2)} , \\ \quad i \text{ \& t are first associates} \\ \frac{\sigma^2 k \lambda_2}{\lambda_0 (\lambda_0 + v \lambda_2)} , \\ \quad i \text{ \& t are second associates} \end{cases}$$

Furthermore, the correlation between estimates are the following,

$$\text{corr.}\{(\hat{\alpha}_0 - \hat{\alpha}_i), (\hat{\alpha}_0 - \hat{\alpha}_t)\} = \begin{cases} \frac{\lambda_1 (\lambda_0 + \frac{v}{2} \lambda_2) + \frac{v}{2} \lambda_2^2}{(\lambda_0 + \lambda_1 + \frac{v}{2} \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_2) - \frac{v}{2} (\frac{v}{2} - 1) \lambda_2^2} \\ = \rho_1 (\text{say}), i \text{ \& t are first assoc.} \\ \frac{\lambda_2 (\lambda_0 + \frac{v}{2} \lambda_2) + \frac{v}{2} \lambda_2^2}{(\lambda_0 + \lambda_1 + \frac{v}{2} \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_2) - \frac{v}{2} (\frac{v}{2} - 1) \lambda_2^2} \\ = \rho_2 (\text{say}), i \text{ \& t are second assoc.} \end{cases}$$

and trace is given by,

$$\text{Tr}(M^{-1}) = \frac{vk \{ (\lambda_0 + \lambda_1 + \frac{v}{2} \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_2) - \frac{v}{2} (\frac{v}{2} - 1) \lambda_2^2 \}}{\lambda_0 (\lambda_0 + v \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_1 + \frac{v}{2} \lambda_2)}$$

#### 2.1.4 INADMISSIBLE AND EQUIVALENT DESIGNS:

Having defined the terms  $\eta^2$ ,  $\rho_1$  and  $\rho_2$  in the Section 2.1.3 and for fixed  $\lambda_0$ , it can be verified that  $\eta^2$  is a decreasing and  $\rho_1$  and  $\rho_2$  are increasing functions of  $\lambda_1$  and  $\lambda_2$ .

Consider two designs  $D_1$  and  $D_2$  with the following parameters,

$$D_1 \equiv ( b_1, \lambda_0^{(1)}, \lambda_1^{(1)}, \lambda_2^{(1)} )$$

$$D_2 \equiv ( b_2, \lambda_0^{(2)}, \lambda_1^{(2)}, \lambda_2^{(2)} )$$

DEFINITION 1:  $D_1$  and  $D_2$  are equivalent if,

$$b_1 = b_2, \lambda_0^{(1)} = \lambda_0^{(2)}, \lambda_1^{(1)} = \lambda_1^{(2)}, \lambda_2^{(1)} = \lambda_2^{(2)}$$

DEFINITION 2:  $D_2$  is strongly inadmissible with respect to  $D_1$  if.

$$b_1 \leq b_2, \eta_1^2 \leq \eta_2^2, \rho_1^{(1)} \geq \rho_1^{(2)}, \rho_2^{(1)} \geq \rho_2^{(2)}$$

with at least one inequality strict.

Since,  $\eta^2$  is decreasing and  $\rho_1$  and  $\rho_2$  are increasing functions of  $\lambda_1$  and  $\lambda_2$ , the above conditions for inadmissibility of  $D_2$  with respect to  $D_1$  can be written as,

$$b_1 \leq b_2, \lambda_0^{(1)} = \lambda_0^{(2)}, \lambda_1^{(1)} \geq \lambda_1^{(2)}, \lambda_2^{(1)} \geq \lambda_2^{(2)}$$

with at least one inequality strict.

A design which is inadmissible can be discarded because there exists a better design (i.e. with less variance of estimates) with lesser number of blocks.

Furthermore, only one design can be selected from each class of equivalent designs.

#### 2.1.5 SOME LEMMAS:

LEMMA 2.1 For given  $(k, v)$  a group divisible treatment incomplete block design with  $(b, \lambda_0, \lambda_1, \lambda_2)$  satisfy the following inequality,

$$\frac{4v\lambda_0 + v(v-2)\lambda_1 + v^2\lambda_2}{2k(k-1)} \leq b \leq \frac{4v\lambda_0 + v(v-2)\lambda_1 + v^2\lambda_2}{4(k-1)}$$

Furthermore, the lower inequality is an equality iff the design is completely binary (i.e.  $r_{ij} = 0$  or  $1$ ,  $0 \leq i \leq v$ ,  $1 \leq j \leq b$ ).

#### PROOF:

Consider information matrix  $M$  given in Section 1.2.4 and Section 2.1.3. One can write,

$$m_{00} = r_0 - \frac{1}{k} \sum_{j=1}^b r_{0j}^2 = \frac{v\lambda_0}{k}$$

which gives,

$$k r_0 = v\lambda_0 + \sum_{j=1}^b r_{0j}^2 \tag{2.1.5}$$

and also,

$$m_{ii} = r_i - \frac{1}{k} \sum_{j=1}^b r_{ij}^2 = \frac{1}{k} \left\{ \lambda_0 + \frac{v-2}{2} \lambda_1 + \frac{v}{2} \lambda_2 \right\}$$

which gives,

$$k r_i = \lambda_0 + \frac{v-2}{2} \lambda_1 + \frac{v}{2} \lambda_2 + \sum_{j=1}^b r_{ij}^2 \quad 2.1.6$$

Adding 2.1.5 and 2.1.6,

$$k \sum_{i=0}^v r_i = k^2 b = 2v \lambda_0 + \frac{v(v-2)}{2} \lambda_1 + \frac{v^2}{2} \lambda_2 + \sum_{i=0}^v \sum_{j=1}^b r_{ij}^2 \quad 2.1.7$$

it can be easily verified that the minimum value of  $\sum_{i=0}^v \sum_{j=1}^b r_{ij}^2$  is  $bk$  and the maximum value is  $b(k^2 - 2k + 2)$ .

Substituting these two values of  $\sum_{i=0}^v \sum_{j=1}^b r_{ij}^2$  in eq.2.1.7

give the lower and the upper bounds on  $b$  in the Lemma.

REMARK: For  $(k=3, v=4)$  an E MB GD TIB design with parameters  $(b, \lambda_0, \lambda_1, \lambda_2 = \lambda_1 + \delta)$ , where  $\delta = -1, 0, 1$ , satisfy the following,

$$b \geq \frac{2}{3} ( 2 \lambda_0 + \lambda_1 + 2 \lambda_2 )$$

LEMMA 2.2 For given  $(k, v)$  consider a group divisible treatment incomplete block design with  $(b, \lambda_0, \lambda_1, \lambda_2)$ . If the design is binary in terms of test-treatments then the test-treatments are replicated equally in the design



with a common replicate size,

$$r = \frac{2\lambda_0 + (v-2)\lambda_1 + v\lambda_2}{2(k-1)}$$

PROOF: For a design that is binary in terms of test-treatments, eq. 2.1.6 gives,

$$kr_i - r_i = \lambda_0 + \frac{v-2}{2}\lambda_1 + \frac{v}{2}\lambda_2, \quad \left[ \sum_{j=1}^b r_{ij}^2 = \sum_{j=1}^b r_{ij} = r_i \right]$$

So, 
$$r_i = \frac{2\lambda_0 + (v-2)\lambda_1 + v\lambda_2}{2(k-1)}, \quad i = 1, 2, \dots, v$$

It, therefore, follows that,

$$r_1 = r_2 = \dots = r_v = r = \frac{2\lambda_0 + (v-2)\lambda_1 + v\lambda_2}{2(k-1)}$$

REMARK: For  $(k=3, v=4)$  and a E MB GD TIB design with parameters  $(b, \lambda_0, \lambda_1, \lambda_2 = \lambda_1 + \delta)$ , where  $\delta = -1, 0, 1$ , Lemma 2.2 gives,

$$r = \frac{1}{2} (\lambda_0 + \lambda_1 + 2\lambda_2)$$

LEMMA 2.3 For any E MB GD TIB design for  $k = 3, v \geq 3$ , we have,

$$b \geq \frac{v\lambda_0}{2}$$

PROOF: Let  $r_0$  units are allocated to the control and  $\underline{m} = (m_1, m_2)$  represents the allocation vector

(i.e. control appears "1" (  $l = 1, 2$  ) times in " $m_1$ " blocks in the design ). Then,

$$\sum_{l=1}^2 m_1 = m_1 + m_2 \leq b \quad 2.1.8$$

and 
$$\sum_{l=1}^2 l m_1 = m_1 + 2m_2 = r_0 \quad 2.1.9$$

and 
$$\sum_{l=1}^2 l(l-1)m_1 = 2 m_2 \quad 2.1.10$$

Eqs. 2.1.8 and 2.1.9 gives,

$$r_0 - m_2 \leq b \quad 2.1.11$$

Using FACT 2.1 and FACT 2.5(ii),

$$v \lambda_0 = \sum_{i=1}^v \lambda_{oi} = 2(r_0 - m_2)$$

Using 2.1.11 gives,

$$b \geq \frac{v \lambda_0}{2}$$

REMARK: for any E MB GD TIB design for  $k = 3, v = 4,$  we have,

$$b \geq 2 \lambda_0$$

LEMMA 2.4 In a binary GD TIB design with  $k > 2$  odd, the quantities  $v \lambda_0$  and  $\lambda_0 + \frac{v-2}{2} \lambda_1 + \frac{v}{2} \lambda_2$  are even integers.

PROOF: Recall that,

$$r_0 = \sum_{l=1}^{k-1} l m_1 ,$$

$$\sum_{j=1}^b r_{0j}^2 = \sum_{l=1}^{k-1} l^2 m_1 ,$$

$$\sum_{j=1}^b r_{ij}^2 = r_i$$

From eq. 2.1.5, we can obtain the following relation,

$$v \lambda_0 = \sum_{l=1}^{k-1} l(k-1)m_1$$

Now, when  $k$  is odd,  $l(k-1)$  is even. Therefore,  $v \lambda_0$  is also even. Furthermore, from eq. 2.1.6, we can get,

$$\lambda_0 + \frac{v-2}{2} \lambda_1 + \frac{v}{2} \lambda_2 = (k-1)r_i$$

So, when  $k$  is odd,  $k-1$  is even and hence,  $(k-1)r_i$  is even i.e.  $\lambda_0 + \frac{v-2}{2} \lambda_1 + \frac{v}{2} \lambda_2$  is even.

REMARK: For an E MB GD TIB design with  $(b, \lambda_0, \lambda_1, \lambda_1 + \delta)$ , where  $\delta = -1, 0, 1$  and for  $k = 3$  and  $v = 4$ , the quantity,  $\lambda_0 + \lambda_1 + 2 \lambda_2$  is even.

LEMMA 2.5 In a GD TIB design with  $k = 3, v = 4$ , if  $\lambda_0 = 0$ , then  $\lambda_1 = \lambda_2$ .

PROOF: In such a design  $(1,2), (3,4)$  are first associates, whereas,  $(1,3), (1,4), (2,3), (2,4)$  are second associates.

In any block (e.g. 1,2,3 or 1,3,4 etc.) of such a design, we have,

number of TT-pairs of first associates type = 1

and, number of TT-pairs of second associates type = 2

Therefore, in  $b$  blocks,

total number of TT-pairs of first associates type =  $b$

total number of TT-pairs of second associates type =  $2b$

Using the relations in FACT 2.3, one can write,

$$b = \frac{1}{4} \cdot 4 \cdot 2 \cdot \lambda_1 \quad \Rightarrow \quad b = 2 \lambda_1$$

$$\text{and,} \quad 2b = \frac{1}{4} \cdot 4^2 \cdot \lambda_2 \quad \Rightarrow \quad b = 2 \lambda_2$$

Hence,  $\lambda_1 = \lambda_2$ .

REMARK: In an E MB GD TIB design for  $k = 3, v = 3$  and with  $(b, \lambda_0 = 0, \lambda_1, \lambda_2 = \lambda_1 + \delta)$ , where  $\delta = -1, 0, 1$ , the only designs which can exist are those for which  $\delta = 0$  (i.e.  $\lambda_1 = \lambda_2$ ).

LEMMA 2.6 In a GD TIB design with  $k = 3, v = 4$ , if  $\lambda_0 = 1$ , then  $\lambda_2 \neq \lambda_1$ .

PROOF: In such a design (1,2), (3,4) are first associates and (1,3), (1,4), (2,3), (2,4) are second associates. If  $\lambda_0 = 1$ , the following two situations may arise.

(i) two blocks contain the control once in each block, and each of the remaining  $b - 2$  blocks, contain test-treatments only.

BLOCK		1	2	3 .....	b
PLOT	1	0	0	test-treatments only	
	2	1	2		
	3	3	4		
		$B_1$		$B_2$	

number of TT-pairs of first associates type in  $B_1 = 0$

number of TT-pairs of first associates type in  $B_2 = b-2$   
(see Lemma 2.5)

total number of TT-pairs of first associates type =  $b-2$

number of TT-pairs of second associates type in  $B_1=2$

number of TT-pairs of second associates type in  $B_2=2(b-2)$   
(see Lemma 2.5)

total number of TT-pairs of second associates type =  $2b-2$

Using the relations in FACT 2.3, we can write,

$$b - 2 = 2 \lambda_1$$

and,  $2b - 2 = 4 \lambda_2$

Hence,

$$\lambda_2 = \lambda_1 + \frac{1}{2}, \quad \text{i.e. } \lambda_2 \neq \lambda_1$$

(ii) Same situation as in (i) above but, the following configuration,

BLOCK		1	2	3 .....	b
PLOT	1	0	0	test-treatments only	
	2	1	3		
	3	2	4		
		$B_1$		$B_2$	

The following are, therefore, true:

total number of TT-pairs of first associates type = b

total number of TT-pairs of second associates type=2(b-2)

Using FACT 2.3, we can write,

$$b = 2 \lambda_1,$$

and,  $2(b-2) = 4 \lambda_1$

Hence,  $\lambda_2 = \lambda_1 - 1$  i.e.  $\lambda_2 \neq \lambda_1$

The proof of the Lemma is, therefore, complete.

**LEMMA 2.7** In an GD TIB design with  $k = 4, v = 4$  and  $(b, \lambda_0, \lambda_1, \lambda_2 = \lambda_1 + \delta)$ ,  $\delta = -1, 0, 1$ , the following are true,

(i) if  $\lambda_0 = 0 \pmod{3}$ , then  $\lambda_2 = \lambda_1$

(ii) if  $\lambda_0 = 1 \pmod{3}$ , then  $\lambda_2 = \lambda_1 + 1$

(iii) if  $\lambda_0 = 2 \pmod{3}$ , then  $\lambda_2 = \lambda_1 - 1$

**PROOF:** Let the design contain  $j$  th BB combination. Then by FACT 2.6 there are  $U = \lambda_0 - 3j + 3$  4-BBs and  $V = 4j - 4$  3-BBs ( $j = 1, 2, \dots, \lceil \lambda_0/3 \rceil + 1$ ) if  $\lambda_0$  is even, or  $U = \lambda_0 - 3j$  4-BBs and  $V = 4j$  3-BBs ( $j = 1, 2, \dots, \lceil \lambda_0/3 \rceil$ ) if  $\lambda_0$  is odd.

Let there are "q" (an integer) 4-BBs in which x & y are first associates (see Definition 4 of Section 1.4.2). Hence, in the rest  $U - q$  4-BBs, x & y are second associates.

The following facts are, therefore, noted.

number of TT-pairs of first associates type in "q" 4-BBs = q

number of TT-pairs of second associates type in "q"  $4\text{-BBs} = 0$   
 number of TT-pairs of first associate type in "U-q"  $4\text{-BBs} = 0$   
 number of TT-pairs of second associates type in "U - q"  
 $4\text{-BBs} = U - q$

number of TT-pairs of first associates type in "V"  $3\text{-BBs} = V$   
 number of TT-pairs of second associates type in "V"  $3\text{-BBs}$   
 $= 2 v$

number of TT-pairs of first associates type in  $b - b_j$   
 remaining blocks  $= 2 (b - b_j)$

number of TT-pairs of second associates type in  $b - b_j$   
 remaining blocks  $= 4 (b - b_j)$

where,  $b_j = U + V$ . So, the

total number of TT-pairs of first associates type in the  
 design  $= q + V + 2 (b - b_j)$

total number of TT-pairs of second associates type in the  
 design  $= U - q + 2 V + 4 (b - b_j)$

Using FACT 2.3, one can get,

$$q = \begin{cases} \frac{\lambda_0}{3} - j + 1 - \frac{4}{3} \delta, & \text{if } \lambda_0 \text{ is even} \\ \frac{\lambda_0}{3} - j - \frac{4}{3} \delta, & \text{if } \lambda_0 \text{ is odd} \end{cases}$$

Consider the following cases,

CASE 1:  $\lambda_0 = 0 \pmod{3} = 3m, m = 0, 1, 2, \dots$

$$q = \begin{cases} m - j + 1 - \frac{4}{3} \delta, & \text{if } \lambda_0 \text{ is even} \\ m - j - \frac{4}{3} \delta, & \text{if } \lambda_0 \text{ is odd} \end{cases}$$

Now,  $q$  is integer only when  $\delta = 0$ , i.e. when  $\lambda_2 = \lambda_1$ . Part (i) of the Lemma is, therefore, proved.

CASE 2:  $\lambda_0 = 1 \pmod{3} = 3m + 1, m = 0, 1, 2, \dots$

$$q = \begin{cases} m - j - \frac{4}{3} (\delta - 1), & \text{if } \lambda_0 \text{ is even} \\ m - j - \frac{1}{3} (4\delta - 1), & \text{if } \lambda_0 \text{ is odd} \end{cases}$$

The  $q$  is integer only for  $\delta = 1$ , i.e., for  $\lambda_2 = \lambda_1 + 1$ . This proves part (ii) of the Lemma.

CASE 3:  $\lambda_0 = 2 \pmod{3} = 3m + 2, m = 0, 1, 2, \dots$

$$q = \begin{cases} m - j - \frac{1}{3} (4\delta - 5), & \text{if } \lambda_0 \text{ is even} \\ m - j - \frac{2}{3} (2\delta - 1), & \text{if } \lambda_0 \text{ is odd} \end{cases}$$

Here,  $q$  is integer if  $\delta = -1$ , therefore,  $\lambda_2 = \lambda_1 - 1$ . This proves part (iii) of the Lemma.

This completes the proof of Lemma 2.7

## 2.2 METHOD OF CONSTRUCTION:

consider a design with the parameters  $k, v, b$  and allocation of  $r_0$  units to the control.

From FACT 2.1,

$$\lambda_0 = \sum_{i=1}^v \lambda_{0i} / v \quad 2.2.1$$

where,  $\sum_{i=1}^v \lambda_{0i}$  is given by eq. 2.1.3



Also from FACT 2.3,

$$\lambda_1 + \frac{v}{v-2} \lambda_2 = \frac{4}{v(v-2)} \sum_{i \neq t} \lambda_{it} \quad 2.2.2$$

where,  $\sum_{i \neq t} \lambda_{it}$  is given by eq. 2.1.4.

For the construction of the design the following algorithm is used.

1. check that the value of  $r$  calculated from eq. 2.1.1 is an integer.
2. select vector  $\mathbf{m} = (m_1, m_2, \dots, m_{k-1})$  which satisfy eq.2.1.2, and compute  $\sum \lambda_{0i}$  and  $\sum \lambda_{it}$  from eqs. 2.1.3 and 2.1.4 respectively.
3. check whether  $\lambda_0$  from eq. 2.2.1 is an integer.
4. try to find some integer values of  $\lambda_1$  and  $\lambda_2$  which satisfy eq. 2.2.2.
5. verify that the design is not non-existent by applying Theorem in Section 2.1.2.

If all the conditions mentioned above are satisfied a design with these parameters may be constructed. It is, however, warned that the satisfaction of these conditions does not imply the actual existence of the design. Trial to actually construct the design are always required. Usually, with few trials it becomes clear that whether or not the treatments can be allocated to the plots according to the plan suggested by the parameters.

**EXAMPLE:** Consider a design with the following parameters and follow the algorithm.

$$k = 4, v = 8, b = 23, r_0 = 36$$

1.  $r = (bk - r_0)/v = 7$  ( an integer )
2. let  $m = (m_1, m_2, m_3) = (11, 11, 1)$ . So,

$$\sum_{l=1}^3 lm_l = 36 \quad (\text{eq. 2.1.2 satisfied})$$

and, 
$$\sum_{l=1}^3 l(1-l)m_l = 28$$

Hence, from eq. 2.1.3 and eq. 2.1.4, we have,

$$\sum_{i=1}^8 \lambda_{0i} = 80, \quad \text{and} \quad \sum_{i,t=1}^8 \lambda_{it} = 44$$

3. from eq. 2.2.1  $\lambda_0 = 80/8 = 10$  ( an integer )
4. the values  $\lambda_1 = 1$  and  $\lambda_2 = 2$ , satisfy eq 2.2.2
5. note that,  $\lambda_1 < \lambda_2$  and  $r > \lambda_1 + \frac{v(\lambda_2 - \lambda_1)}{2}$ .

According to the Theorem of non-existence in Section 2.1.2, the design is non-existent if  $b < v$ . Since,  $b > v$ , the design may, therefore, exist .

In fact, the design exists, with the following structure.

BLOCK	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	2	2	2	
PLOT 1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
2	1	1	1	1	1	2	2	2	3	3	4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
3	2	3	4	5	6	3	4	7	4	5	6	1	2	2	2	3	3	3	4	4	4	5	0			
4	5	6	7	8	7	8	6	8	5	7	8	8	5	6	7	7	6	8	5	7	8	6	1			

CHAPTER 3

**3.1 MINIMAL COMPLETE CLASS OF GENERATOR  
DESIGNS ( MCCGD ) FOR  $k = 2, v = 4$  :**

**THEOREM 3.1**

For  $k = 2$  and  $v = 4$  the minimal complete class of generator designs is given in Table 3.1.

TABLE 3.1 MCCGD FOR  $k = 2, v = 4$  .

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
$D_1$	2	0	1	0	0	1	0	1 3 2 4
$D_2$	4	0	2	0	0	0	1	1 1 2 2 3 4 3 4
$D_3$	4	4	1	4	1	0	0	0 0 0 0 1 2 3 4

**PROOF:**

The above mentioned designs  $D_1, D_2, D_3$  are the only possible generator designs and since they are distinct and neither of them is strongly inadmissible with respect to the other one, they constitute the minimal complete set (see Section 1.3.3) for  $k = 2, v = 4$ .

Each of the following theorems has the similar proof as given in the case of Theorem 3.1.

3.2 MCCGD FOR  $k = 2, v = 6$  :

THEOREM 3.2

For  $k = 2, v = 6$  the minimal complete class of generator designs is given in Table 3.2.

TABLE 3.2 MCCGD FOR  $k = 2, v = 6$ .

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE											
$D_1$	6	0	2	0	0	1	0	1	1	2	4	4	5	2	3	3	5	6	6
$D_2$	6	6	1	6	1	0	0	0	0	0	0	0	0	1	2	3	4	5	6
$D_3$	9	0	3	0	0	0	1	1	1	2	2	2	3	3	3	3	3	3	3

3.3 MCCGD FOR  $k = 2, v = 8$  :

THEOREM 3.3

For  $k = 2$  and  $v = 8$  the minimal complete class of generator designs is given in Table 3.3.

**TABLE 3.3 MCCGD FOR  $k = 2, v = 8$ .**

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
$D_1$	8	8	1	8	1	0	0	0 0 0 0 0 0 0 0 1 2 3 4 5 6 7 8
$D_2$	12	0	3	0	0	1	0	1 1 1 2 2 3 5 5 6 6 7 2 3 4 3 4 4 6 8 7 8 8
$D_3$	16	0	4	0	0	0	1	1 1 1 1 2 2 2 2 3 3 3 3 4 5 6 7 8 5 6 7 8 5 6 7 8 5 ----- 4 4 4 6 7 8

**3.4 MCCGD FOR  $k = 2, v = 10$  :**

**THEOREM 3.4**

For  $k = 2$  and  $v = 10$  the minimal complete class of generator designs is given in Table 3.4.

**TABLE 3.4 MCCGD FOR  $k = 2, v = 10$**

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
$D_1$	10	10	1	10	1	0	0	0 0 0 0 0 0 0 0 0 0 1 2 3 4 5 6 7 8 9 10
$D_2$	20	0	4	0	0	1	0	1 1 1 1 2 2 2 3 3 4 6 6 6 2 3 4 5 3 4 5 4 5 5 7 8 9 ----- 6 7 7 7 8 8 9 10 8 9 10 9 10 10

**TABLE 3.4 (CONTINUED)**

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE									
D <sub>3</sub>	25	0	5	0	0	0	1	1	1	1	1	1	2	2	2	2	
								6	7	8	9	10	6	7	8	9	
								2	3	3	3	3	3	4	4	4	
								10	6	7	8	9	10	6	7	8	
4	4	5	5	5	5	5											
9	10	6	7	8	9	10											



CHAPTER 4

**4.1 GENERAL METHOD OF PROVING A CLASS  
OF GENERATOR DESIGNS TO BE MCCGD :**

Let us suppose that for a given  $(v, k)$ ,  $\{ D_1, D_2, \dots, D_n \}$  is a set of generator designs. Also, consider an arbitrary design "D" with Parameters  $( b, \lambda_0, \lambda_1, \lambda_2 )$ , where  $\lambda_2 = \lambda_1 + \delta, \delta = -1, 0, 1$ . To show that the set of generator designs is MCCGD, we show that there exists a design  $D^*$  with parameters  $( b^*, \lambda_0^*, \lambda_1^*, \lambda_2^* )$ , such that,

$$D^* = \bigcup_{i=1}^n f_i D_i.$$

$$\lambda_0^* = \lambda_0, \lambda_1^* = \lambda_1, \lambda_2^* = \lambda_2 \quad \text{and} \quad b^* \leq b$$

where,

$$\lambda_0^* = \sum_{i=1}^n f_i \lambda_0^{(i)}, \lambda_1^* = \sum_{i=1}^n f_i \lambda_1^{(i)}, \lambda_2^* = \sum_{i=1}^n f_i \lambda_2^{(i)}, b^* = \sum_{i=1}^n f_i b_i$$

$b_i, \lambda_0^{(i)}, \lambda_1^{(i)}, \lambda_2^{(i)}$  are the parameters of generator design  $D_i$  ( $i = 1, 2, \dots, n$ ).

The design D is, therefore, either equivalent to or strongly inadmissible with respect to  $D^*$  (Section 1.3.3 and Section 2.1.4). The proof is completed by finally noting that the set  $\{ D_1, D_2, \dots, D_n \}$  consists of non-equivalent generator designs.

In the next section, we give MCCGD for E MB GD TIB design for  $k = 3$ ,  $v = 4$  and prove the minimal complete nature of the class. In other sections MCCGD are given for  $k = 3$  and  $v = 6, 8, 10$ , without the accompanying proofs which are very long and tedious, but similar to the one given in the next section.

4.2 MCCGD FOR  $k = 3, v = 4$  :

THEOREM 4.2

For  $k = 3$  and  $v = 4$  the minimal complete class of generator designs is given in Table 4.1.

TABLE 4.2 MCCGD FOR  $k = 3, v = 4$ .

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE					
$D_1$	2	2	1	2	1	1	0	0	0				
				0				1	3				
								2	4				
$D_2$	4	0	3	0	0	2	2	1	1	1	2		
				0				2	2	3	3		
								3	4	4	4		
$D_3$	4	8	1	0	2	0	0	0	0	0	0		
				4				0	0	0	0		
								1	2	3	4		



TABLE 4.2 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE									
D <sub>4</sub>	4	4	2	4	2	0	1	0	0	0	0	0	0	0	0	0	0
				0				1	1	2	2						
								3	4	3	4						
D <sub>5</sub>	6	6	3	2	2	2	1	0	0	0	0	1	2				
				2				1	1	0	0	3	3				
								2	2	3	4	4	4				
D <sub>6</sub>	7	5	4	3	2	2	2	0	0	0	0	1	1	2			
				1				1	1	2	0	2	3	3			
								2	3	3	4	4	4	4			
D <sub>7</sub>	8	8	4	4	3	1	2	0	0	0	0	0	0	1	1		
				2				1	2	2	2	0	0	2	3		
								4	3	4	4	1	3	3	4		
D <sub>8</sub>	11	9	6	3	3	3	3	0	0	0	0	0	0	1	1	1	
				3				1	2	3	0	0	0	2	2	2	
								4	4	4	1	2	3	3	3	4	
								1	2								
								3	3								
								4	4								

**PROOF:**

Before going in the detail of the proof, note that  $\lambda_0 - 3\lambda_1$  and  $3\lambda_1 - \lambda_0$  are even integers ( Remark of LEMMA 2.4 ). This implies that  $\lambda_1$  is even if  $\lambda_0$  is even, and  $\lambda_1$  is odd if  $\lambda_0$  is odd.

CASE 1.  $\lambda_0 = 0 \pmod{3}$ .

(a)  $\lambda_2 = \lambda_1 - 1$

Note that  $\lambda_1 \geq 1$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$  with,

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1 + 2}{2}, \quad f_4 = \lambda_1 - 1$$

Note the following,

- the quantity  $\lambda_0 - 3\lambda_1 + 2$  is even (Remark in Lemma 2.4), therefore,  $f_3$  is an integer.
- $f_3 \geq 0$  if  $\lambda_1 \leq \lambda_0/3$
- $b^* = f_1 \cdot 2 + f_3 \cdot 4 + f_4 \cdot 4 = 2\lambda_0$

Therefore,  $b \geq b^*$  (Remark in Lemma 2.3)

Hence,  $D$  is either equivalent (if  $b = b^*$ ) or  $S$ -inadmissible (if  $b > b^*$ ) with respect to  $D^*$  for,

$$\lambda_0 = 0 \pmod{3}, \quad \lambda_2 = \lambda_1 - 1, \quad \lambda_1 \leq \lambda_0/3.$$

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4 \cup f_6 D_6$ , with

$$f_1 = \frac{\lambda_0}{3}, \quad f_2 = \frac{3\lambda_1 - \lambda_0 - 6}{6}, \quad f_4 = \frac{\lambda_0 - 3}{3}, \quad f_6 = 1.$$

Note that,

- quantities  $\lambda_0$  and  $\lambda_0 - 3$  are both divisible by 3. Hence,  $f_1$  and  $f_4$  are integers.
- $f_2 \geq 0$ , if  $\lambda_1 \geq \frac{\lambda_0}{3} + 2$ .
- $\lambda_0 = 3 \pmod{6}$  implies that  $\lambda_1 = 1 \pmod{2}$  for  $f_2$  to be integer, and  $\lambda_0 = 0 \pmod{6} > 0$  implies that

$\lambda_1 = 0 \pmod{2}$  for  $f_2$  to be integer.

This includes all possible values of  $\lambda_0$  and  $\lambda_1$  for the design except  $\lambda_0 = 0, \lambda_1 = 0 \pmod{2}$ .

$$\square \quad b^* = f_1 \cdot 2 + f_2 \cdot 4 + f_4 \cdot 4 + 1.7 = \frac{2}{3} (2\lambda_0 + 3\lambda_1) - 1$$

$\square$  Now, by Remark of Lemma 2.1

$$b \geq \frac{2}{3} (2\lambda_0 + 3\lambda_1) - \frac{4}{3}$$

It is easy to see that  $\frac{2}{3}(2\lambda_0 + 3\lambda_1)$  is an integer. Therefore, right hand side of the above inequality is not integer. The next higher integer is the same as  $b^*$ . Hence,

$$b \geq b^*$$

The design  $D$  is , therefore, equivalent or  $S$ -inadmissible with respect to  $D^*$ , except in the following case,

$$\square \quad \lambda_0 = 0 \text{ and } \lambda_1 = 0 \pmod{2} \geq 2$$

In such a case Lemma 2.5 is applicable and hence, the only designs exist are those for which  $\lambda_2 = \lambda_1$ . So, No design with the above mentioned parameter values and  $\lambda_2 = \lambda_1 - 1$  exists. So, the possibility of  $\lambda_0 = 0$  is ruled out.

Hence,  $D$  is equivalent or  $S$ -admissible with respect to  $D^*$  for,

$$\lambda_0 = 0 \pmod{3}, \lambda_2 = \lambda_1 - 1, \lambda_1 \geq \frac{\lambda_0}{3} + 2.$$

Only one case when  $\lambda_1 = 1 + \lambda_0/3$  is also not covered here. It can be shown by using Remark of Lemma 2.2

that  $r$  is not integer ( $r = \lambda_0 + 1/2$ ) in such a case and hence, there do not exist designs with  $\lambda_1 = 1 + \frac{\lambda_0}{3}$  and  $\lambda_2 = \lambda_1 - 1$ .

(b)  $\underline{\lambda_2 = \lambda_1}$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1}{2}, \quad f_4 = \lambda_1$$

Note the following,

- $\lambda_0 - 3\lambda_1$  is even by Lemma 2.4. Therefore,  $f_3$  is integer.
- $\lambda_1 \leq \lambda_0/3$  in order to have  $f_3 \geq 0$ .
- $b^* = f_1 \cdot 2 + f_3 \cdot 4 + f_4 \cdot 4 = 2\lambda_0$

Therefore,  $b \geq b^*$  by Lemma 2.3.

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4$ , with

$$f_1 = \lambda_0/3, \quad f_2 = \frac{3\lambda_1 - \lambda_0}{6}, \quad f_4 = \lambda_0/3$$

Note the following,

- $\lambda_1 \geq \lambda_0/3$  in order to have  $f_2 \geq 0$ .
- $\lambda_0 = 0 \pmod{6}$  implies that  $\lambda_1 = 0 \pmod{2}$  for  $f_2$  to be integer, and  $\lambda_0 = 3 \pmod{6}$  implies that  $\lambda_1 = 1 \pmod{2}$  for  $f_2$  to be integer.
- By Lemma 2.1,  $b^*$  requires the minimum number of blocks because  $D_1, D_2,$  and  $D_4$  all are completely binary. Hence,  $b \geq b^*$ .

By (i) and (ii) above it is clear that  $D$  is either equivalent to or  $S$ -inadmissible with respect to a design which is union of copies of designs from Table 4.2 for,

$$\lambda_0 = 0 \pmod{3} \text{ and } \lambda_2 = \lambda_1 .$$

(c)  $\underline{\lambda_2 = \lambda_1 + 1}$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1 - 2}{2}, \quad f_4 = \lambda_1 + 1$$

Note the following,

- $\lambda_0 - 3\lambda_1 - 2$  is even by Lemma 2.4. It implies that  $f_3$  is integer.
- $f_3 \geq 0$  iff  $\lambda_1 \leq \frac{\lambda_0 - 2}{3}$ . But  $\frac{\lambda_0 - 2}{3}$  is not integer, therefore,

$$\lambda_1 \leq \frac{\lambda_0 - 3}{3} \quad (\text{next lower integer})$$

□  $b^* = 2\lambda_0$

Therefore,  $b \geq b^*$  (Lemma 2.3)

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2$ , with,

$$f_1 = \lambda_0, \quad f_2 = \frac{\lambda_1 - \lambda_0}{2}$$

Note the following,

- $\lambda_1 - \lambda_0$  is even by Lemma 2.4  $\Rightarrow f_3$  is integer.
- $f_2 \geq 0 \Rightarrow \lambda_1 \geq \lambda_0$
- $D^*$  contains minimum number of blocks because  $D_1$  and  $D_2$  are completely binary ( Lemma 2.1).

Hence,  $b \geq b^*$ .

(iii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4 \cup f_7 D_7$ , with

$$f_1 = \frac{\lambda_0 - 3}{3}, f_2 = \frac{3\lambda_1 - \lambda_0}{6}, f_4 = \frac{\lambda_0 - 3}{3}, f_7 = 1$$

Note the following,

- $\lambda_0 - 3$  is divisible by 3  $\Rightarrow f_1$  and  $f_4$  are integers
- $f_2 \geq 0$  iff  $\lambda_1 \geq \lambda_0/3$
- $\lambda_0 = 3 \pmod{6} \Rightarrow \lambda_1 = 1 \pmod{2}$  for  $f_2$  to be integer
- $\lambda_0 = 0 \pmod{6} > 0 \Rightarrow \lambda_1 = 0 \pmod{2}$  for  $f_2$  to be integer.
- the cases  $\lambda_0 = 0, \lambda_1 \geq 0$  are covered in (ii) above.
- $b^* = \frac{2}{3} ( 2\lambda_0 + 3\lambda_1 + 2 )$

Whereas,  $b \geq \frac{2}{3} ( 2\lambda_0 + 3\lambda_1 + 2 )$  by Lemma 2.1.

Therefore,  $b \geq b^*$

Hence,  $D$  is either equivalent or  $S$ -inadmissible with respect to a design which is a union of copies of designs from Table 4.1 for,

$$\lambda_0 = 0 \pmod{3}, \lambda_2 = \lambda_1 + 1$$

So, (a), (b) and (c) covers all possibilities for,

$$\lambda_0 = 0 \pmod{3}.$$

CASE 2.  $\lambda_0 = 1 \pmod{3}$

(a)  $\lambda_2 = \lambda_1 - 1$

Note that  $\lambda_1 \geq 1$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1 + 2}{2}, \quad f_4 = \lambda_1 - 1$$

Note the following,

□  $\lambda_0 - 3\lambda_1 + 2$  is even (Lemma 2.4)  $\Rightarrow f_3$  is integer

□  $f_3 \geq 0$  iff  $\lambda_1 \leq \lambda_0/3$

Since  $\lambda_0/3$  is not integer,  $\lambda_1 \leq \frac{\lambda_0 - 1}{3}$

□  $b^* = 2\lambda_0$

Hence,  $b \geq b^*$  (Lemma 2.3)

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4$ , with

$$f_1 = \frac{\lambda_0 + 2}{3}, \quad f_2 = \frac{3\lambda_1 - \lambda_0 - 2}{6}, \quad f_4 = \frac{\lambda_0 - 1}{3}$$

Note the following,

□ Since,  $\lambda_0 + 2$  and  $\lambda_0 - 1$  are both divisible by 3,  $f_1$  and  $f_4$  are integers.

□  $f_2 \geq 0$  iff  $\lambda_1 \leq \frac{\lambda_0 + 2}{3}$

- $\lambda_0 = 1 \pmod{6}$  and  $\lambda_1 = 1 \pmod{2} \Rightarrow f_2$  is integer.
- $\lambda_0 = 4 \pmod{6}$  and  $\lambda_1 = 0 \pmod{2} \geq 2 \Rightarrow f_2$  is integer.
- $b^*$  requires minimum number of blocks because  $D_1, D_2$  and  $D_4$  all are completely binary (Lemma 2.1).

Hence,  $b \geq b^*$ .

Cases (i) and (ii) covers the whole range for,

$$\lambda_0 = 1 \pmod{3} \text{ and } \lambda_2 = \lambda_1 - 1.$$

$$(b) \quad \underline{\lambda_2 = \lambda_1}$$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1}{2}, \quad f_4 = \lambda_1$$

Note the following,

- $\lambda_0 - 3\lambda_1$  is even (Lemma 2.4)  $\Rightarrow f_3$  is integer.
- $f_3 \geq 0$  iff  $\lambda_1 \leq \frac{\lambda_0 - 1}{3}$  (in fact  $\lambda_1 \leq \lambda_0/3$ , but  $\lambda_0/3$  is not integer)
- $b^* = 2\lambda_0$

Hence,  $b \geq b^*$  (Lemma 2.3)

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4 \cup f_7 D_7$ , with

$$f_1 = \frac{\lambda_0 - 1}{3}, \quad f_2 = \frac{3\lambda_1 - \lambda_0 - 2}{6}, \quad f_4 = \frac{\lambda_0 - 4}{3}, \quad f_7 = 1$$



Note the following,

- Since,  $\lambda_0 - 1$  and  $\lambda_0 - 4$  are divisible by 3,  $f_1$  and  $f_4$  are integers.
- $f_2 \geq 0$  iff  $\lambda_1 \geq \frac{\lambda_0 + 2}{3}$
- $\lambda_0 = 1 \pmod{6} > 1$ ,  $\lambda_1 = 1 \pmod{2} > 1$  for  $f_2$  to be integer.
- $\lambda_0 = 4 \pmod{6}$ ,  $\lambda_1 = 0 \pmod{2}$  for  $f_2$  to be integer.
- The case when  $\lambda_0 = 1$  is not covered above. But it is clear from Lemma 2.6 that for  $\lambda_0 = 1$  a design with  $\lambda_2 = \lambda_1$  does not exist.
- $b^* = \frac{2}{3} ( 2 \lambda_0 + 3 \lambda_1 + 1 )$ .  
This is the next higher integer than the lower bound  $\frac{2}{3} ( 2 \lambda_0 + 3 \lambda_1 )$  on  $b$  given by Lemma 2.1.  
Hence,  $b \geq b^*$ .

Cases (i), (ii) covers the whole range for,

$$\lambda_0 = 1 \pmod{3} \text{ and } \lambda_2 = \lambda_1$$

(c)  $\lambda_2 = \lambda_1 + 1$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3 \lambda_1 - 2}{2}, \quad f_4 = \lambda_1 + 1$$

Note the following,

- $f_3$  is integer because  $\lambda_0 - 3 \lambda_1 - 2$  is even (Lemma 2.4)

$$\square \quad f_3 \geq 0 \text{ iff } \lambda_1 \leq \frac{\lambda_0 - 4}{3} \quad (\text{an integer})$$

$$\square \quad b^* = 2\lambda_0$$

Hence,  $b \geq b^*$  (Lemma 2.3)

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4 \cup f_6 D_6$ , with

$$f_1 = \frac{\lambda_0 - 4}{3}, \quad f_2 = \frac{3\lambda_1 - \lambda_0 - 2}{6}, \quad f_4 = \frac{\lambda_0 - 1}{3}, \quad f_6 = 1$$

Note that,

$\square \quad \lambda_0 - 4$  and  $\lambda_0 - 1$  are both divisible by 3, hence,  $f_1$  and  $f_4$  are integers.

$$\square \quad f_2 \geq 0 \text{ iff } \lambda_1 \geq \frac{\lambda_0 + 2}{3}$$

$\square \quad \lambda_0 = 1 \pmod{6}, \lambda_1 = 1 \pmod{2}$  for  $f_2$  to be integer

$\lambda_0 = 4 \pmod{6}, \lambda_1 = 0 \pmod{2} \geq 2$  for  $f_2$  to be integer.

$$\square \quad b^* = \frac{2}{3} (2\lambda_0 + 3\lambda_1 - 8)$$

$b^*$  is smaller than the lower bound  $\frac{2}{3}(2\lambda_0 + 3\lambda_1 + 2)$  on  $b$  given by Lemma 2.1.

Hence,  $b \geq b^*$

The case which is not covered above is when  $\lambda_1 = \frac{\lambda_0 - 1}{3}$ .

In such a case  $r = \lambda_0 + 1/2$  (from Lemma 2.2) which is not integer, and therefore, such a design does not exist.

All the possibilities are, therefore, covered by (i), (ii)

above for,

$$\lambda_0 = 1 \pmod{3} \text{ and } \lambda_2 = \lambda_1 + 1$$

CASE 3.  $\lambda_0 = 2 \pmod{3}$

(a)  $\lambda_2 = \lambda_1 - 1$

Note that  $\lambda_1 \geq 1$

(i) let  $D^* = f_2 D_2 \cup f_5 D_5$ , with

$$f_2 = \frac{\lambda_1 - 2}{2}, \quad f_5 = 1$$

Note the following,

- $f_2 \geq 0$  and integer if  $\lambda_1 = 0 \pmod{2} \geq 2$
- $\lambda_0 = 2$
- $b^* = 2(\lambda_1 + 1)$
- $b \geq 2\lambda_1 + 4/3$  (Lemma 2.1)

This lower bound on  $b$  is not an integer. therefore, the next higher possible integer is  $2(\lambda_1 + 2)$  which is larger than  $b^*$ .

Hence,  $b > b^*$

(ii) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1 + 2}{2}, \quad f_4 = \lambda_1 - 1$$

Note the following,

- $\lambda_0 - 3\lambda_1 + 2$  is even by Lemma 2.4, i.e.  $f_3$  is integer.

$$\square \quad f_3 \geq 0 \text{ iff } \lambda_1 \leq \frac{\lambda_0 - 2}{3}$$

$$\square \quad b^* = 2\lambda_0$$

Therefore,  $b \geq b^*$  (Lemma 2.3)

(iii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4 \cup f_7 D_7$ , with

$$f_1 = \frac{\lambda_0 + 1}{3}, f_2 = \frac{3\lambda_1 - \lambda_0 - 4}{6}, f_4 = \frac{\lambda_0 - 5}{3}, f_7 = 1$$

Note the following,

$\square \quad \lambda_0 + 1$  and  $\lambda_0 - 5$  are divisible by 3, i.e.  $f_1$  and  $f_4$  are integers.

$$\square \quad f_2 \geq 0 \text{ iff } \lambda_1 \geq \frac{\lambda_0 + 4}{3}$$

$\square \quad \lambda_0 = 2 \pmod{6} > 2, \lambda_1 = 0 \pmod{2} > 2$  for  $f_2$  to be integer.

$\lambda_0 = 5 \pmod{6}, \lambda_1 = 1 \pmod{2} > 1$  for  $f_2$  to be integer.

$\square \quad$  A case not covered above is for  $\lambda_0 = 2$  and  $\lambda_1 = 0 \pmod{2}$ . This is already covered in (i) above.

$$\square \quad b^* = \frac{2}{3}(2\lambda_0 + 3\lambda_1 - 1)$$

$\square \quad b \geq \frac{2}{3}(2\lambda_0 + 3\lambda_1 - 2)$  by Lemma 2.1. This lower bound is not integer and  $b^*$  is smaller than the next higher integer  $\frac{2}{3}(2\lambda_0 + 3\lambda_1 + 2)$ .

Hence,  $b > b^*$

The case for  $\lambda_1 = \frac{\lambda_0 + 1}{3}$  can be discarded as

$r = \lambda_0 - 1/2$  (Lemma 2.2) is not an integer and such designs do not exist.

This covers all the possibilities for,

$$\lambda_0 = 2 \pmod{3} \text{ and } \lambda_2 = \lambda_1 - 1$$

(b)  $\underline{\lambda_2 = \lambda_1}$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, \quad f_3 = \frac{\lambda_0 - 3\lambda_1}{2}, \quad f_4 = \lambda_1$$

Note that,

- $\lambda_0 - 3\lambda_1$  is even, therefore,  $f_3$  is integer
- $f_3 \geq 0$  iff  $\lambda_1 \leq \frac{\lambda_0 - 2}{3}$  (an integer)
- $b^* = 2\lambda_0$

Hence,  $b \geq b^*$  (Lemma 2.3)

(iii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4 \cup f_6 D_6$ , with

$$f_1 = \frac{\lambda_0 - 2}{3}, \quad f_2 = \frac{3\lambda_1 - \lambda_0 - 4}{6}, \quad f_4 = \frac{\lambda_0 - 2}{3}, \quad f_6 = 1$$

Note that,

- $\lambda_0 - 2$  is divisible by 3, hence,  $f_1$  and  $f_4$  are integers.
- $f_2 \geq 0$  iff  $\lambda_1 \geq \frac{\lambda_0 + 4}{3}$

□  $\lambda_0 = 2 \pmod{6}, \lambda_1 = 0 \pmod{2} > 0$  for  $f_2$  to be integer.

$\lambda_0 = 5 \pmod{6}, \lambda_1 = 1 \pmod{3} > 1$  for  $f_2$  to be integer.

□  $b^* = \frac{2}{3} ( 2 \lambda_0 + 3 \lambda_1 + \frac{1}{2} )$

□  $b \geq \frac{2}{3} ( 2 \lambda_0 + 3 \lambda_1 )$  by Lemma 2.1, but the lower bound is not an integer. The next higher integer than this lower bound is  $\frac{2}{3} ( 2 \lambda_0 + 3 \lambda_1 + 2 )$  which is larger than  $b^*$ .

Hence,  $b \geq b^*$

The uncovered case for  $\lambda_1 = \frac{\lambda_0 + 1}{3}$  gives  $r = \lambda_0 + 1/2$

(Lemma 2.2) which is not an integer and, therefore, the designs do not exist.

All the possibilities has, therefore, been considered for,

$$\lambda_0 = 2 \pmod{3}, \lambda_2 = \lambda_1$$

(c)  $\underline{\lambda_2 = \lambda_1 + 1}$

(i) let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \lambda_1, f_3 = \frac{\lambda_0 - 3 \lambda_1 - 2}{2}, f_4 = \lambda_1 + 1$$

Note the following,

□  $\lambda_0 - 3 \lambda_1 - 2$  is even (Lemma 2.4), i.e.  $f_3$  is integer.

□  $f_3 \geq 0$  iff  $\lambda_1 \leq \frac{\lambda_0 - 2}{3}$

$$b^* = 2\lambda_0$$

therefore,  $b \geq b^*$  (Lemma 2.3)

(ii) let  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4$ , with

$$f_1 = \frac{\lambda_0 - 2}{3}, \quad f_2 = \frac{3\lambda_1 - \lambda_0 + 2}{6}, \quad f_4 = \frac{\lambda_0 + 1}{3}$$

Note that,

- $\lambda_0 - 2$  and  $\lambda_0 + 1$  are divisible by 3, i.e.  $f_1$  and  $f_4$  are integers.
- $f_2 \geq 0$  iff  $\lambda_1 \geq \frac{\lambda_0 - 2}{3}$
- $\lambda_0 = 2 \pmod{6}, \lambda_1 = 0 \pmod{2}$  for  $f_2$  to be integer.
- $\lambda_0 = 5 \pmod{6}, \lambda_1 = 1 \pmod{2}$  for  $f_2$  to be integer.
- $b^* = \frac{2}{3} (2\lambda_0 + 3\lambda_1 + 2)$
- $b \geq \frac{2}{3} (2\lambda_0 + 3\lambda_1 + 2)$  (Lemma 2.1)

Hence,  $b \geq b^*$

The cases (i), (ii) cover all the possibilities for,

$$\lambda_0 = 2 \pmod{3}, \lambda_2 = \lambda_1 + 1$$

CASE 1, CASE 2, and CASE 3 cover all the possible values of  $\lambda_0$  ( $= 0, 1, 2, 3, \dots$ ), of  $\lambda_1$  ( $= 0, 1, 2, 3, \dots$ ) and of  $\lambda_2$  ( $= \lambda_1 + \delta, \delta = -1, 0, 1$ ) and in each case it is shown that an arbitrary design  $D$  is either equivalent or  $S$ -inadmissible with respect to a design  $D^*$  which is the union of copies of generator designs from Table 4.2. Furthermore, it can be noted that none of the

designs in Table 4.2 is equivalent to any other design in the table, and all the designs in Table 4.2, except the design  $D_8$ , are used in the construction of design  $D$  for all possible combinations of parameters  $\lambda_0, \lambda_1, \lambda_2$ . The design  $D_8$  in Table 4.2 is a generator design and since it is not equivalent or  $S$ -inadmissible with respect to the union of any number of copies of any combination of other generator designs, it has to be included in the list.

It is, therefore, concluded that the designs in Table 4.2 constitute a minimal complete class of generator designs for  $k = 3, v = 4$ . Hence, the Theorem is proved.

#### 4.3 MCCGD FOR $k = 3, v = 6$ :

##### THEOREM 4.3

For  $k = 3$  and  $v = 6$  the minimal complete class of generator designs is given in Table 4.3.

TABLE 4.3 MCCGD FOR  $k = 3, v = 6$ .

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE	
$D_1$	2	0	1	0	0	1	0	1	4
				0				2	5
								3	6



TABLE 4.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
D <sub>2</sub>	6	12	1	0	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
D <sub>3</sub>	6	6	2	6	2	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0
D <sub>4</sub>	7	3	3	3	1	1	1	0	0	0	1	1	2	2	1	3	4	3	4	5
D <sub>5</sub>	9	9	3	9	3	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0
D <sub>6</sub>	10	0	5	0	0	2	2	1	1	1	1	1	2	2	2	2	2	2	3	3
				0																
D <sub>7</sub>	12	6	5	6	2	1	2	0	0	0	0	0	0	0	1	1	1	1	1	1
				0																
D <sub>8</sub>	15	3	7	3	1	2	3	0	0	0	1	1	1	1	1	1	1	1	1	1
				0																
D <sub>9</sub>	15	9	6	3	2	2	2	3	3	0	0	0	0	0	1	1	1	1	1	1
				3																

TABLE 4.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE								
D <sub>10</sub>	18	0	9	0	0	3	4	1	1	1	1	1	1	1	1	
								2	2	2	3	3	3	4	4	5
								4	5	6	4	5	6	5	6	6
								2	2	2	2	2	2	3	3	3
								3	3	3	4	4	5	4	4	5
								4	5	6	5	6	6	5	6	6

4.4 MCCGD FOR k = 3, v = 8 :

THEOREM 4.4

For k = 3 and v = 8 the minimal complete class of generator designs is given in Table 4.4.

TABLE 4.4 MCCGD FOR k = 3, v = 8.

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE								
D <sub>1</sub>	8	16	1	0	2	0	0	0	0	0	0	0	0	0	0	
								0	0	0	0	0	0	0	0	
								1	2	3	4	5	6	7	8	
D <sub>2</sub>	12	4	4	4	1	1	1	0	0	0	0	1	1	1	2	2
								1	3	5	7	3	4	6	3	4
								2	4	6	8	5	8	7	7	6
								2	3	4						
								5	6	5						
								8	8	7						
D <sub>3</sub>	12	12	3	12	3	1	0	0	0	0	0	0	0	0	0	0
								1	1	1	2	2	3	5	5	5
								2	3	4	3	4	4	6	7	8
								0	0	0						
								6	6	7						
								7	8	8						

TABLE 4.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE							
D <sub>4</sub>	16	16	4	16 0	4	0	1	0 0 0 0 0 0 0 0 0 1 1 1 1 2 2 2 2 3 5 6 7 8 5 6 7 8 5 ----- 0 0 0 0 0 0 0 3 3 3 4 4 4 4 6 7 8 5 6 7 8							
D <sub>5</sub>	18	14	5	10 2	3	1	1	0 0 0 0 0 0 0 0 0 1 5 2 2 2 3 3 3 6 4 6 4 7 8 6 4 8 7 ----- 0 0 0 1 1 1 2 4 4 7 0 0 2 3 5 3 5 6 8 1 5 6 7 8 5 7 8							
D <sub>6</sub>	19	9	6	7 1	2	2	1	0 0 0 0 0 0 0 0 1 1 1 2 5 5 6 7 0 2 2 3 3 6 7 8 8 4 4 ----- 1 1 1 2 2 2 3 3 4 3 5 7 3 5 6 5 6 5 4 6 8 4 7 8 8 7 8 ----- 4 6 7							
D <sub>7</sub>	23	13	7	11 1	3	1	2	0 0 0 0 0 0 0 0 0 1 1 1 2 2 2 3 3 3 5 6 7 6 7 8 5 7 8 ----- 0 0 0 1 1 1 1 2 2 4 5 0 2 3 7 4 3 4 6 8 4 5 6 8 8 7 8 ----- 2 3 3 4 4 5 4 6 6 5 6 5 8 7 7							
D <sub>8</sub>	23	21	6	19 1	5	1	1	0 0 0 0 0 0 0 0 0 1 1 1 1 1 2 2 2 2 4 5 6 7 8 4 5 6 7 ----- 0 0 0 0 0 0 0 0 0 2 3 3 3 3 3 4 4 5 8 4 5 6 7 8 5 6 7 ----- 0 0 1 5 4 6 0 2 6 7 7 8 3 8 8							

TABLE 4.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE								
D <sub>9</sub>	24	0	9	0	0	2	3	1	1	1	1	1	1	1	1	1
								2	2	3	3	4	4	5	5	6
								5	6	7	8	6	7	8	7	8
D <sub>10</sub>	25	11	8	5	2	2	2	0	0	0	0	0	0	0	0	1
								4	4	5	5	6	0	0	0	2
								6	7	7	8	8	1	2	3	5
D <sub>11</sub>	30	10	10	6	2	2	3	1	1	1	1	1	1	2	2	2
								2	3	3	4	4	7	3	3	6
								6	5	6	7	8	8	7	8	7
D <sub>12</sub>	32	0	12	0	0	4	3	0	0	0	0	0	0	0	0	1
								1	1	2	2	3	3	0	0	2
								5	6	7	5	6	7	4	8	7

TABLE 4.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
								4	5	5	5	6								
								6	6	6	7	7								
								7	7	8	8	8								
D <sub>13</sub>	38	10	13	6	2	4	3	0	0	0	0	0	0	0	0	0	0	0	1	
				2				1	1	2	5	5	6	0	0	2				
								2	3	3	6	7	7	4	8	3				
								1	1	1	1	1	1	1	1	1				
								2	2	3	4	4	5	5	6	6				
								3	4	4	8	5	6	7	7	8				
								1	2	2	2	2	2	2	2	2				
								7	3	4	4	5	6	6	5	7				
								8	4	5	7	6	7	8	8	8				
								3	3	3	3	3	3	3	3	4	4			
								4	4	5	5	5	6	7	5	6				
								6	6	7	8	7	8	8	8	7				
								4	5											
								7	6											
								8	8											

4.5 MCCGD FOR  $k = 3, v = 10$  :

THEOREM 4.5

For  $k = 3$  and  $v = 10$  the minimal complete class of generator designs is given in Table 4.5.

TABLE 4.5 MCCGD FOR  $k = 3, v = 10$  .

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
D <sub>1</sub>	10	20	1	0	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
				10				0	0	0	0	0	0	0	0	0	0	0	0	0
								1	2	3	4	5	6	7	8	9				

TABLE 4.5 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE
								0 0 10
D <sub>2</sub>	14	12	3	8 2	2	1	0	0 0 0 0 0 0 0 0 0 1 1 2 3 6 6 7 8 0 2 3 4 4 7 8 9 9 5 ----- 0 1 2 6 7 0 4 3 9 8 10 5 5 10 10
D <sub>3</sub>	20	20	4	20 0	4	1	0	0 0 0 0 0 0 0 0 0 1 1 1 1 2 2 2 3 3 2 3 4 5 3 4 5 4 5 ----- 0 0 0 0 0 0 0 0 0 4 6 6 6 6 7 7 7 8 5 7 8 9 10 8 9 10 9 ----- 0 0 8 9 10 10
D <sub>4</sub>	25	5	7	5 0	1	2	1	0 0 0 0 0 1 1 1 1 1 3 6 8 5 2 3 4 5 2 4 7 9 10 3 4 5 6 ----- 1 1 2 2 2 2 2 3 3 7 9 3 4 4 7 8 5 6 8 10 5 5 6 9 10 8 9 ----- 3 4 4 5 6 6 6 7 7 9 7 7 8 8 10 8 10 9 10 9 10
D <sub>5</sub>	25	15	6	15 0	3	1	1	0 0 0 0 0 0 0 0 0 1 1 1 2 2 2 3 3 4 3 5 8 4 7 8 5 9 6 ----- 0 0 0 0 0 0 1 1 1 4 5 6 6 7 7 2 4 6 10 10 9 10 8 9 10 9 7 ----- 2 2 3 3 4 5 8 3 5 4 7 5 6 9 6 9 8 10 7 8 10

TABLE 4.5 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE													
D <sub>6</sub>	25	25	5	25	5	0	1	0	0	0	0	0	0	0	0	0	0	0	0		
								1	1	1	1	1	2	2	2	2	2	2	2		
								6	7	8	9	10	6	7	8	9	10	6	7	8	9
								0	0	0	0	0	0	0	0	0	0	0	0	0	0
D <sub>7</sub>	30	0	9	0	0	2	2	1	1	1	1	1	1	1	1	1	1	1	1		
								2	2	2	3	4	5	6	7	9	10	10	10	10	
								3	6	7	8	5	8	9	10	10	10	10	10	10	10
								2	2	2	2	2	2	2	2	3	3	3	3	3	3
D <sub>8</sub>	30	10	8	10	2	1	2	0	0	0	0	0	0	0	0	0	0	0	0		
								1	1	2	2	3	3	4	4	5	5	5	5	5	
								6	7	8	9	10	6	7	8	9	10	6	7	8	9
								0	1	1	1	1	1	1	1	2	2	2	2	2	2
D <sub>9</sub>	32	16	8	14	3	2	1	0	0	0	0	0	0	0	0	0	0	0	0		
								1	1	1	2	2	3	5	5	5	5	5	5	5	
								2	3	4	3	4	4	6	7	8	8	8	8	8	8
								0	1	1	1	1	1	1	1	2	2	2	2	2	2
D <sub>9</sub>	32	16	8	14	3	2	1	0	0	0	0	0	0	0	0	0	0	0	0		
								1	1	1	2	2	3	5	5	5	5	5	5	5	
								2	3	4	3	4	4	6	7	8	8	8	8	8	8
								0	1	1	1	1	1	1	1	2	2	2	2	2	2

TABLE 4.5 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE
								0 0 0 0 0 0 1 1 1 6 6 7 7 9 0 2 3 4 9 8 9 8 10 10 5 5 6 ----- 1 1 2 2 2 2 3 3 3 7 8 3 4 5 8 4 7 8 9 10 6 7 9 10 5 10 9 ----- 4 4 6 6 6 5 8 7 7 9 10 9 8 10 10
D <sub>10</sub>	32	26	7	24 1	5	1	1	0 0 0 0 0 0 0 0 0 1 1 1 1 1 2 2 2 2 3 4 5 8 9 4 5 7 8 ----- 0 0 0 0 0 0 0 0 0 2 3 3 3 3 4 4 4 5 9 5 7 9 10 6 7 9 6 ----- 0 0 0 0 0 0 0 1 1 5 6 6 6 7 8 0 2 6 7 8 9 10 8 10 10 10 7 ----- 2 3 4 5 7 3 4 5 8 9 6 8 10 9 10
D <sub>11</sub>	37	11	10	9 1	2	2	2	0 0 0 0 0 0 0 0 0 1 1 2 2 3 4 6 7 8 3 5 4 7 5 6 9 8 9 ----- 0 1 1 1 1 1 1 1 1 0 2 2 3 4 4 5 6 6 10 8 8 10 9 9 10 7 7 ----- 2 2 2 2 2 2 3 3 3 3 3 4 5 5 7 4 4 5 6 6 10 9 9 10 8 8 10 ----- 3 3 4 4 4 5 5 6 7 7 7 5 5 6 6 6 9 8 9 9 7 7 10 8 8 10 10 ----- 8 9 10



TABLE 4.5 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE																						
D <sub>12</sub>	37	21	9	19	4	1	2	0	0	0	0	0	0	0	0	0	0	0	0											
								1	1	1	1	2	2	2	2	2	3													
								6	7	8	9	6	8	9	10	7														
								0	0	0	0	0	0	0	0	0														
								3	3	3	4	4	4	4	5	5														
								8	9	10	6	7	8	5	6	7														
D <sub>13</sub>	44	32	10	28	6	1	2	0	0	0	0	0	0	0	0	0	0	0	0											
								1	1	1	1	1	1	2	2	2														
								6	7	8	9	10	6	6	7	8														
								0	0	0	0	0	0	0	0	0														
								2	2	2	3	3	3	3	3	3														
								9	10	7	6	7	8	9	10	8														
D <sub>14</sub>	49	17	13	13	3	2	3	0	0	0	0	0	0	0	0	0	0	0												
								1	1	1	2	2	2	3	3	3														
								6	7	8	7	8	9	8	9	10														
								0	0	0	0	0	0	1	1	1														
								4	4	4	5	0	0	2	2	3														
								6	7	9	6	5	10	9	10	6														

TABLE 4.5 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE									
								1	1	1	1	1	1	1	2	2	
								3	4	4	5	5	6	8	3	3	
								7	8	9	10	7	9	10	6	7	
								2	2	2	2	2	2	3	3	3	
								4	4	5	5	6	8	4	4	5	
								8	9	10	6	7	10	10	6	8	
								3	3	3	4	4	4	4	5	5	
								5	7	9	5	5	6	7	6	7	
								9	8	10	7	10	8	10	8	9	
								5	6	6	7						
								8	7	9	8						
								9	10	10	9						
D <sub>15</sub>	56	18	15	12	3	3	3	0	0	0	0	0	0	0	0	0	
				3				1	2	3	4	4	4	8	5	5	
								7	7	7	8	5	10	9	9	8	
								0	0	0	0	0	0	1	1	1	
								6	6	6	0	0	0	2	2	2	
								10	9	10	1	2	3	6	8	9	
								1	1	1	1	1	1	1	1	1	
								3	3	3	4	4	4	5	5	5	
								10	6	8	7	6	8	7	9	10	
								1	2	2	2	2	2	2	2	2	
								9	3	3	3	4	4	4	5	5	
								10	9	10	7	10	7	9	6	8	
								2	2	3	3	3	3	3	3	3	
								5	8	4	4	4	5	5	5	8	
								6	10	6	8	9	7	6	10	9	
								4	4	4	5	6	6	6	6	7	
								5	5	6	8	7	7	7	8	8	
								7	10	9	9	8	8	9	10	10	
								7	7								
								9	9								
								10	10								

CHAPTER 5

5.1 MCCGD FOR  $k = 4, v = 4$  :

THEOREM 5.1

For  $k = 4$  and  $v = 4$  the minimal complete class of generator designs is given in Table 5.1.

TABLE 5.1 MCCGD FOR  $k = 4, v = 4$  .

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE				
$D_1$	1	0	1	0	0	1	1	1				
				0				2				
				0				3				
				0				4				
$D_2$	2	4	1	0	2	1	0	0	0			
				2				0	0			
				0				1	3			
				0				2	4			
$D_3$	4	8	2	0	4	0	1	0	0	0	0	0
				4				0	0	0	0	0
				0				1	1	2	2	
				0				3	4	3	4	
$D_4$	4	4	3	4	3	2	2	0	0	0	0	
				0				1	1	1	2	
				0				2	2	3	3	
				0				3	4	4	4	

PROOF:

The procedure for proving the class to be minimal class is same as discussed in Section 4.1. The lower bound on B is, however, obtained by using the

FACT 2.6 and FACT 2.7, i.e. by finding an index  $j_0$  which minimizes  $|N - M_{j_0}|$ , where  $N$  is the total number of TT-pairs in the design and  $M_{j_0}$  is the number of TT-pairs in the  $j_0$  th BB combination with  $b_{j_0}$  blocks in it.

CASE 1:  $\lambda_0 = 0 \pmod{3}$

According to the Lemma 2.7,  $\lambda_2 = \lambda_1$ , so there is no need to consider the cases when  $\lambda_2 = \lambda_1 \pm 1$ . Further note that,  $N = 6\lambda_1$  (FACT 2.3).

(a) Consider  $3\lambda_1 - 2\lambda_0 \geq 0$ .

Let  $D^* = f_1 D_1 \cup f_3 D_3$ , with

$$f_1 = \frac{3\lambda_1 - 2\lambda_0}{3}, \quad f_3 = \lambda_0/3$$

Since  $\lambda_0 = 0 \pmod{3}$ ,  $f_1$  and  $f_3$  are both integers.

Furthermore,

$$b^* = \frac{3\lambda_1 + 2\lambda_0}{3}$$

By Lemma 2.1,

$$b \geq \frac{3\lambda_1 + 2\lambda_0}{3} = b^*$$

Hence,  $D$  is either equivalent or  $S$ -inadmissible with respect to  $D^*$ .

(b) Consider  $3\lambda_1 - 2\lambda_0 < 0$ .

Note that  $6\lambda_1 - \lambda_0 = L \pmod{9}$ ,  $L = 0, 3, 6$ .

We will take these cases separately.

$$(i) \quad 6\lambda_1 - \lambda_0 = 0 \pmod{9} = 9m, \quad m = 0, 1, 2, \dots$$

Note that  $N = \lambda_0 + 9m$  (FACT 2.3)

Let  $D^* = f_2 D_2 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_2 = \frac{2\lambda_0 - 3\lambda_1}{9}, \quad f_3 = \frac{6\lambda_1 - \lambda_0}{9}, \quad f_4 = f_2$$

Note that  $f_2, f_3$  are non-negative integers.

$$\text{so, } b^* = \lambda_0 + m$$

take,

$$j_0 = \begin{cases} m & , \text{ if } \lambda_0 \text{ is odd} \\ m + 1 & , \text{ if } \lambda_0 \text{ is even} \end{cases}$$

$$\text{So, } M_{j_0} = \lambda_0 + 9m \quad (\text{FACT 2.6})$$

$$\text{and } b_{j_0} = \lambda_0 + m$$

Therefore,

$$b \geq b_{j_0} = b^* \quad (\text{FACT 2.7})$$

$$(ii) \quad 6\lambda_1 - \lambda_0 = 3 \pmod{9} = 9m + 3, \quad m = 0, 1, 2, \dots$$

Note that,  $N = \lambda_0 + 9m + 3$  (FACT 2.3)

Take,

$$j_0 = \begin{cases} m & , \text{ if } \lambda_0 \text{ is odd} \\ m + 1 & , \text{ if } \lambda_0 \text{ is even} \end{cases}$$

So,  $M_{j_0} = \lambda_0 + 9m$  is closest to  $N$ , with  $b_{j_0} = \lambda_0 + m$

But  $b \geq b_{j_0} = \lambda_0 + m$ ,

Now consider  $D^*$  with  $(b^*, \lambda_0^*, \lambda_1^*, \lambda_2^*)$ , where  $\lambda_0^* = \lambda_0$ ,  $\lambda_1^* = \lambda_1 + 1$ , and  $\lambda_2^* = \lambda_2$ . Then

$$6\lambda_1^* - \lambda_0^* = 6\lambda_1 - \lambda_0 + 6 = 0 \pmod{9}$$

Therefore, from Subsection b(i) above,  $b^* = \lambda_0 + m$ .

Hence,  $D$  is with  $(b \geq \lambda_0 + m, \lambda_0, \lambda_1, \lambda_2)$  and  $D^*$  is with  $(b^* = \lambda_0 + m, \lambda_0, \lambda_1 + 1, \lambda_1 + 1)$ . That means any  $D$  with  $\lambda_0, \lambda_1, \lambda_2$  and  $6\lambda_1 - \lambda_0 = 3 \pmod{9}$  is  $S$ -inadmissible with respect to  $D^*$ .

$$(iii) \quad 6\lambda_1 - \lambda_0 = 6 \pmod{9} = 9m + 6, \quad m = 0, 1, 2, \dots$$

Note that  $N = 6\lambda_1 = \lambda_0 + 9m + 6$  (FACT 2.3)

Let,  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = 1, \quad f_2 = f_4 = \frac{2\lambda_0 - 3\lambda_1 + 3}{9}, \quad f_3 = \frac{6\lambda_1 - \lambda_0 - 6}{9}$$

Note that  $f_2, f_3$ , and  $f_4$  are non-negative integers.

Now,  $b^* = \lambda_0 + m + 1$ .

Take,

$$j_0 = \begin{cases} m + 1, & \text{if } \lambda_0 \text{ is odd} \\ m + 2, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

So,  $M_{j_0} = \lambda_0 + 9m + 9$ , which is closest to  $N$ , and

$$b_{j_0} = \lambda_0 + m + 1 \quad (\text{FACT 2.6})$$

$$\text{Now, } b \geq b_{j_0} = b^* \quad (\text{FACT 2.7})$$

CASE 2 :  $\lambda_0 = 1 \pmod{3}$

According to Lemma 2.7,  $\lambda_2 = \lambda_1 + 1$  and  
therefore,  $N = 6\lambda_1 + 4$  (FACT 2.3)

(a)  $3\lambda_1 - 2\lambda_0 \geq -8$ . Therefore,  $N = 4\lambda_0 + 6m - 12$ .

Let  $D^* = f_1 D_1 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = \frac{3\lambda_1 - 2\lambda_0 + 8}{3}, \quad f_3 = 1, \quad f_4 = \frac{\lambda_0 - 4}{3}$$

Note that,  $f_1$  and  $f_4$  are non-negative integers.

$$\text{So, } b^* = \frac{4\lambda_0 - 4}{3} + m = 5m$$

Take,

$$j_0 = \begin{cases} \frac{\lambda_0 - 1}{3} + m, & \text{if } \lambda_0 \text{ is odd} \\ \frac{\lambda_0 + 2}{3} + m, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

So,

$M_{j_0} = 4\lambda_0 + 9m - 3$  is closest to  $N$  and

$$b_{j_0} = 5m + 1 \quad (\text{FACT 2.6})$$

Therefore,  $b \geq b_{j_0} > b^*$  (FACT 2.7)

$$(b) 3\lambda_1 - 2\lambda_0 < -8$$

Consider the cases when,

$$6\lambda_1 - \lambda_0 = L \pmod{9}, L = 2, 5, 8.$$

$$(i) 6\lambda_1 - \lambda_0 = 2 \pmod{9} = 9m + 2, m = 0, 1, 2, \dots$$

$$\text{Then, } N = \lambda_0 + 9m + 6$$

Let,  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = 1, f_2 = \frac{2\lambda_0 - 3\lambda_1 - 5}{9}, f_3 = \frac{2\lambda_0 - 3\lambda_1 + 4}{9}, f_4 = \frac{6\lambda_1 - \lambda_0 - 2}{9}$$

Note that,  $f_2, f_3, f_4$  are non-negative integers.

$$\text{So, } b^* = \lambda_0 + m + 1.$$

Take,

$$j_0 = \begin{cases} m + 1, & \text{if } \lambda_0 \text{ is odd} \\ m + 2, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

Therefore,  $M_{j_0} = \lambda_0 + 9m + 9$  is closest to  $N$ , with

$$b_{j_0} = \lambda_0 + m + 1 \quad (\text{FACT 2.6})$$

So,

$$b \geq b_{j_0} = b^* \quad (\text{FACT 2.7})$$

$$(ii) 6\lambda_1 - \lambda_0 = 5 \pmod{9} = 9m + 5, m = 0, 1, 2, \dots$$

Note that,  $N = \lambda_0 + 9m + 9$ .

Take,  $D^* = f_2 D_2 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_2 = \frac{2\lambda_0 - 3\lambda_1 - 8}{9}, f_3 = \frac{2\lambda_0 - 3\lambda_1 + 1}{9}, f_4 = \frac{6\lambda_1 - \lambda_0 + 4}{9}$$

It can be seen that  $f_2, f_3, f_4$  are non-negative integers.



Also,  $b^* = \lambda_0 + m + 1$

Take,

$$j_0 = \begin{cases} m + 1, & \text{if } \lambda_0 \text{ is odd} \\ m + 2, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

Then,  $M_{j_0} = \lambda_0 + 9m + 9$  is equal to  $N$  with

$$b_{j_0} = \lambda_0 + m + 1 \quad (\text{FACT 2.6})$$

So,  $b \geq b_{j_0} = b^*$  (FACT 2.7)

$$(iii) \quad 6\lambda_1 - \lambda_0 = 8 \pmod{9} = 9m + 8, \quad m = 0, 1, 2, \dots$$

Take,  $j_0$  as in Subsection (ii) above.

So,  $b \geq b_{j_0} = \lambda_0 + m + 1$

Consider,  $D^*$  with  $(b^*, \lambda_0^* = \lambda_0, \lambda_1^* = \lambda_1 + 1, \lambda_2^* = \lambda_2)$

Now,  $6\lambda_1^* - \lambda_0^* = 6\lambda_1 - \lambda_0 + 6 = 9m + 14 = 5 \pmod{9}$

Therefore, by Subsection (ii) above.

$$b^* = \lambda_0 + m + 1$$

Hence,  $D$  is with  $(b \geq \lambda_0 + m + 1, \lambda_0, \lambda_1, \lambda_2)$  and  $D^*$  is with  $(b^* = \lambda_0 + m + 1, \lambda_0, \lambda_1 + 1, \lambda_1 + 1)$ . Therefore, any  $D$  with  $\lambda_0, \lambda_1, \lambda_2$  and  $6\lambda_1 - \lambda_0 = 5 \pmod{9}$  is  $S$ -inadmissible with respect to  $D^*$ .

CASE 3 :  $\lambda_0 = 2 \pmod{3}$

According to Lemma 2.7,  $\lambda_2 = \lambda_1 - 1$ , and therefore,  $N = 6\lambda_1 - 4$  (FACT 2.3)

$$(a) \quad 3\lambda_1 - 2\lambda_0 \geq -1$$

$$N = 4\lambda_0 + 6m - 6$$

Let,  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_4 D_4$ , with

$$f_1 = \frac{3\lambda_1 - 2\lambda_0 + 1}{3}, \quad f_2 = 1, \quad f_4 = \frac{\lambda_0 - 2}{3}$$

Note that,  $f_1$  and  $f_4$  are non-negative integers.

$$\text{So,} \quad b^* = \frac{4\lambda_0 + 3m - 2}{3}$$

Take,

$$j_0 = \begin{cases} \frac{\lambda_0 - 2}{3} + m, & \text{if } \lambda_0 \text{ is odd} \\ \frac{\lambda_0 + 1}{3} + m, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

So,  $M_{j_0} = 4\lambda_0 + 9m - 6$ , and

$$b_{j_0} = \frac{4\lambda_0 + 3m - 2}{3}$$

$$\text{So,} \quad b \geq b_{j_0} = b^*$$

$$(b) \quad 3\lambda_1 - 2\lambda_0 < -1$$

Consider  $6\lambda_1 - \lambda_0 = L \pmod{9}$ ,  $L = 1, 4, 7$ .

$$(i) \quad 6\lambda_1 - \lambda_0 = 1 \pmod{9} \geq 10 \quad (= 9m + 1, \\ m = 0, 1, 2, \dots)$$

Note that,  $N = \lambda_0 + 9m + 6$

Let,  $D^* = f_1 D_1 \cup f_2 D_2 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_1 = 1, f_2 = \frac{2\lambda_0 - 3\lambda_1 + 11}{9}, f_3 = \frac{2\lambda_0 - 3\lambda_1 + 2}{9}, f_4 = \frac{6\lambda_1 - \lambda_0 - 10}{9}$$

So,  $b^* = \lambda_0 + m$

Take,

$$j_0 = \begin{cases} m + 1, & \text{if } \lambda_0 \text{ is odd} \\ m + 2, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

So,  $M_{j_0} = \lambda_0 + 9m + 9$  is closest to  $N$ , and

$$b_{j_0} = \lambda_0 + m + 1$$

Therefore,

$$b \geq b_{j_0} > b^*$$

$$(ii) \quad 6\lambda_1 - \lambda_0 = 4 \pmod{9} = 9m + 4, \quad m = 0, 1, 2, \dots$$

Therefore,  $N = \lambda_0 + 9m - 3$

Let,  $D^* = f_2 D_2 \cup f_3 D_3 \cup f_4 D_4$ , with

$$f_2 = \frac{2\lambda_0 - 3\lambda_1 + 8}{9}, f_3 = \frac{2\lambda_0 - 3\lambda_1 - 1}{9}, f_4 = \frac{6\lambda_1 - \lambda_0 - 4}{9}$$

This gives,  $b^* = \lambda_0 + m$

Take,

$$j_0 = \begin{cases} m, & \text{if } \lambda_0 \text{ is odd} \\ m + 1, & \text{if } \lambda_0 \text{ is even} \end{cases}$$

So,  $M_{j_0} = \lambda_0 + 9m$ , with

$$b_{j_0} = \lambda_0 + m$$

Therefore,  $b \geq b_{j_0} = b^*$ .

$$(iii) \quad 6\lambda_1 - \lambda_0 = 7 \pmod{9} = 9m + 7, \quad m = 0, 1, 2, \dots$$

Note that,  $N = \lambda_0 + 9m + 3$

Take,

$$j_0 = \begin{cases} m & , \text{ if } \lambda_0 \text{ is odd} \\ m + 1 & , \text{ if } \lambda_0 \text{ is even} \end{cases}$$

So that,  $M_{j_0} = \lambda_0 + 9m$  and  $b_{j_0} = \lambda_0 + m$

Therefore,  $b \geq \lambda_0 + m$

Let,  $D^*$  is a design with  $(b^* \lambda_0^* = \lambda_0, \lambda_1^* = \lambda_1 + 1, \lambda_2^* = \lambda_2)$

Here,  $6\lambda_1^* - \lambda_0^* = 6\lambda_1 - \lambda_0 + 6 = 4 \pmod{9}$

This case is covered in b(ii) above, and hence,

$$b^* = \lambda_0 + m$$

Therefore,  $D$  is with  $(b \geq \lambda_0 + m, \lambda_0, \lambda_1, \lambda_2)$  and

$D^*$  is with  $(b^* = \lambda_0 + m, \lambda_0, \lambda_1 + 1, \lambda_2^*)$

Hence,  $D$  has smaller  $\lambda_1$  value than the same for  $D^*$ .

Therefore, any design  $D$  for which  $6\lambda_1 - \lambda_0 = 7 \pmod{9}$ , is  $S$ -inadmissible with respect to  $D^*$ .

$$(iv) \quad 6\lambda_1 - \lambda_0 = 1$$

Therefore,  $\lambda_0 = 1 \pmod{6} = 6m + 1, \quad m = 0, 1, 2, \dots$

From Lemma 2.2,

$$r = \frac{\lambda_0 + 3\lambda_1 - 2}{3} = 2m + \lambda_1 - 1/3,$$

which is not an integer.

Therefore, such designs cannot exist.

Note that  $6\lambda_1 - \lambda_0$  is not allowed to take all possible values in the above cases. It will be shown below that for all such uncovered cases any design  $D$  will be  $S$ -inadmissible with respect to a design  $D^*$ , which is union of copies of designs from Table 5.1.

CASE 1: Let  $\lambda_0$  is even.

From FACT 2.7 ,  $b \geq b_{j_0} = \lambda_0 + j_0 - 1, j = 1, 2, \dots$   
 $\dots\dots\dots, \lceil \lambda_0/3 \rceil + 1$

Therefore,  $b \geq \lambda_0$

(i) Let  $\lambda_0 = 0 \pmod{6}$  and  $6\lambda_1 - \lambda_0 < 0 \Rightarrow \lambda_1 < \lambda_0/6$

$$\text{Let, } D^* = \frac{\lambda_0}{6} D_2 \cup \frac{\lambda_0}{6} D_3$$

$$\text{So, } b^* = \lambda_0, \lambda_0^* = \lambda_0, \lambda_1^* = \lambda_0/6$$

Therefore,  $D^*$  is with  $(b^* = \lambda_0, \lambda_0^* = \lambda_0, \lambda_1^* = \lambda_0/6, \lambda_2^*)$

Whereas,  $D$  is with  $(b \geq \lambda_0, \lambda_0, \lambda_1 < \lambda_0/6, \lambda_2)$

So,  $D$  has smaller  $\lambda_1$  than  $\lambda_1^*$  of  $D^*$ . Hence,  $D$  is  $S$ -inadmissible with respect to  $D^*$ .

(ii) Let  $\lambda_0 = 2 \pmod{6}, 6\lambda_1 - \lambda_0 < 4 \Rightarrow$

$$\lambda_1 < \frac{\lambda_0 + 4}{6}$$

$$\text{Let, } D^* = \frac{\lambda_0 + 4}{6} D_2 \cup \frac{\lambda_0 - 2}{6} D_3$$

$$\text{So, } b^* = \lambda_0, \lambda_0^* = \lambda_0, \text{ and } \lambda_1^* = \frac{\lambda_0 + 4}{6}$$

So,  $D^*$  is with  $(b^* = \lambda_0, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 4}{6}, \lambda_2^*)$

and  $D$  is with  $(b \geq \lambda_0, \lambda_0, \lambda_1 < \frac{\lambda_0 + 4}{6}, \lambda_2)$

Hence,  $D$  is  $S$ -inadmissible with respect to  $D^*$ .

(iii) Let  $\lambda_0 = 4 \pmod{6}$ ,  $6\lambda_1 - \lambda_0 < -4 \Rightarrow$

$$\lambda_1 < \frac{\lambda_0 - 4}{6}$$

$$\text{Let, } D^* = \frac{\lambda_0 - 4}{6} D_2 \cup \frac{\lambda_0 + 2}{6} D_3$$

Therefore,  $b^* = \lambda_0, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 - 4}{6}$

So,  $D^*$  is with  $(b^* = \lambda_0, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 - 4}{6}, \lambda_2^*)$

and  $D$  is with  $(b \geq \lambda_0, \lambda_0, \lambda_1 < \frac{\lambda_0 - 4}{6}, \lambda_2)$

Hence,  $D$  is  $S$ -inadmissible with respect to  $D^*$ .

CASE 2: Let  $\lambda_0$  is odd

From FACT 2.7,  $b \geq b_{j_0} = \lambda_0 + j_0$ ,  $j_0 = 1, 2, \dots, \lfloor \lambda_0/3 \rfloor$

Therefore,  $b \geq \lambda_0 + 1$ .

(i) Let  $\lambda_0 = 1 \pmod{6}$  and  $6\lambda_1 - \lambda_0 < 5 \Rightarrow$

$$\lambda_1 < \frac{\lambda_0 + 5}{6}$$

$$\text{Let, } D^* = \frac{\lambda_0 - 7}{6} D_2 \cup \frac{\lambda_0 - 1}{6} D_3 \cup D_4$$

Then,  $b^* = \lambda_0 + 1, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 5}{6}$

So,  $D^*$  is with  $(b^* = \lambda_0 + 1, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 5}{6}, \lambda_2^*)$

and  $D$  is with  $(b \geq \lambda_0 + 1, \lambda_0, \lambda_1 < \frac{\lambda_0 + 5}{6}, \lambda_2)$

Hence,  $D$  is  $S$ -inadmissible with respect to  $D^*$ , for  
 $6\lambda_1 - \lambda_0 < 5$ .

(ii) Let  $\lambda_0 = 3 \pmod{6}$  and  $6\lambda_1 - \lambda_0 < 9 \Rightarrow$   
 $\lambda_1 < \frac{\lambda_0 + 9}{6}$ .

Let,  $D^* = \frac{\lambda_0 - 3}{6} D_2 \cup \frac{\lambda_0 - 3}{6} D_3 \cup D_4$

So,  $b^* = \lambda_0 + 1, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 9}{6}$ , i.e.,

$D^*$  is with  $(b^* = \lambda_0 + 1, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 9}{6}, \lambda_2^*)$

and  $D$  is with  $(b \geq \lambda_0 + 1, \lambda_0, \lambda_1 < \frac{\lambda_0 + 5}{6}, \lambda_2)$

So,  $D$  is  $S$ -inadmissible with respect to  $D^*$  for  
 $6\lambda_1 - \lambda_0 < 9$ .

(iii) Let  $\lambda_0 = 5 \pmod{6}$  and  $6\lambda_1 - \lambda_0 < 13 \Rightarrow$   
 $\lambda_1 < \frac{\lambda_0 + 13}{6}$

Let  $D^* = \frac{\lambda_0 + 1}{6} D_2 \cup \frac{\lambda_0 - 5}{6} D_3 \cup D_4$

So,  $b^* = \lambda_0 + 1, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 13}{6}$

So,  $D^*$  is with  $(b^* = \lambda_0 + 1, \lambda_0^* = \lambda_0, \lambda_1^* = \frac{\lambda_0 + 13}{6}, \lambda_2^*)$

and  $D$  is with  $(b \geq \lambda_0 + 1, \lambda_0, \lambda_1 < \frac{\lambda_0 + 13}{6}, \lambda_2)$

Hence,  $D$  is  $S$ -inadmissible with respect to  $D^*$  for  
 $6\lambda_1 - \lambda_0 < 13$ .

This makes the three cases considered before exhaustive. Hence proof of the theorem is complete.

5.2 MCCGD FOR  $k = 4, v = 6$  :

THEOREM 5.2

For  $k = 4$  and  $v = 6$  the minimal complete class of generator designs is given in Table 5.2

TABLE 5.2 MCCGD FOR  $k = 4, v = 6$ .

DESIGN	b	$r_0$	r	m	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE							
$D_1$	2	2	1	2	1	1	0	0	0	1	2	3	4	5	6
$D_2$	6	0	4	0	0	3	2	1	1	1	1	2	3	4	4
				0				2	2	2	4	4	4	5	5
				0				3	3	3	5	5	5	6	6
				0				4	5	6	6	6	6	6	6
$D_3$	6	18	1	0	3	0	0	0	0	0	0	0	0	0	0
				0				0	0	0	0	0	0	0	0
				6				1	2	3	4	5	6	6	6
$D_4$	6	12	2	0	4	1	0	0	0	0	0	0	0	0	0
				6				0	0	0	0	0	0	0	0
				0				1	1	2	4	4	5	6	6
				0				2	3	3	5	6	6	6	6
$D_5$	7	4	4	4	2	2	2	0	0	0	0	1	1	2	2
				0				1	1	2	4	2	3	3	3
				0				2	3	3	5	4	4	5	5
				0				6	5	4	6	5	6	6	6



TABLE 5.2 ( CONTINUED ) .

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE									
D <sub>6</sub>	7	10	3	4	4	1	1	0	0	0	0	0	0	0	0	0	0
				3				1	1	2	4	0	0	0			
				0				2	3	3	5	1	2	3			
								6	5	4	6	4	5	6			
D <sub>7</sub>	9	0	6	0	0	3	4	1	1	1	1	1	1	2	2	2	2
				0				2	2	2	3	3	3	3	3	3	3
				0				4	4	5	4	4	5	4	4	5	5
				0				5	6	6	5	6	6	5	6	6	6
D <sub>8</sub>	9	6	5	0	2	2	3	0	0	0	1	1	1	1	2	2	2
				3				0	0	0	2	2	3	3	3	3	3
				0				1	2	3	4	5	4	5	4	4	4
								4	6	5	5	6	6	6	5	6	6
D <sub>9</sub>	9	18	3	0	6	0	1	0	0	0	0	0	0	0	0	0	0
				9				0	0	0	0	0	0	0	0	0	0
				0				1	1	1	2	2	2	3	3	3	3
								4	5	6	4	5	6	4	5	6	6
D <sub>10</sub>	10	10	5	10	5	2	2	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	2	2	2	3	3
				0				2	2	3	3	4	3	3	4	5	5
								5	6	4	6	5	4	5	6	6	6
								0									
								4									
								5									
								6									
D <sub>11</sub>	12	6	7	6	3	3	4	0	0	0	0	0	0	1	1	1	1
				0				1	1	1	2	2	3	2	2	3	3
				0				2	3	4	3	4	5	4	5	4	4
								5	4	5	6	6	6	6	6	6	6
								1	2	2							
								3	3	3							
								5	4	4							
								6	5	5							
D <sub>12</sub>	12	12	6	6	5	2	3	0	0	0	0	0	0	0	0	0	0
				3				1	1	1	2	2	3	0	0	0	0
				0				2	3	5	3	4	4	1	2	3	3
								4	6	6	5	6	5	4	5	6	6

TABLE 5.2 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
								1	1	2										
								2	3	3										
								5	4	4										
								6	5	6										
D <sub>13</sub>	12	12	6	12	6	3	2	0	0	0	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	1	2	2	2				
				0				2	2	2	3	4	5	3	4	5				
								3	3	4	5	6	6	6	6	5	6			
								0	0	0										
								3	3	4										
								4	4	5										
								6	5	6										
D <sub>14</sub>	12	18	5	6	7	1	2	0	0	0	0	0	0	0	0	0	0	0	0	0
				6				1	1	1	2	2	3	0	0	0				
				0				2	3	5	3	4	4	1	2	3				
								6	5	6	4	6	5	4	5	6				
								0	0	0										
								0	0	0										
								1	2	3										
								4	5	6										
D <sub>15</sub>	15	12	8	12	6	3	4	0	0	0	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	1	2	2	2				
				0				2	2	3	3	4	4	3	3	4				
								6	5	4	6	5	6	4	5	6				
								0	0	0	1	1	2							
								2	3	3	2	3	3							
								5	4	5	4	5	4							
								6	5	6	5	6	6							
D <sub>16</sub>	15	18	7	12	8	2	3	0	0	0	0	0	0	0	0	0	0	0	0	0
				3				1	1	1	1	1	1	2	2	2				
				0				2	2	3	3	4	5	3	3	4				
								4	5	5	6	6	6	4	6	6				
								0	0	0	0	0	0							
								2	3	3	0	0	0							
								5	4	4	1	2	3							
								6	5	5	4	5	6							
D <sub>17</sub>	18	18	9	18	9	3	4	0	0	0	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	1	1	1	1				
				0				2	2	2	3	3	3	4	4	5				
								4	5	6	4	5	6	5	6	6				

**TABLE 5.2 ( CONTINUED ).**

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE										
								0	0	0	0	0	0	0	0	0	0	0
								2	2	2	2	2	2	3	3	3	3	3
								3	3	3	4	4	5	4	4	4	5	5
								4	5	6	5	6	6	5	6	6	6	6

**5.3 MCCGD FOR k = 4, v = 8 :**

**THEOREM 5.3**

For k = 4 and v = 8 the minimal complete class of generator designs is given in Table 5.3.

**TABLE 5.3 MCCGD FOR k = 4, v = 8 .**

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE											
D <sub>1</sub>	2	0	1	0	0	1	0	1	5										
				0				2	6										
				0				3	7										
				0				4	8										
D <sub>2</sub>	8	24	1	0	3	0	0	0	0	0	0	0	0	0	0	0	0	0	
				0				0	0	0	0	0	0	0	0	0	0		
				0				0	0	0	0	0	0	0	0	0	0		
				8				1	2	3	4	5	6	7	8				
D <sub>3</sub>	10	8	4	0	2	2	1	0	0	0	0	1	1	1	2	5			
				4				0	0	0	0	2	2	3	3	6			
				0				1	2	3	4	3	4	4	4	7			
								5	6	7	8	8	7	6	5	8			
								5											
	6																		
	7																		
	8																		

TABLE 5.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE								
D <sub>4</sub>	12	0	6	0	0	2	3	1	1	1	1	1	1	2	2	2
								2	2	3	3	4	4	3	3	4
				0				5	7	5	6	5	6	5	7	5
				0				6	8	7	8	8	7	6	8	7
								-----								
								2	3	3						
								4	4	4						
								6	5	6						
								8	8	7						
D <sub>5</sub>	12	16	4	8	5	1	1	0	0	0	0	0	0	0	0	0
								1	1	1	2	2	3	4	6	0
				4				2	3	4	3	5	5	5	7	1
				0				7	8	6	4	8	6	7	8	5
								-----								
								0	0	0						
								0	0	0						
								2	3	4						
								6	7	8						
D <sub>6</sub>	12	24	3	0	6	1	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0
				12				1	1	1	2	2	3	5	5	5
				0				2	3	4	3	4	4	6	7	8
								-----								
								0	0	0						
								0	0	0						
								6	6	7						
								7	8	8						
D <sub>7</sub>	15	12	6	8	4	2	2	0	0	0	0	0	0	0	0	0
								1	1	2	2	3	3	4	4	0
				2				5	6	5	6	5	6	5	6	1
				0				7	8	8	7	7	8	8	7	2
								-----								
								0	1	1	1	2	2			
								0	2	3	4	3	4			
								3	3	5	7	7	5			
								4	4	6	8	8	6			
D <sub>8</sub>	16	8	7	8	3	2	3	0	0	0	0	0	0	0	0	1
								1	1	1	2	2	3	3	4	3
				0				2	4	5	3	7	4	5	6	6
				0				6	8	7	5	8	7	6	8	7
								-----								
								1	1	1	2	2	2	3		
								2	3	4	4	3	4	4		
								5	7	5	6	6	5	5		
								8	8	6	7	8	7	8		

TABLE 5.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE																
D <sub>9</sub>	16	32	4	0	8	0	1	0	0	0	0	0	0	0	0	0	0							
								0	0	0	0	0	0	0	0	0	0	0						
								1	1	1	1	2	2	2	2	2	3	3						
								5	6	7	8	5	6	7	8	5								
								-----								0	0	0	0	0	0	0		
								0	0	0	0	0	0	0										
								3	3	3	4	4	4	4										
								6	7	8	5	6	7	8										
								-----								0	0	0	0	0	0	0	1	1
								-----								1	2	3	0	0	0	0	2	2
-----								5	6	7	1	2	3	0	3	4								
-----								8	8	8	7	5	6	4	7	8								
-----								1	1	1	1	2	2	2	3	4								
-----								2	3	3	4	3	4	3	4	6								
-----								6	4	5	5	5	5	4	5	7								
-----								7	8	6	6	8	7	6	7	8								
-----								5																
-----								6																
-----								7																
-----								8																
D <sub>11</sub>	19	20	7	16	7	2	2	0	0	0	0	0	0	0	0	0	0							
								1	1	1	1	1	2	2	2	2	2							
								3	4	5	6	7	3	4	5	5								
								5	6	7	8	8	7	8	6	8								
								-----								0	0	0	0	0	0	0	0	0
								2	3	3	3	4	4	4	0	0								
								6	6	6	5	5	5	7	1	3								
								7	8	7	8	6	7	8	2	4								
								-----								1								
								-----								2								
-----								3																
-----								4																
D <sub>12</sub>	19	28	6	8	8	2	1	0	0	0	0	0	0	0	0	0	0							
								1	1	1	1	2	2	3	6	0								
								2	2	3	4	3	5	4	7	1								
								3	4	5	6	4	7	8	8	7								
								-----								0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0								
								1	2	2	3	3	4	4	5	5								
								8	6	8	6	7	5	7	6	8								
								-----								0	0	0	0	0	0	0	0	0
								-----								1	2	3	3	4	4	5	5	
-----								8	6	8	6	7	5	7	6	8								

TABLE 5.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE									
								5 6 7 8									
D <sub>13</sub>	20	8	9	0 4 0	2	3	4	0 0 1 5	0 0 2 6	0 0 3 7	0 0 4 8	1 2 4 6	1 2 5 7	1 2 7 8	1 3 5 8	1 3 6 8	1 3 5 7
								1 3 6 8	1 4 5 6	1 4 7 8	2 3 5 6	2 3 5 8	2 4 5 8	2 4 7 8	2 6 7 8	2 6 7 8	3 4 5 7
D <sub>14</sub>	20	16	8	16 0 0	6	2	3	0 1 2 6	0 1 2 7	0 1 3 8	0 1 4 5	0 1 5 8	0 2 7 8	0 2 8 7	0 2 5 6	0 2 4 5	0 2 3 6
								0 2 6 8	0 3 4 5	0 3 5 6	0 3 6 7	0 4 6 7	0 4 7 8	0 4 8 7	0 5 6 7	0 5 7 8	1 3 4 6
D <sub>15</sub>	23	36	7	11 11 1	10	1	2	0 1 2 5	0 1 3 6	0 1 4 7	0 1 5 8	0 2 6 7	0 2 7 8	0 3 8 7	0 3 6 8	0 4 7 8	0 4 5 6
								0 3 5 7	0 4 6 8	0 4 7 8	0 5 6 7	0 6 7 8	0 6 8 7	0 7 8 6	0 7 8 6	0 8 6 7	0 8 5 6
								0 4 5	0 4 6	0 4 7	0 5 6	0 6 7	0 7 8	0 7 8	0 8 6	0 8 5	0 5 6

TABLE 5.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE									
D <sub>16</sub>	23	44	6	3 19 1	11	1	1	0	0	0	0	0	0	0	0	0	0
								1	1	6	0	0	0	0	0	0	0
								2	4	7	1	1	1	2	2	2	2
								3	5	8	6	7	8	4	5	6	6
								-----									
								0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0
								2	2	3	3	3	3	3	3	4	4
								7	8	4	5	6	7	8	6	7	7
								-----									
								0	0	0	0	0					
								0	0	0	0	0					
4	5	5	5	0													
8	6	7	8	1													
D <sub>17</sub>	24	24	9	24 0 0	9	2	3	0	0	0	0	0	0	0	0	0	0
								1	1	1	1	1	1	1	1	1	1
								2	2	3	3	4	4	5	5	7	7
								7	8	5	6	6	8	6	7	8	8
								-----									
								0	0	0	0	0	0	0	0	0	0
								2	2	2	2	2	2	2	2	3	3
								3	3	4	4	5	5	6	6	4	4
								7	8	5	6	6	8	7	7	5	5
								-----									
								0	0	0	0	0	0				
								3	3	3	4	4	4				
5	6	6	5	6	7												
8	8	7	7	8	8												
D <sub>18</sub>	30	40	10	22 6 2	12	2	3	0	0	0	0	0	0	0	0	0	0
								1	1	1	1	1	1	1	1	1	1
								2	2	3	3	4	4	5	6	7	7
								5	8	5	6	8	6	7	7	8	8
								-----									
								0	0	0	0	0	0	0	0	0	0
								2	2	2	2	2	2	2	2	3	3
								3	3	4	4	5	6	7	7	4	4
								7	8	5	6	6	7	8	5	7	7
								-----									
								0	0	0	0	0	0	0	0	0	0
								3	3	4	4	0	0	0	0	0	0
5	6	5	6	3	3	4	4	5	5								
8	8	7	8	6	7	8	7	6	6								
-----																	
0	0	0															
0	0	0															
5	0	0															
8	1	2															

TABLE 5.3 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE								
D <sub>19</sub>	32	32	12	32	12	4	3	0	0	0	0	0	0	0	0	0
								1	1	1	1	1	1	1	1	1
								2	2	2	2	3	3	3	3	4
								5	6	7	8	5	6	7	8	5
								0	0	0	0	0	0	0	0	0
								1	1	1	2	2	2	2	2	2
								4	4	4	3	3	3	3	4	4
								6	7	8	5	6	7	8	5	6
								0	0	0	0	0	0	0	0	0
								2	2	3	3	3	3	5	5	5
								4	4	4	4	4	4	6	6	6
								7	8	5	6	7	8	7	7	8
								0	0	0	0	0				
								5	5	5	6	6				
								6	7	7	7	7				
								8	8	8	8	8				

5.4 MCCGD FOR  $k = 4, v = 10$  :

THEOREM 5.4

For  $k = 4$  and  $v = 10$  the minimal complete class of generator designs is given in Table 5.4

TABLE 5.4 MCCGD FOR  $k = 4, v = 10$ .

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE								
D <sub>1</sub>	10	30	1	0	3	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0
								1	2	3	4	5	6	7	8	9
								0								
								0								
								0								
								10								



TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE											
D <sub>2</sub>	13	12	4	6	3	1	1	0	0	0	0	0	0	0	0	0	0	0	0
								1	2	3	4	5	6	0	0	0			
								9	7	8	7	8	7	1	3	5			
								10	8	10	9	9	10	2	4	6			
								1	1	2	2								
								3	4	3	4								
								5	6	6	5								
7	8	9	10																
D <sub>3</sub>	14	26	3	4	5	1	0	0	0	0	0	0	0	0	0	0	0	0	0
								1	1	6	6	0	0	0	0	0			
								2	4	7	9	2	2	3	3	7			
								3	5	8	10	4	5	4	5	9			
								0	0	0	0	0							
								0	0	0	0	0							
								7	8	8	0	0							
10	9	10	1	6															
D <sub>4</sub>	15	0	6	0	0	2	2	1	1	1	1	1	1	2	2	2			
								2	2	3	4	4	5	3	4	6			
								3	5	6	7	6	7	4	5	8			
								8	9	9	8	10	10	7	6	10			
								2	3	3	3	4	6						
								7	5	4	5	5	7						
								9	6	9	8	8	8						
10	7	10	10	9	9														
D <sub>5</sub>	15	10	5	0	2	2	1	0	0	0	0	0	1	1	1	1			
								0	0	0	0	0	2	3	4	6			
								1	3	5	7	9	3	5	5	8			
								2	4	6	8	10	4	9	7	10			
								2	2	2	3	4	7						
								3	4	6	6	6	8						
								5	5	7	7	8	9						
8	10	9	10	9	10														
D <sub>6</sub>	19	26	5	8	6	1	1	0	0	0	0	0	0	0	0	0	0	0	
								1	1	2	2	3	3	4	4	0			
								7	8	5	8	5	6	6	7	1			
								10	9	9	10	10	9	10	9	5			
								0	0	0	0	0	0	0	0	1			
								0	0	0	0	0	0	0	0	2			
								1	2	2	3	3	4	4	9	3			
6	6	7	7	8	8	5	10	4											

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE																																																																																																																														
								5 6 7 8																																																																																																																														
D <sub>7</sub>	20	40	4	0 20 0	8	1	0	<table border="0"> <tr><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>1</td><td>1</td><td>2</td><td>2</td><td>2</td><td>3</td><td>3</td></tr> <tr><td>2</td><td>3</td><td>4</td><td>5</td><td>3</td><td>4</td><td>5</td><td>4</td><td>5</td></tr> <tr><td colspan="9">-----</td></tr> <tr><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td><td>0</td></tr> <tr><td>4</td><td>6</td><td>6</td><td>6</td><td>6</td><td>7</td><td>7</td><td>7</td><td>8</td></tr> <tr><td>5</td><td>7</td><td>8</td><td>9</td><td>10</td><td>8</td><td>9</td><td>10</td><td>9</td></tr> <tr><td colspan="9">-----</td></tr> <tr><td>0</td><td>0</td><td colspan="7"></td></tr> <tr><td>0</td><td>0</td><td colspan="7"></td></tr> <tr><td>8</td><td>9</td><td colspan="7"></td></tr> <tr><td>10</td><td>10</td><td colspan="7"></td></tr> </table>	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	2	2	2	3	3	2	3	4	5	3	4	5	4	5	-----									0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	4	6	6	6	6	7	7	7	8	5	7	8	9	10	8	9	10	9	-----									0	0								0	0								8	9								10	10							
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TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
D <sub>10</sub>	25	50	5	0	10	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0
				25				0	0	0	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	2	2	2	2	2	2	2	2
								6	7	8	9	10	6	7	8	9				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	3	3	3	3	3	3	4	4	4	4	4	4
								10	6	7	8	9	10	6	7	8				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0
								4	4	5	5	5	5	5	5	5	5	5	5	5
								9	10	6	7	8	9	10						
D <sub>11</sub>	25	40	6	10	9	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0
				15				1	1	1	2	2	3	3	5	6				
				0				2	4	7	4	5	4	5	6	8				
								3	9	8	6	9	8	10	7	10				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								7	0	0	0	0	0	0	0	0	0	0	0	0
								9	1	1	1	2	2	2	3	3				
								10	5	6	10	7	8	10	6	7				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0
								3	4	4	4	5	6	8						
								9	5	7	10	8	8	9						
D <sub>12</sub>	25	20	8	20	6	2	2	0	0	0	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	1	2	2	2				
				0				3	3	4	4	5	9	3	3	4				
				0				8	9	7	10	6	10	8	10	7				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	2	3	3	4	4	5	5				
								4	5	6	5	6	5	6	6	8				
								9	10	9	7	7	8	8	10	9				
								0	0	1	1	3	3	7						
								6	7	2	2	4	4	8						
								7	8	5	6	5	6	9						
								9	10	7	8	9	10	10						
D <sub>13</sub>	25	30	7	20	8	2	1	0	0	0	0	0	0	0	0	0	0	0	0	0
				5				1	1	1	1	1	1	2	2	2				
				0				2	2	3	4	5	9	3	4	6				
								3	4	5	7	8	10	10	8	7				

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE													
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								2	3	3	3	4	5	6	6	6	6	6	6	6	
								5	4	4	7	5	6	8	8	8	9	9	9	9	
								9	5	6	9	10	7	9	10	10					
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								7	7	0	0	0	0	0	0	0	0	0	0	0	
								8	8	1	2	3	4	7							
								9	10	6	5	8	9	10							
D <sub>14</sub>	26	14	9	12	4	2	3	0	0	0	0	0	0	0	0	0	0	0	0	0	
				1				1	1	1	1	2	2	2	2	2	2	2	2	3	
				0				3	4	5	7	3	4	5	7	4	4	4	4	4	
								6	8	10	9	10	6	9	8	7					
								0	0	0	0	1	1	1	1	1	1	1	1	1	
								3	4	5	0	2	2	3	4	5	5	5	5	5	
								5	6	6	9	6	8	6	9	7					
								8	8	7	10	7	10	9	10	8					
								2	2	2	3	3	3	4	4						
								3	4	5	4	5	6	5	5						
								7	8	6	7	8	8	6	7						
								9	9	10	10	9	10	9	10						
D <sub>15</sub>	27	18	9	14	5	3	2	0	0	0	0	0	0	0	0	0	0	0	0	0	
				2				1	1	1	1	1	2	2	2	2	2	2	2	3	
				0				2	2	3	4	6	3	5	6	4	4	4	4	4	
								3	4	5	7	10	4	7	10	9					
								0	0	0	0	0	0	0	0	1	1	1	1	1	
								3	4	6	6	7	0	0	2	3	3	3	3	3	
								6	5	7	8	8	5	8	5	7					
								7	8	9	9	10	10	9	8	8					
								1	1	2	2	2	3	3	4	7					
								4	5	3	4	7	4	6	5	8					
								9	6	5	6	9	5	8	6	9					
								10	9	9	8	10	10	10	10	7	10				
D <sub>16</sub>	30	40	8	20	10	1	2	0	0	0	0	0	0	0	0	0	0	0	0	0	
				10				1	1	1	1	1	1	2	2	2	2	2	2	2	2
				0				2	3	4	5	6	8	3	4	5	5	5	5	5	5
								7	9	9	10	8	10	6	10	8					
								0	0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	3	3	3	3	4	4	4	4	4	4	4	4
								6	9	4	5	6	7	5	6	7					
								9	10	8	7	10	10	6	7	8					

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
								0	0	0	0	0	0	0	0	0	0	0	0	0
								5	5	0	0	0	0	0	0	0	0	0	0	0
								7	8	1	2	3	4	5	1	2				
								9	9	6	7	8	9	10	7	8				
								0	0	0										
								0	0	0										
								3	4	5										
								9	10	6										
D <sub>17</sub>	30	30	9	30	9	2	2	0	0	0	0	0	0	0	0	0	0	0	0	0
				0				1	1	1	1	1	1	1	1	1	1	1	1	1
				0				2	2	3	3	4	4	5	5	6	6	6	6	6
				0				6	7	7	8	8	9	10	10	9				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	2	2	2	2	2	3	3				
								3	3	4	4	5	5	6	4	4				
								7	10	9	10	8	8	9	6	8				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								3	3	3	4	4	4	5	6	6	6	6	6	6
								5	5	6	5	5	7	6	7	8	8	8	8	8
								9	9	10	6	7	10	7	8	10				
								0	0	0										
								7	7	8										
								8	9	9										
								9	10	10										
D <sub>18</sub>	31	24	10	22	7	2	3	0	0	0	0	0	0	0	0	0	0	0	0	0
				1				1	1	1	1	1	1	1	2	2				
				0				2	3	3	4	5	6	6	3	4				
								7	8	10	8	7	9	9	9	8				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	2	2	3	3	3	3	4				
								4	5	6	6	4	5	7	7	5				
								10	10	7	8	6	9	8	10	9				
								0	0	0	0	0	1	1	1	2				
								4	4	5	8	0	2	4	5	3				
								7	9	7	9	5	6	7	8	7				
								8	10	10	10	6	10	9	10	9				
								2	3	3	4									
								5	4	5	5									
								8	6	6	6									
								9	10	8	7									
D <sub>19</sub>	32	58	7	7	12	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0
				24				1	1	1	2	3	4	5	0	0				
				1				2	4	6	6	7	8	8	1	1				
								3	5	7	9	10	10	9	8	9				

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
								0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0
								1	2	2	2	2	2	3	3	3	3	3	3	3
								10	4	5	7	8	10	4	5	6				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0
								3	3	4	4	4	5	5	5	6				
								8	9	6	7	9	6	7	10	8				
								0	0	0	0	0								
								0	0	0	0	0								
								6	7	7	9	0								
								10	8	9	10	1								
D <sub>20</sub>	32	48	8	17	11	2	1	0	0	0	0	0	0	0	0	0	0	0	0	0
				14				1	1	1	1	1	2	2	2	3				
				1				2	2	3	4	5	3	5	5	4				
								3	4	6	8	7	4	6	9	5				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								3	4	6	6	6	7	7	7	0				
								5	6	8	8	9	8	8	9	1				
								10	7	9	10	10	9	10	10	5				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0
								1	1	2	2	2	3	3	3	4				
								9	10	7	8	10	7	8	9	5				
								0	0	0	0	0								
								0	0	0	0	0								
								4	4	5	6	0								
								9	10	8	7	6								
D <sub>21</sub>	32	28	10	24	8	3	2	0	0	0	0	0	0	0	0	0	0	0	0	0
				2				1	1	1	1	1	1	1	1	2				
				0				2	2	2	3	3	4	5	6	3				
								3	4	9	5	10	7	9	8	4				
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	2	2	3	3	3	4	4				
								3	4	5	6	4	5	6	5	6				
								7	5	6	10	5	8	8	10	9				
								0	0	0	0	0	0	0	0	1				
								4	5	6	6	7	8	0	0	4				
								7	6	7	9	8	9	7	8	6				
								8	7	10	10	9	10	9	10	8				

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE													
								1	2	2	3	3									
								5	5	7	4	6									
								7	8	8	9	7									
								10	9	10	10	9									
D <sub>22</sub>	34	16	12	8	4	3	4	0	0	0	0	0	0	0	0	0	0	0	0	0	0
				4				1	1	1	1	2	4	5	8	0					
				0				2	3	4	5	3	7	6	9	2					
								7	8	10	6	10	9	10	10	6					
								0	0	1	1	1	1	1	1	1					
								0	0	2	2	3	3	4	4	5					
								4	5	6	7	6	7	6	8	7					
								8	9	8	10	10	9	9	9	8					
								1	2	2	2	2	2	2	2	3					
								5	3	3	4	4	4	5	5	4					
								9	5	8	6	7	9	6	7	5					
								10	9	9	8	10	10	9	8	8					
								3	3	3	3	4	5								
								4	4	6	6	5	7								
								5	6	7	8	6	8								
								10	7	9	10	7	10								
D <sub>23</sub>	36	34	11	32	10	2	3	0	0	0	0	0	0	0	0	0	0	0	0	0	0
				1				1	1	1	1	1	1	1	1	1					
				0				2	3	3	4	4	5	5	6	7					
								6	7	8	8	10	6	9	10	9					
								0	0	0	0	0	0	0	0	0					
								1	2	2	2	2	2	2	2	2					
								8	3	3	4	4	5	5	6	7					
								9	6	9	7	8	9	10	9	8					
								0	0	0	0	0	0	0	0	0					
								2	3	3	3	3	3	3	4	4					
								8	4	4	5	5	6	7	5	6					
								10	9	10	7	8	10	10	6	7					
								0	0	0	0	0	0	1	3	4					
								4	4	5	5	5	0	2	6	5					
								6	7	6	7	8	9	7	8	9					
								8	9	7	8	10	10	10	9	10					
D <sub>24</sub>	37	58	9	14	13	1	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0
				22				1	1	1	1	1	2	2	2	2					
				0				2	3	4	7	8	3	4	6	8					
								5	6	10	9	10	7	10	10	9					

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE													
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								3	3	3	4	4	0	0	0	0	0	0	0	0	
								4	5	9	5	6	1	1	1	1	1	1	1	1	
								8	9	10	7	9	6	7	8	9					
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								2	2	2	2	3	3	3	3	3	3	3	4	4	
								6	7	8	9	6	7	8	10	6					
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								0	0	0	0	0	0	0	0	0	0	0	0	0	
								4	4	4	5	5	5	5	5	5	5	7	7	7	
								7	8	9	6	8	9	10	10	10	10	10	10	10	
								5													
								6													
								7													
								8													
D <sub>25</sub>	37	48	10	27	12	2	2	0	0	0	0	0	0	0	0	0	0	0	0	0	
				9				1	1	1	1	1	1	1	1	1	1	1	1	1	1
				1				2	2	3	3	4	4	5	5	6	6	6	6	6	6
								9	9	10	10	8	8	6	7	7	7	7	7	7	7
								0	0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	2	2	2	2	3	3	3	3	3	3	3	3
								3	3	4	4	7	7	6	6	6	6	6	6	6	6
								5	5	6	6	10	10	8	8	7	7	7	7	7	7
								0	0	0	0	0	0	0	0	0	0	0	0	0	0
								3	4	4	5	5	6	6	7	7	7	7	7	7	7
								4	5	5	8	8	9	9	8	8	8	8	8	8	8
								7	9	9	10	10	10	10	9	9	9	9	9	9	9
								0	0	0	0	0	0	0	0	0	0	0	0	0	0
								0	0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	3	3	4	4	5	5	6	6	6	6	6	6
								8	8	9	9	10	10	6	7	7	7	7	7	7	7
								0													
								0													
								0													
								1													
D <sub>26</sub>	37	38	11	34	11	3	2	0	0	0	0	0	0	0	0	0	0	0	0	0	
				2				1	1	1	1	1	1	1	1	1	1	1	1	1	1
				0				2	2	2	3	3	4	4	5	5	5	5	5	5	5
								3	4	5	7	9	7	10	8	10	8	10	8	10	10



TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE														
								0	0	0	0	0	0	0	0	0	0	0	0	0		
								1	1	2	2	2	2	2	2	2	2	2	2	2		
								6	6	3	3	4	5	6	6	6	6	6	8	8		
								8	9	4	5	7	10	7	8	9						
								0	0	0	0	0	0	0	0	0	0	0	0	0		
								2	3	3	3	3	3	3	3	4	4	4	4	4		
								9	4	4	5	6	6	8	5	5						
								8	5	8	7	10	10	9	8	9						
								0	0	0	0	0	0	0	0	0	0	0	0	0		
								4	4	5	5	6	7	7	0	0						
								6	6	6	6	7	8	9	7	8						
								9	10	7	9	8	10	10	9	10						
								7														
								8														
								9														
								10														
D <sub>27</sub>	42	48	12	34	13	2	3	0	0	0	0	0	0	0	0	0	0	0	0	0		
				7				1	1	1	1	1	1	1	1	1	1	1	1	1	1	
				0				2	2	3	3	4	4	5	5	7	7	7	7	7	7	
								6	10	6	8	7	9	8	9	9	9	9	9	9	9	
								0	0	0	0	0	0	0	0	0	0	0	0	0	0	
								1	1	2	2	2	2	2	2	2	2	2	2	2	2	
								7	8	3	3	4	4	5	5	6	6	6	6	6	6	
								10	10	6	9	8	10	7	10	8						
								0	0	0	0	0	0	0	0	0	0	0	0	0	0	
								2	2	3	3	3	3	3	3	3	3	3	3	3	3	
								7	8	4	4	5	5	7	8	9	9	9	9	9	9	
								9	9	6	7	9	10	8	10	10						
								0	0	0	0	0	0	0	0	0	0	0	0	0	0	
								4	4	4	4	4	5	6	0	0						
								5	5	6	7	8	6	9	1	2						
								6	10	9	10	9	7	10	6	7						
								0	0	0	0	0	5									
								0	0	0	0	0	6									
								3	4	5	5	6	7									
								7	8	8	9	10	8									
D <sub>28</sub>	54	56	16	48	16	3	4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
				4				1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
				0				2	2	2	3	3	3	4	4	4	4	4	4	4	4	4
								6	8	9	6	9	10	7	9	10						

TABLE 5.4 ( CONTINUED ).

DESIGN	b	r <sub>0</sub>	r	m	λ <sub>0</sub>	λ <sub>1</sub>	λ <sub>2</sub>	STRUCTURE												
								0	0	0	0	0	0	0	0	0	0	0	0	0
								1	1	1	1	1	1	1	1	1	2	2	2	2
								5	5	5	6	6	7	8	3	3	3	3	3	3
								7	8	10	7	8	9	10	7	9	7	9	9	9
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	2	2	2	2	2	2	2	2	2	2	2
								3	4	4	4	5	5	5	6	7	7	7	7	7
								10	6	7	8	6	8	9	10	9	10	9	10	9
								0	0	0	0	0	0	0	0	0	0	0	0	0
								2	2	3	3	3	3	3	3	3	3	3	3	3
								7	8	4	4	4	5	5	5	6	6	6	6	6
								10	10	6	8	10	7	9	10	7	9	10	7	7
								0	0	0	0	0	0	0	0	0	0	0	0	0
								3	3	3	4	4	4	4	4	4	4	4	4	4
								6	7	8	5	5	5	6	6	6	7	7	7	7
								8	8	9	7	9	10	9	10	8	8	8	8	8
								0	0	0	0	0	0	0	0	5	7	7	7	7
								4	5	5	0	0	0	0	0	6	8	8	8	8
								8	6	6	5	6	7	9	9	9	9	9	9	9
								9	7	8	8	9	10	10	10	10	10	10	10	10

## CHAPTER 6.

### 6.1 A-OPTIMAL DESIGNS :

Consider an E GD TIB design with parameters  $(k, v, b, \lambda_0, \lambda_1, \lambda_2)$ . Then from Section 2.1.3, the trace of the inverse of information matrix  $M$  is given as

$$\text{Tr}(M^{-1}) = \frac{vk \left[ (\lambda_0 + \lambda_1 + \frac{v}{2} \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_2) - \frac{v}{2} (\frac{v}{2} - 1) \lambda_2^2 \right]}{\lambda_0 (\lambda_0 + v \lambda_2) (\lambda_0 + \frac{v}{2} \lambda_1 + \frac{v}{2} \lambda_2)}$$

The design, among all the E GD MB TIB designs for given  $(k, v, b)$ , which minimizes the  $\text{Tr}(M^{-1})$ , given above, is called A-optimal. In fact, such a design, if exists, has a least average variance of the BLUE's of treatment contrasts  $\{\alpha_0 - \alpha_i\}, i = 1, 2, \dots, v$  (the sum of diagonal elements of the variance-covariance matrix of the estimates).

Majumdar (1981) and Majumdar and Notz (1983) showed that if there is a design that minimizes  $\text{Tr}(M^{-1})$  then it is binary in the test-treatments and the number of units allocated to the control are less than half of the the total units in the design, i.e.

$$r_0 < \frac{b k}{2}$$

The actual minimization of the  $\text{Tr}(M^{-1})$ , given above, is done for  $k = 2, 3, 4$  and  $v = 4, 6, 8, 10$ , and for each  $b < 50$ . The A-optimal designs are obtained and given in the following tables with their possible structure through the corresponding ( for that  $k$  and  $v$  ) MCCGD given in the previous chapter.

**6.2 A-OPTIMAL E MB GD TIB DESIGNS FOR  $k = 2$  :**

**TABLE 6.2 A-OPTIMAL E MB GD TIB DESIGNS FOR  $k = 2$ .**

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
2	4	4	1	0	0	$D_3$ (Table 3.1)
		6	1	1	0	$D_1^U D_3$
		8	1	0	1	$D_2^U D_3$
		10	2	1	0	$D_1^U 2D_3$
		12	2	0	1	$D_2^U 2D_3$
		14	2	1	1	$D_1^U D_2^U 2D_3$
		16	3	0	1	$D_1^U 3D_3$
		18	3	1	1	$D_1^U D_2^U 3D_3$
		20	3	2	1	$2D_1^U D_2^U 3D_3$
		22	4	1	1	$D_1^U D_2^U 4D_3$
		24	4	2	1	$2D_1^U D_2^U 4D_3$
		26	4	1	2	$D_1^U 2D_2^U 4D_3$
		28	5	2	1	$2D_1^U D_2^U 5D_3$
		30	5	1	2	$D_1^U 2D_2^U 5D_3$

TABLE 6.2 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
2	4	32	5	2	2	$2D_1^U \ 2D_2^U \ 5D_3$
		34	6	1	2	$D_1^U \ 2D_2^U \ 6D_3$
		36	6	2	2	$2D_1^U \ 2D_2^U \ 6D_3$
		38	6	3	2	$3D_1^U \ 2D_2^U \ 6D_3$
		40	7	2	2	$2D_1^U \ 2D_2^U \ 7D_3$
		42	7	3	2	$3D_1^U \ 2D_2^U \ 7D_3$
		44	8	2	2	$2D_1^U \ 2D_2^U \ 8D_3$
		46	8	3	2	$3D_1^U \ 2D_2^U \ 8D_3$
		48	8	2	3	$2D_1^U \ 3D_2^U \ 8D_3$
		50	8	3	3	$3D_1^U \ 3D_2^U \ 8D_3$
2	6	6	1	0	0	$D_1$ (Table 3.2)
		9	0	0	1	$D_3$
		12	1	1	0	$D_1^U \ D_2$
		15	1	0	1	$D_2^U \ D_3$
		18	2	1	0	$D_1^U \ 2D_2$
		21	2	0	1	$2D_2^U \ D_3$
		24	3	1	0	$D_1^U \ 3D_2$
		27	3	0	1	$3D_2^U \ D_3$
		30	4	1	0	$D_1^U \ 4D_2$
		33	3	1	1	$D_1^U \ 3D_2^U \ D_3$
		36	5	1	0	$D_1^U \ 5D_2$
		39	4	1	1	$D_1^U \ 4D_2^U \ D_3$
		42	3	1	2	$D_1^U \ 3D_2^U \ 2D_3$

TABLE 6.2 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
2	6	45	5	1	1	$D_1 \cup 5D_2 \cup D_3$
		48	4	1	2	$D_1 \cup 4D_2 \cup 2D_3$
2	8	8	1	0	0	$D_1$ (Table 3.3)
		12	0	1	0	$D_2$
		16	2	0	0	$2D_1$
		20	1	1	0	$D_1 \cup D_2$
		24	1	0	1	$D_1 \cup D_3$
		28	2	1	0	$2D_1 \cup D_2$
		32	2	0	1	$2D_1 \cup D_3$
		36	3	1	0	$3D_1 \cup D_2$
		40	3	0	1	$3D_1 \cup D_3$
		44	2	1	1	$2D_1 \cup D_2 \cup D_3$
		48	4	0	1	$4D_1 \cup D_3$
2	10	10	1	0	0	$D_1$ (Table 3.4 )
		20	2	0	0	$2D_1$
		25	0	0	1	$D_3$
		30	1	1	0	$D_1 \cup D_2$
		35	1	0	1	$D_1 \cup D_3$
		40	2	1	0	$2D_1 \cup D_2$
		45	2	0	1	$2D_1 \cup D_3$
		50	3	1	0	$3D_1 \cup D_2$

**6.3 A-OPTIMAL E MB GD TIB DESIGNS FOR  $k = 3$  :**

**TABLE 6.3 A-OPTIMAL E MB GD TIB DESIGNS FOR  $k = 3$ .**

$k$	$v$	$b$	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
3	4	2	1	1	0	$D_1$ (Table 4.2)
		4	2	0	1	$D_4$
		6	3	1	1	$D_1 \cup D_4$
		7	2	2	2	$D_6$
		8	4	2	1	$2D_1 \cup D_4$
		9	3	3	2	$D_1 \cup D_6$
		10	5	1	2	$D_1 \cup 2D_4$
		11	4	2	3	$D_4 \cup D_6$
		12	6	2	2	$2D_1 \cup 2D_4$
		13	5	3	3	$D_1 \cup D_4 \cup D_6$
		14	7	3	2	$3D_1 \cup 2D_4$
		15	6	4	3	$2D_1 \cup D_4 \cup D_6$

TABLE 6.3 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
3	4	16	8	2	3	$2D_1 \cup 3D_4$
		17	7	3	4	$D_1 \cup 2D_4 \cup D_6$
		18	9	3	3	$3D_1 \cup 3D_4$
		19	8	4	4	$2D_1 \cup 2D_4 \cup D_6$
		20	10	4	3	$4D_1 \cup 3D_4$
		21	9	5	4	$3D_1 \cup 2D_4 \cup D_6$
		22	11	3	4	$3D_1 \cup 4D_4$
		23	10	4	5	$2D_1 \cup 3D_4 \cup D_6$
		24	12	4	4	$4D_1 \cup 4D_4$
		25	11	5	5	$3D_1 \cup 3D_4 \cup D_6$
		26	13	5	4	$5D_1 \cup 4D_4$
		27	12	6	5	$4D_1 \cup 3D_4 \cup D_6$
		28	14	4	5	$4D_1 \cup 5D_4$
		29	13	5	6	$3D_1 \cup 4D_4 \cup D_6$
		30	15	5	5	$5D_1 \cup 5D_4$
		31	14	6	6	$4D_1 \cup 4D_4 \cup D_6$
		32	16	6	5	$6D_1 \cup 5D_4$
		33	15	7	6	$5D_1 \cup 4D_4 \cup D_6$
		34	17	5	6	$5D_1 \cup 6D_4$
		35	16	6	7	$4D_1 \cup 5D_4 \cup D_6$
		36	18	6	6	$6D_1 \cup 6D_4$
		37	17	7	7	$5D_1 \cup 5D_4 \cup D_6$



TABLE 6.3 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
3	4	38	19	7	6	$7D_1 \cup 6D_4$
		39	18	8	7	$6D_1 \cup 5D_4 \cup D_6$
		40	20	6	7	$6D_1 \cup 7D_4$
		41	19	7	8	$5D_1 \cup 6D_4 \cup D_6$
		42	21	7	7	$7D_1 \cup 7D_4$
		43	20	8	8	$6D_1 \cup 6D_4 \cup D_6$
		44	22	8	7	$8D_1 \cup 7D_4$
		45	21	9	8	$7D_1 \cup 6D_4 \cup D_6$
		46	23	7	8	$7D_1 \cup 8D_4$
		47	22	8	9	$6D_1 \cup 7D_4 \cup D_6$
		48	24	8	8	$8D_1 \cup 8D_4$
		49	23	9	9	$7D_1 \cup 7D_4 \cup D_6$
		50	25	9	8	$9D_1 \cup 8D_4$
3	6	2	0	1	0	$D_1$ (Table 4.3)
		6	2	1	0	$D_3$
		7	1	1	1	$D_4$
		9	3	0	1	$D_5$
		10	0	2	2	$D_6$
		11	3	1	1	$D_1 \cup D_5$
		12	2	1	2	$D_7$
		13	3	2	1	$2D_1 \cup D_5$
		14	2	2	2	$2D_4$

TABLE 6.3 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
3	6	15	5	1	1	$D_3^U D_5$
		16	4	1	2	$D_4^U D_5$
		17	5	2	1	$D_1^U D_3^U D_5$
		18	4	2	2	$D_3^U D_7$
		19	3	2	3	$D_4^U D_7$
		20	6	1	2	$D_1^U 2D_5$
		21	7	2	1	$2D_3^U D_5$
		22	6	2	2	$2D_1^U 2D_5$
		23	5	2	3	$2D_4^U D_5$
		24	8	1	2	$D_3^U 2D_5$
		25	5	3	3	$D_1^U 2D_4^U D_5$
		26	8	2	2	$D_1^U D_3^U 2D_5$
		27	7	2	3	$D_1^U D_4^U 2D_5$
		28	8	3	2	$2D_1^U D_3^U 2D_5$
		29	7	3	3	$2D_1^U D_4^U 2D_5$
		30	10	2	2	$2D_3^U 2D_5$
		31	9	2	3	$2D_1^U 3D_5$
		32	10	3	2	$D_1^U 2D_3^U 2D_5$
		33	9	3	3	$3D_1^U 3D_5$
		34	8	3	4	$D_1^U 2D_4^U 2D_5$
		35	11	2	3	$D_1^U D_3^U 3D_5$
		36	12	3	2	$3D_3^U 2D_5$
		37	11	3	3	$2D_1^U D_3^U 3D_5$

TABLE 6.3 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
3	6	38	10	3	4	$2D_1^U D_4^U 3D_5$
		39	11	4	3	$3D_1^U D_3^U 3D_5$
		40	10	4	4	$3D_1^U D_4^U 3D_5$
		41	13	3	3	$D_1^U 2D_3^U 3D_5$
		42	12	3	4	$3D_1^U 4D_5$
		43	13	4	3	$2D_1^U 2D_3^U 3D_5$
		44	12	4	4	$4D_1^U 4D_5$
		45	15	3	3	$3D_3^U 3D_5$
		46	14	3	4	$2D_1^U D_3^U 4D_5$
		47	15	4	3	$D_1^U 3D_3^U 3D_5$
		48	14	4	4	$3D_1^U D_3^U 4D_5$
		49	13	4	5	$3D_1^U D_4^U 4D_5$
		50	14	5	4	$4D_1^U D_3^U 4D_5$
3	8	8	2	0	0	$D_1$ (Table 4.4)
		12	3	1	0	$D_3$
		16	4	0	1	$D_4$
		18	3	1	1	$D_5$
		19	2	2	1	$D_6$
		20	5	1	0	$D_1^U D_3$
		23	5	1	1	$D_8$
		24	4	2	1	$D_2^U D_3$
		25	2	2	2	$D_{10}$
		28	5	1	2	$D_2^U D_4$

**TABLE 6.3 ( CONTINUED ).**

<b>k</b>	<b>v</b>	<b>b</b>	$\lambda_0$	$\lambda_1$	$\lambda_2$	<b>STRUCTURE</b>
3	8	30	6	2	1	$D_3^U D_5$
		31	3	3	2	$D_2^U D_6$
		32	8	0	1	$2D_1^U D_4$
		34	7	1	2	$D_4^U D_5$
		35	6	2	2	$D_2^U D_8$
		36	5	3	2	$2D_2^U D_3$
		37	3	3	3	$D_2^U D_{10}$
		38	2	4	3	$D_{13}$
		39	9	1	2	$D_4^U D_8$
		40	8	2	2	$D_2^U D_3^U D_4$
		42	7	3	2	$D_6^U D_8$
		43	4	4	3	$2D_2^U D_6$
		44	11	1	2	$D_3^U 2D_4$
		46	10	2	2	$2D_8$
		47	9	3	2	$2D_3^U D_7$
		48	6	4	3	$3D_2^U D_3$
		49	4	4	4	$2D_2^U D_{10}$
		50	3	5	4	$D_2^U D_{13}$
3	10	10	2	0	0	$D_1$ (Table 4.5)
		14	2	1	0	$D_2$
		20	4	1	0	$D_3$
		25	3	1	1	$D_5$

TABLE 6.3 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
3	10	30	2	1	2	$D_8$
		32	5	1	1	$D_{10}$
		35	7	0	1	$D_1 \cup D_6$
		37	4	1	2	$D_{12}$
		39	5	2	1	$D_2 \cup D_5$
		40	8	1	0	$2D_1 \cup D_3$
		44	6	1	2	$D_{13}$
		45	7	2	1	$D_3 \cup D_5$
		49	3	2	3	$D_{14}$
		50	8	1	2	$D_5 \cup D_6$

6.4 A-OPTIMAL E MB GD TIB DESIGNS FOR  $k = 4$ .

TABLE 6.4 A-OPTIMAL E MB GD TIB DESIGNS FOR  $k = 4$ .

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	4	1	0	1	1	$D_1$ (Table 5.1)
		2	2	1	0	$D_2$
		3	2	2	1	$D_1^{\cup} D_2$
		4	3	2	2	$D_4$
		5	3	3	3	$D_1^{\cup} D_4$
		6	5	3	2	$D_2^{\cup} D_4$
		7	5	4	3	$D_1^{\cup} D_2^{\cup} D_4$
		8	6	4	4	$2D_4$
		9	7	3	4	$D_1^{\cup} D_3^{\cup} D_4$
		10	8	5	4	$D_2^{\cup} 2D_4$
		11	8	6	5	$D_1^{\cup} D_2^{\cup} 2D_4$
		12	9	6	6	$3D_4$
		13	10	5	6	$D_1^{\cup} D_3^{\cup} 2D_4$
		14	11	7	6	$D_2^{\cup} 3D_4$
		15	11	8	7	$D_1^{\cup} D_2^{\cup} 3D_4$
		16	12	8	8	$4D_4$
		17	13	7	8	$D_1^{\cup} D_3^{\cup} 3D_4$
		18	14	9	8	$D_2^{\cup} 4D_4$
		19	14	10	9	$D_1^{\cup} D_2^{\cup} 4D_4$
		20	15	10	10	$5D_4$
		21	16	9	10	$D_1^{\cup} D_3^{\cup} 4D_4$

TABLE 6.4 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	4	22	17	11	10	$D_2^U 5D_4$
		23	17	12	11	$D_1^U D_2^U 5D_4$
		24	18	12	12	$6D_4$
		25	19	11	12	$D_1^U D_3^U 5D_4$
		26	20	13	12	$D_2^U 6D_4$
		27	20	14	13	$D_1^U D_2^U 6D_4$
		28	21	14	14	$7D_4$
		29	22	13	14	$D_1^U D_3^U 6D_4$
		30	23	15	14	$D_2^U 7D_4$
		31	23	16	15	$D_1^U D_2^U 7D_4$
		32	24	16	16	$8D_4$
		33	25	15	16	$D_1^U D_3^U 7D_4$
		34	26	17	16	$D_2^U 8D_4$
		35	26	18	17	$D_1^U D_2^U 8D_4$
		36	27	18	18	$9D_4$
		37	28	17	18	$D_1^U D_3^U 8D_4$
		38	29	19	18	$D_2^U 9D_4$
		39	29	20	19	$D_1^U D_2^U 9D_4$
		40	30	20	20	$10D_4$
		41	31	19	20	$D_1^U D_3^U 9D_4$
		42	32	21	20	$D_2^U 10D_4$
		43	32	22	21	$D_1^U D_2^U 10D_4$
		44	33	22	22	$11D_4$

TABLE 6.4 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	4	45	34	21	22	$D_1^U D_3^U 10D_4$
		46	35	23	22	$D_2^U 11D_4$
		47	35	24	23	$D_1^U D_2^U 11D_4$
		48	36	24	24	$12D_4$
		49	37	23	24	$D_1^U D_3^U 11D_4$
		50	38	25	24	$D_2^U 12D_4$
4	6	2	1	1	0	$D_1$ (Table 5.2)
		7	4	1	1	$D_6$
		9	5	2	1	$D_1^U D_6$
		10	5	2	2	$D_{10}$
		11	7	1	1	$D_1^U D_9$
		12	6	3	2	$D_{13}$
		13	8	2	1	$2D_1^U D_9$
		14	6	3	3	$D_1^U D_{12}$
		15	8	2	3	$D_{16}$
		16	9	3	2	$2D_1^U D_{14}$
		17	9	3	3	$D_1^U D_{16}$
		18	9	3	4	$D_{17}$
		19	10	4	3	$2D_1^U D_{16}$
		20	10	4	4	$D_1^U D_{17}$
		21	12	3	3	$D_1^U D_6^U D_{16}$
		22	11	5	4	$2D_1^U D_{17}$
		23	13	4	3	$2D_1^U D_6^U D_{14}$



TABLE 6.4 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	6	24	13	4	4	$D_1^U D_6^U D_{16}$
		25	13	4	5	$D_6^U D_{17}$
		26	14	5	4	$D_6^U D_{16}$
		27	14	5	5	$D_1^U D_6^U D_{17}$
		28	14	5	6	$D_{10}^U D_{17}$
		29	15	6	5	$2D_1^U D_6^U D_{17}$
		30	15	6	6	$D_1^U D_{10}^U D_{17}$
		31	17	5	5	$D_1^U 2D_6^U D_{16}$
		32	16	7	6	$2D_1^U D_{10}^U D_{17}$
		33	18	6	5	$2D_1^U 2D_5^U D_{15}$
		34	18	6	6	$D_1^U 2D_6^U D_{17}$
		35	18	6	7	$D_6^U D_{10}^U D_{17}$
		36	19	7	6	$2D_1^U 2D_6^U D_{17}$
		37	19	7	7	$D_1^U D_6^U D_{10}^U D_{17}$
		38	19	7	8	$2D_{10}^U D_{17}$
		39	20	8	7	$2D_1^U D_6^U D_{10}^U D_{17}$
		40	20	8	8	$D_1^U 2D_{10}^U D_{17}$
		41	22	7	7	$D_1^U 3D_6^U D_{17}$
		42	21	9	8	$2D_1^U 2D_{10}^U D_{17}$
		43	23	8	7	$2D_1^U 3D_6^U D_{17}$
		44	23	8	8	$D_1^U 2D_6^U D_{10}^U D_{17}$
		45	23	8	9	$D_6^U 2D_{10}^U D_{17}$

TABLE 6.4 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	6	46	24	9	8	$2D_1^U \cup 2D_6^U \cup D_{10}^U \cup D_{17}$
		47	24	9	9	$D_1^U \cup D_6^U \cup 2D_{10}^U \cup D_{17}$
		48	24	9	10	$3D_{10}^U \cup D_{17}$
		49	25	10	9	$2D_1^U \cup D_6^U \cup 2D_{10}^U \cup D_{17}$
		50	25	10	10	$D_1^U \cup 3D_{10}^U \cup D_{17}$
4	8	2	0	1	0	$D_1$ (Table 5.3)
		10	2	2	1	$D_3$
		12	5	1	1	$D_5$
		14	5	2	1	$D_1^U \cup D_5$
		15	4	2	2	$D_7$
		16	3	2	3	$D_8$
		17	4	3	2	$D_1^U \cup D_7$
		19	7	2	2	$D_{11}$
		20	6	2	3	$D_{14}$
		21	7	3	2	$D_1^U \cup D_{11}$
		22	6	3	3	$D_1^U \cup D_{14}$
		23	10	1	2	$D_{15}$
		24	9	2	3	$D_{17}$
		26	9	3	3	$D_1^U \cup D_{17}$
		27	4	4	5	$D_4^U \cup D_7$
		28	9	4	3	$2D_1^U \cup D_{17}$
		29	4	5	5	$D_1^U \cup D_4^U \cup D_7$
		30	12	2	3	$D_{18}$

TABLE 6.4 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	8	31	12	3	3	$D_5^U D_{11}$
		32	12	4	3	$D_{19}$
		33	7	5	5	$D_1^U D_4^U D_{11}$
		34	11	4	4	$D_1^U D_5^U D_{14}$
		35	10	4	5	$D_7^U D_{14}$
		36	14	3	4	$D_5^U D_{17}$
		37	10	5	5	$D_1^U D_7^U D_{14}$
		38	14	4	4	$D_1^U D_5^U D_{17}$
		39	13	4	5	$D_7^U D_{17}$
		40	14	5	4	$2D_1^U D_5^U D_{17}$
		41	13	5	5	$D_1^U D_7^U D_{17}$
		42	12	5	6	$D_1^U D_8^U D_{17}$
		43	16	4	5	$D_{11}^U D_{17}$
		44	17	5	4	$D_5^U D_{19}$
		45	16	5	5	$D_1^U D_{11}^U D_{17}$
		46	15	5	6	$D_1^U D_{14}^U D_{17}$
		47	16	6	5	$D_7^U D_{19}$
		48	15	6	6	$2D_1^U D_{14}^U D_{17}$
		49	10	7	8	$D_1^U D_4^U D_7^U D_{14}$
		50	18	5	6	$D_1^U 2D_{17}$
4	10	10	3	0	0	$D_1$ (Table 5.4)
		13	3	1	1	$D_2$

TABLE 6.4 ( CONTINUED ).

k	v	b	$\lambda_0$	$\lambda_1$	$\lambda_2$	STRUCTURE
4	10	14	5	1	0	$D_3$
		15	2	2	1	$D_5$
		20	5	2	1	$D_9$
		25	6	2	2	$D_{12}$
		26	4	2	3	$D_{14}$
		27	5	3	2	$D_{15}$
		28	3	3	3	$D_2^U D_4$
		30	9	2	2	$D_{17}$
		31	7	2	3	$D_{18}$
		32	8	3	2	$D_{21}$
		33	6	3	3	$D_2^U D_8$
		34	4	3	4	$D_{22}$
		36	10	2	3	$D_{23}$
		37	11	3	2	$D_{26}$
		38	9	3	3	$D_2^U D_{12}$
		39	7	3	4	$D_2^U D_{14}$
		40	8	4	3	$D_2^U D_{15}$
		41	4	4	5	$D_4^U D_{14}$
		42	13	2	3	$D_{27}$
		43	12	3	3	$D_2^U D_{17}$
		44	10	3	4	$D_2^U D_{18}$
		45	11	4	3	$D_2^U D_{21}$

**TABLE 6.4 ( CONTINUED ) .**

<b>k</b>	<b>v</b>	<b>b</b>	$\lambda_0$	$\lambda_1$	$\lambda_2$	<b>STRUCTURE</b>
4	10	46	7	4	5	$D_4^U D_{18}$
		47	8	5	4	$D_4^U D_{21}$
		48	6	5	5	$D_2^U D_4^U D_8$
		49	13	3	4	$D_2^U D_{23}$
		50	12	4	4	$D_8^U D_{17}$



## CHAPTER 7

### 7.1 COMPUTER ALGORITHM FOR FINDING ALL BINARY E MB GD TIB DESIGNS :

The program " DSGN " (See Appendix-A) uses an algorithm based on the discussion in Section 2.2 and generates all possibly existing, admissible, and binary E MB GD TIB designs for a given (v,k). The details are as follows.

INPUT: An input disk file named " PARAMS " ( abbreviation of parameters of the design ) is created and saved by the user with the following entries.

K (represents the block size k) in the first column.

V (represents the number of test-treatments v) in the next three columns, right justified.

MAXB (represents the maximum number of blocks for which the designs are required, i.e.  $1 \leq b \leq \text{MAXB}$ ) in the next three columns, right justified.

#### ALGORITHM:

Step 1. Generate a value B [ B = 1(1)MAXB ]

Step 2. Call Subroutine " VECTOR " which generates the vector M [represents the vector  $\underline{m} = ( m_1, m_2, \dots, m_{k-1} )$ ].

If all possible vectors  $M$  have been generated then go to Step 1 to generate the next value of  $B$ , otherwise continue.

Step 3. Compute  $R$  (represents  $r$ , the replications of test-treatments) and check whether it is an integer.

If not an integer then go to Step 2 to generate next possible vector  $M$ , otherwise continue.

Step 4. Compute  $LD_0$  (represents  $\lambda_0$ ) and check whether it is an integer.

If not an integer then go to Step 2, otherwise continue.

Step 5. Compute  $LD_1$  (represents  $\lambda_1$ ) and check whether it is an integer.

If not an integer then go to Step 2, otherwise continue.

Step 6 Compute  $LD_2$  (represents  $\lambda_2$ )

Step 7 Call Subroutine " EXIST " which checks for the non-existence of the design.

If the design is non-existent go to Step 2, otherwise continue.

Step 8 Call Subroutine " DTYP " which decides the type of design as BIB, BTIB, or MBTIB designs.

Step 9 Call Subroutine " ADMSB " ( see remark below ) which checks whether the design is admissible or not ( see Section 2.1.4 ).

If inadmissible go to Step 2, otherwise continue.

Step 10 Call Subroutine " OPTIM " which finds the optimal design. Then go to Step 2.

The program will automatically stop when the B values are exhausted ( i.e. when  $B > \text{MAXB}$  ) by giving a message " JOB COMPLETED ! " .

OUTPUT:

1. All designs are stored in a disk file (Created by the program) named as " DESIGN " .

NOTE: The file "DESIGN" can be written (properly and automatically formatted) on the printer by running a program "PRT1" given in Appendix-D.

The sample output of this file is given in Appendix-G.

2. All optimal designs classified in BIB, BTIB, or MBTIB designs are stored in a disk file (created by the program) named as "OPTIMAL".

NOTE: The file "OPTIMAL" can be printed (properly and automatically formatted) by running a program "PRT2" given in Appendix-E.

A sample output of this file is given in Appendix-H.

REMARK: In order to check the admissibility of a design a matrix named "ADM" is created by the program. The  $i$  th row of this matrix corresponds to  $\lambda_0 = i - 1$ . The parameters  $b, \lambda_1, \lambda_2$  are entered into the



$(\lambda_0 + 1)$  th row of this matrix. So, "ADM" looks as follows ( for example ).

	$b$	$\lambda_1$	$\lambda_2$	
row 1	...	...	...	$(\lambda_0 = 0)$
2	...	...	...	$(\lambda_0 = 1)$
⋮	⋮	⋮	⋮	⋮
8	7	2	1	$(\lambda_0 = 8)$
⋮	⋮	⋮	⋮	⋮

Hence, the program knows that for  $\lambda_0 = 8$  (for example) a design with  $b = 7, \lambda_1 = 2, \lambda_2 = 1$  is already available. So, if the program find a candidate of the form  $(b, \lambda_0 = 8, \lambda_1 = 2, \lambda_2 = 1)$  with  $b \geq 7$  it discard it treating as equivalent or inadmissible design.

In this way all equivalent or inadmissible designs can be discarded.

**7.2 COMPUTER ALGORITHM/PROGRAM FOR FINDING A POSSIBLE SET OF DESIGNS TO BE MCCGD, AND THE STRUCTURED CONSTRUCTION OF ALL DESIGNS :**

---

The program "SEARCH" ( given in Appendix-B ), helps in finding a MCCGDs for a given  $(v,k)$ .

INPUT:1. Disk file "PARAMS" created by the user and updated by the program "DSGN" ( Appendix-A)

2. Disk file "DESIGN" created by the program "DSGN"  
(Appendix-A)

OUTPUT:Disk file "GD" created by the program "SEARCH"  
contains the possible candidates for the MCCGDs for the  
(v,k).

REMARK: Each design in the file "GD" is tried to be  
constructed manually. The design which cannot be  
constructed ( a little trial and error is  
required ) should be deleted from the file "GD"  
by the user. In this way the file "GD" is updated  
and then saved containing only the MCCGDs.

The program "STRUCT" ( given in  
Appendix-C ) is used to find the structure of a design by  
finding the possible union of replicates ( copies ) of  
generator designs which is equivalent to the design  
structured.

INPUT:1. Disk file "PARAMS" updated by the program  
"DSGN" ( Appendix-A ).

2. Disk file "DESIGN" created by the program "DSGN"  
( Appendix-A ).

3. Disk file "GD" updated by the user as MCCGDs.

OUTPUT:Disk file "STRC" created by the program "STRUCT"  
contains the parameters of the designs structured  
followed by the labels of the generator designs used and

the number of copies of them required.

NOTE: The file "STRC" can be written on a printer (properly and automatically formatted) by running the program "PRT3" given in Appendix-F. A sample output of the file is given in Appendix-I.

REMARK: All the program given here are written in standard FORTRAN and are run on Burroughs large system main frame of the Middle East Technical University.



## CHAPTER 8

### CONCLUSIONS:

An Extended Class of Treatment Incomplete Block Designs is introduced. This class contains all binary Balanced Incomplete Block (BIB) designs, Balanced Treatment Incomplete Block (BTIB) designs, and the Most Balanced Group Divisible Treatment Incomplete Block (MB GD TIB) designs. The minimal complete classes of generator designs (MCCGDs) are constructed for the practical range of  $k$  (the size of blocks) and  $v$  (number of test-treatments) values and the minimal complete nature of such classes is proved through case by case enumeration. A general non-existence theorem for GD TIB designs is given and the usefulness of the theorem is explained through examples. A general method/algorithm for the construction of all admissible E MB GD TIB designs is suggested and applied in order to construct the designs for various parameters values.

The A-optimality criterion is applied in order to obtain optimal designs for a given  $(v, k, b)$ . Elaborate computer programs/subroutines in standard FORTRAN are prepared which are not only useful in finding the MCCGD but also suggest possible structure of these designs. The program sorts out the A-optimal designs as well. Some trial and error is required in finding the actual structure of these designs. These efforts are, however, minimized through the use of the computer program provided here.

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APPENDICES:

APPENDIX-A: ( PROGRAM " DSGN " )

```
FILE 7(KIND=DISK,TITLE="PARAMS",FILETYPE=7)
FILE 8(KIND=DISK,TITLE="ADMSB",FILETYPE=7)
FILE 9(KIND=DISK,TITLE="DESIGN",FILETYPE=7)
FILE 10(KIND=DISK,TITLE="OPTIMAL",FILETYPE=7)
INTEGER V,U,W,STD,STO,CTN,PB,PM(8),B,R0,R,TCP,TTP,S3,
*      M(8),ADM(100,12)
REAL DSG(1700,20),OPT(320,20)
READ(7,2000)K,V,MAXB,MLDO,LB,U,W,STD,STO,CTN,PB,
*(PM(J),J=1,K-1)
CTN=CTN+1
IF(LB.GT.MAXB) GO TO 1800
IF(LB.NE.0) GO TO 1900
DO 10 ROW = 1,100
DO 10 COL = 1,12
10 ADM(ROW,COL)=-1
PB=1
STD=1
STO=1
MINB=1
W=1
50 DO 200 B = MINB,MAXB
PRINT/,"DESIGN IN PROCESS B=",B
PB=B
CALL VECTOR(K,B,R0,PM,M,I,N,S3,MAXL,&100,&700,&1700)
100 IF(R0.GT.(K-1)*B) GO TO 600
IF(MOD(B*K-R0,V).EQ.0) GO TO 400
GO TO 600
150 LB=PB
IF(OPT(1+(K+7)*(STO-1),4*W-3).EQ.0) GO TO 200
W=W+1
IF(W.EQ.6) GO TO 300
200 CONTINUE
PRINT/,"JOB COMPLETED !"
GO TO 1700
300 W=1
STO=STO+1
GO TO 200
400 R=(B*K-R0)/V
500 TCP=(K-1)*R0-S3
IF(MOD(TCP,V).EQ.0) GO TO 900
600 CALL VECT(K,B,R0,PM,M,I,N,S3,MAXL,&100,&700,&1700)
700 DO 800 J = 1,K-1
PM(J)=0
800 CONTINUE
GO TO 150
```

```

900 LD0=TCP/V
   IF(LD0.LT.MAXL) GO TO 600
   IF(LD0.GT.MLD0) MLD0=LD0
   TTP=K*(K-1)*B/2-S3/2-TCP
   N1=4*TTP
   N2=2*V*(V-1)
   IF(MOD(N1,N2).EQ.0) GO TO 910
   IF(MOD(N1-V**2,N2).EQ.0) GO TO 920
   IF(MOD(N1+V**2,N2).EQ.0) GO TO 930
   GO TO 600
910 LD1=N1/N2
   LD2=LD1
   GO TO 1200
920 LD1=(N1-V**2)/N2
   LD2=LD1+1
   GO TO 1200
930 LD1=(N1+V**2)/N2
   LD2=LD1-1
1200 IF(LD2.LT.0.OR.LD1.GT.R) GO TO 600
   CALL EXIST(V,B,R,LD1,LD2,&1300,&600)
1300 IF(LD1.EQ.LD2) GO TO 1320
   IF(LD1.EQ.LD2+1.OR.LD1.EQ.LD2-1) GO TO 1310
   DSN="  PB"
   GO TO 1360
1310 DSN="MBTIB"
   GO TO 1360
1320 IF(LD0.EQ.LD1) GO TO 1340
1330 DSN=" BTIB"
   GO TO 1360
1340 IF(R.EQ.R0) GO TO 1350
   GO TO 1330
1350 DSN="  BIB"
1360 IF(DSN.EQ."  PB") GO TO 600
   IF(LD0.EQ.0) GO TO 1370
   NUM=(LD0+LD1+V*LD2/2.)*(LD0+V*LD2/2.)-V*(V-2)*LD2**2
   * /4.
   DEN=LD0*(LD0+V*LD2)*(LD0+V*LD1/2.+V*LD2/2.)
   TRC=V*K*NUM/DEN
   GO TO 1380
1370 TRC=-1
1380 CALL ADMSB(K,B,AD,LD0,LD1,LD2,DSN,ADM,M,PM,&1400,&600)
1400 U=U+1
   LB=B
   DSG(1+(K+8)*(STD-1),U)=B
   DSG(2+(K+8)*(STD-1),U)=R0
   DSG(3+(K+8)*(STD-1),U)=R
   DO 1410 ROW = 1,K-1
1410 DSG(ROW+3+(K+8)*(STD-1),U)=M(ROW)
   DSG((K+8)*STD-5,U)=LD0
   DSG((K+8)*STD-4,U)=LD1
   DSG((K+8)*STD-3,U)=LD2

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        DSG((K+8)*STD-2,U)=TRC
        DSG((K+8)*STD-1,U)=DSN
        DSG((K+8)*STD,U)=AD
        IF(U.EQ.20) GO TO 1600
1500 IF(OPT(1+(K+7)*(STO-1),4*W-3).EQ.0) GO TO 1510
        IF(R0.GT.INT(B*K/2.+5)) GO TO 600
1510 CALL OPTIM(K,W,STO,TRC,DSN,M,OPT,B,R0,R,LD0,LD1,LD2)
        GO TO 600
1600 U=0
        STD=STD+1
        GO TO 1500
1700 DO 1750 J = 1,K-1
1750 PM(J)=M(J)
        CALL SAVE(K,V,MAXB,MLD0,LB,U,W,STD,STO,CTN,B,PM,ADM,DSG,
*           OPT)
1800 STOP
1900 CALL RINIT(K,MLD0,STD,STO,ADM,DSG,OPT)
        MINB=PB
        GO TO 50
2000 FORMAT(I1,2(X,I2),X,I3,2(X,I2),X,I1,X,I3,X,I2,X,I3,X,I2,
*           8(X,I3))
        END
        SUBROUTINE VECTOR(K,B,R0,PM,M,I,N,S3,MAXL,*,*,*)
        INTEGER S,R0,PM(8),M(8),S3,S1,S2
        DO 100 I = 1,K-1
        M(I)=PM(I)
100 CONTINUE
        I=0
110 IF(I.EQ.0) GO TO 120
        M(K-1)=-1
        GO TO 150
120 M(K-1)=M(K-1)-1
150 I=K-1
        N=B
        IF(K.EQ.2) GO TO 250
        DO 200 J = 1,K-2
        N=N-M(J)
200 CONTINUE
        ENTRY VECT(K,B,R0,PM,M,I,N,S3,MAXL,*,*,*)
250 M(I)=M(I)+1
        IF(M(I).GT.N) GO TO 700
        IF(TIME(2).GT.1400) RETURN3
        S1=0
        S2=0
        S3=0
        MAXL=0
        DO 300 L = 1,K-1
        S1=S1+M(L)
        S2=S2+L*M(L)
        S3=S3+L*(L-1)*M(L)
        IF(M(L).GT.0) MAXL=L
300 CONTINUE
        IF(S1.GT.B) GO TO 400
        R0=S2

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        RETURN1
400 I=I-1
      IF(I.EQ.0) GO TO 650
      IF(I.EQ.1) GO TO 600
      N=B
      DO 450 J = 1,I-1
      N=N-M(J)
450 CONTINUE
      IF(M(I).EQ.N) GO TO 550
500 M(I)=M(I)+1
      GO TO 110
550 IF(I.EQ.1) GO TO 650
      M(I)=0
      GO TO 400
600 IF(M(I).EQ.B) GO TO 650
      GO TO 500
650 RETURN2
700 DO 750 J = 1,K-1
      PM(J)=M(J)
750 CONTINUE
      GO TO 400
      END
      SUBROUTINE EXIST(V,B,R,LD1,LD2,*,*)
      INTEGER V,B,R
      IF(R.LT.LD1.OR.R.LT.V*(LD2-LD1)/2+LD1) GO TO 300
      IF(B.GE.V) GO TO 250
      IF(R.EQ.LD1) GO TO 200
      IF(LD1.LT.LD2) GO TO 100
      GO TO 300
100 IF(R.EQ.LD1+V*(LD2-LD1)/2) GO TO 150
      GO TO 300
150 IF(B.EQ.V-1) GO TO 250
      GO TO 300
200 IF(LD1.EQ.LD2) GO TO 250
      IF(B.GE.2) GO TO 250
      GO TO 300
250 RETURN1
300 RETURN2
      END
      SUBROUTINE ADMSB(K,B,AD,LD0,LD1,LD2,DSN,ADM,M,PM,*,*)
      INTEGER B,ADM(100,12),M(8),PM(8)
      J=10
      AD="  NAD"
400 IF(B.EQ.ADM(LD0+1,J).AND.LD1.EQ.ADM(LD0+1,J+1).AND.LD2.
*   EQ.ADM(LD0+1,J-2)) GO TO 850
      IF(B.EQ.ADM(LD0+1,J).AND.LD1.LE.ADM(LD0+1,J+1).AND.LD2.
*   LE.ADM(LD0+1,J-2)) GO TO 750
      IF(LD1.LE.ADM(LD0+1,J+1).AND.LD2.LE.ADM(LD0+1,J+2))
*   GO TO 750
450 ADM(LD0+1,J)=B
      ADM(LD0+1,J+1)=LD1
      ADM(LD0+1,J+2)=LD2
      IF(J.EQ.10) GO TO 500

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RETURN1
500 AD=" ADM"
    INDC=2
550 IF(DSN.EQ." BIB") GO TO 700
    IF(DSN.EQ." BTIB") GO TO 650
    J=7
600 IF(INDC.EQ.1) GO TO 400
    GO TO 450
650 J=4
    GO TO 600
700 J=1
    GO TO 600
750 IF(J.EQ.10) GO TO 800
    RETURN2
800 INDC=1
    GO TO 550
850 NN=0
    DO 900 JJ = 1,K-1
        IF(M(JJ).EQ.PM(JJ)) NN=NN+1
900 CONTINUE
    IF(NN.EQ.K-1) RETURN2
    GO TO 450
END
SUBROUTINE OPTIM(K,W,STO,TRC,DSN,M,OPT,B,R0,R,LD0,LD1,
*           LD2)
    INTEGER W,STO,M(8),B,R0,R,ROW
    REAL OPT(320,20)
    J=1
100 IF(TRC.LT.OPT((K+7)*STO-1,J+4*(W-1))) GO TO 400
150 IF(J.EQ.1) GO TO 200
    RETURN
200 IF(DSN.EQ." BIB") GO TO 350
    IF(DSN.EQ." BTIB") GO TO 300
    IF(DSN.EQ."MBTIB") GO TO 250
    IF(J.EQ.1) RETURN
    GO TO 100
250 J=4
    GO TO 100
300 J=3
    GO TO 100
350 J=2
    GO TO 100
400 IF(TRC.LT.0) TRC=999.9
    OPT(1+(K+7)*(STO-1),J+4*(W-1))=B
    OPT(2+(K+7)*(STO-1),J+4*(W-1))=R0
    OPT(3+(K+7)*(STO-1),J+4*(W-1))=R
    DO 450 ROW = 1,K-1
        OPT(ROW+3+(K+7)*(STO-1),J+4*(W-1))=M(ROW)
450 CONTINUE
    OPT((K+7)*STO-4,J+4*(W-1))=LD0
    OPT((K+7)*STO-3,J+4*(W-1))=LD1
    OPT((K+7)*STO-2,J+4*(W-1))=LD2

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```

OPT((K+7)*STO-1,J+4*(W-1))=TRC
OPT((K+7)*STO,J+4*(W-1))=DSN
GO TO 150
END
SUBROUTINE SAVE(K,V,MAXB,MLD0,LB,U,W,STD,STO,CTN,PB,PM,
*
ADM,DSG,OPT)
INTEGER V,U,W,STD,STO,CTN,PB,PM(8),ADM(100,12),ROW,COL,
*
ST
REAL DSG(1700,20),OPT(320,20)
REWIND(7)
WRITE(7,500)K,V,MAXB,MLD0,LB,U,W,STD,STO,CTN,PB,(PM(J),
*
J=1,K-1)
LOCK(7)
REWIND(8)
WRITE(8,600)((ADM(ROW,COL),COL=1,12),ROW=1,MLD0+1)
LOCK(8)
REWIND(9)
DO 200 ST = 1,STD
WRITE(9,700)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20),ROW+1
*
,K+5)
WRITE(9,750)(DSG((K+8)*ST-2,COL),COL=1,20)
WRITE(9,800)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20),ROW=
*
K+7,K+8)
200 CONTINUE
LOCK(9)
REWIND(10)
DO 400 ST = 1,STO
WRITE(10,700)((OPT(ROW+(K+7)*(ST-1),COL),COL=1,20),ROW=
*
1,K+5)
WRITE(10,750)(OPT((K+7)*ST-1,COL),COL=1,20)
WRITE(10,800)(OPT((K+7)*ST,COL),COL=1,20)
400 CONTINUE
LOCK(10)
RETURN
500 FORMAT(I1,2(X,I2),X,I3,2(X,I2),X,I1,X,I3,X,I2,X,I3,X,I2
*
,8(X,I3))
600 FORMAT(4(3(X,I3),2X))
700 FORMAT(20(X,F5.1))
750 FORMAT(20(X,F5.2))
800 FORMAT(20(X,A5))
END
SUBROUTINE RINIT(K,MLD0,STD,STO,ADM,DSG,OPT)
INTEGER STD,STO,ADM(100,12),ROW,COL,ST
REAL DSG(1700,20),OPT(320,20)
READ(8,850)((ADM(ROW,COL),COL=1,12),ROW=1,MLD0+1)
DO 100 ROW = MLD0+2,100
DO 100 COL = 1,12
ADM(ROW,COL)=-1
100 CONTINUE
DO 200 ST =1,STD
READ(9,900)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20),ROW=1,
*
K+5)

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      READ(9,920) (DSG((K+8)*ST-2, COL), COL=1, 20)
      READ(9,950) ((DSG(ROW+(K+8)*(ST-1), COL), COL=1, 20), ROW=
*          K+7, K+8)
200 CONTINUE
      DO 300 ST = 1, STO
      READ(10,900) ((OPT(ROW+(K+7)*(ST-1), COL), COL=1, 20), ROW=1
*          , K+5)
      READ(10,920) (OPT((K+7)*ST-1, COL), COL=1, 20)
      READ(10,950) (OPT((K+7)*ST, COL), COL=1, 20)
300 CONTINUE
      RETURN
850 FORMAT(4(3(X,I3), 2X))
900 FORMAT(20(X,F5.1))
920 FORMAT(20(X,F5.2))
950 FORMAT(20(X,A5))
      END

```

APPENDIX-B: ( PROGRAM "SEARCH" )

```

FILE 7 (KIND=DISK, TITLE="PARAMS", FILETYPE=7)
FILE 9 (KIND=DISK, TITLE="DESIGN", FILETYPE=7)
FILE 10 (KIND=DISK, TITLE="GD", FILETYPE=7)
      INTEGER V, STD, ST, STS, ROW, COL, CL, B, DB, DL0, DL1, DL2, BK(100
*          ), LD0(100), LD1(100), LD2(100)
      REAL DSG(1700, 20)
      READ(7, 2000) K, V, STD
      DO 100 ST = 1, STD
      READ(9, 2500) ((DSG(ROW+(K+8)*(ST-1), COL), COL=1, 20), ROW=1
*          , K+5)
      READ(9, 3000) (DSG(K+6+(K+8)*(ST-1), COL), COL=1, 20)
      READ(9, 3500) ((DSG(ROW+(K+8)*(ST-1), COL), COL=1, 20), ROW+K
*          +7, k+8)
100 CONTINUE
      MAXJ=1
      BK(1)=DSG(1, 1)
      LD0(1)=DSG(K+3, 1)
      LD1(1)=DSG(K+4, 1)
      LD2(1)=DSG(K+5, 1)
      DO 870 ST = 1, STD
      DO 870 COL = 1, 20
      B=DSG(1+(K+8)*(ST-1), COL)
      IF(B.EQ.999.OR.B.EQ.0) GO TO 180
      L0=DSG(K+3+(K+8)*(ST-1), COL)
      L1=DSG(K+4+(K+8)*(ST-1), COL)
      L2=DSG(K+5+(K+8)*(ST-1), COL)
      DO 150 J = 0, MAXJ-1
      IF(B.NE.BK(MAXJ-J)) GO TO 160
      IF(L0.EQ.LD0(MAXJ-J).AND.L1.EQ.LD1(MAXJ-J).AND.L2.EQ.
*          LD2(MAXJ-J)) GO TO 870

```

```

150 CONTINUE
160 DO 850 J = 1,MAXJ
    IF(B-BK(J).GE.DSG(1,1)) GO TO 500
    IF(ST.EQ.1.AND.COL.EQ.1) GO TO 870
    GO TO 750
180 DO 200 J = 1,MAXJ
200 WRITE(10,4000)BK(J),LD0(J),LD1(J),LD2(J)
    LOCK(10)
    PRINT/,"JOB COMPLETED !"
    STOP
500 IF(L0.GE.LD0(J).AND.L1.GE.LD1(J).AND.L2.GE.LD2(J))
    * GO TO 550
    GO TO 850
550 DB=B-BK(J)
    DL0=L0-LD0(J)
    DL1=L1-LD1(J)
    DL2=L2-LD2(J)
    STS=ST
600 IF(DB.GT.DSG(1+(K+8)*(STS-1),1)) GO TO 650
    IF(STS.EQ.1) GO TO 650
    STS=STS-1
    GO TO 600
650 DO 700 CL = 1,20
    IF(DB.EQ.DSG(1+(K+8)*(STS-1),CL)) GO TO 800
    IF(DB.LT.DSG(1+(K+8)*(STS-1),CL)) GO TO 850
700 CONTINUE
720 STS=STS+1
    IF(DB.EQ.DSG(1+(K+8)*(STS-1),1)) GO TO 650
    GO TO 850
750 MAXJ=MAXJ+1
    BK(MAXJ)=B
    LD0(MAXJ)=L0
    LD1(MAXJ)=L1
    LD2(MAXJ)=L2
    GO TO 870
800 IF(DL0.EQ.DSG(K+3+(K+8)*(STS-1),CL).AND.DL1.EQ.DSG(K+4+
    * (K+8)*(STS-1),CL).AND.DL2.EQ.DSG(K+5+(K+8)*(STS-1),
    * CL)) GO TO 870
    GO TO 700
850 CONTINUE
    GO TO 750
870 CONTINUE
    GO TO 180
2000 FORMAT(I1,X,I2,16X,I3)
2500 FORMAT(20(X,F5.1))
3000 FORMAT(20(X,F5.2))
3500 FORMAT(20(X,A5))
4000 FORMAT(4(X,I3))
    END

```

APPENDIX-C: ( PROGRAM "STRUCT" )

```
FILE 7(KIND=DISK,TITLE="PARAMS",FILETYPE=7)
FILE 9(KIND=DISK,TITLE="DESIGN",FILETYPE=7)
FILE 10(KIND=DISK,TITLE="GD",FILETYPE=7)
FILE 11(KIND=DISK,TITLE="STRC",FILETYPE=7)
INTEGER V,STD,ST,STS,ROW,COL,CL,B,DB,DL0,DL1,DL2,BK(100
*) ,LD0(100),LD1(100),LD2(100),STR(10000,20)
REAL DSG(1700,20)
READ(7,2000)K,V,STD
DO 100 ST = 1,STD
READ(9,2500)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20),ROW=1
*) ,K+5)
READ(9,3000)(DSG(K+6+(K+8)*(ST-1),COL),COL=1,20)
READ(9,3500)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20),ROW=K
*) +7,K+8)
100 CONTINUE
N=0
DO 150 J = 1,1000
READ(10,8050,END=200)BK(J),LD0(J),LD1(J),LD2(J)
N=N+1
150 CONTINUE
200 DO 250 ST = 1,100
DO 250 ROW = 1,N+4
DO 250 COL = 1,20
STR(ROW+(N+4)*(ST-1),COL)=999
250 CONTINUE
DO 500 ST = 1,STD
DO 450 COL = 1,20
B=DSG(1+(K+8)*(ST-1),COL)
IF(B.GE.999.OR.B.EQ.0) GO TO 560
L0=DSG(K+3+(K+8)*(ST-1),COL)
L1=DSG(K+4+(K+8)*(ST-1),COL)
L2=DSG(K+5+(K+8)*(ST-1),COL)
STR(1+(N+4)*(ST-1),COL)=B
STR(2+(N+4)*(ST-1),COL)=L0
STR(3+(N+4)*(ST-1),COL)=L1
STR(4+(N+4)*(ST-1),COL)=L2
IF(B.GT.BK(N)) GO TO 350
DO 300 J = 1,N
IF(B.EQ.BK(J)) GO TO 1000
300 CONTINUE
350 DO 400 J = 1,N
DB=B-BK(J)
DL0=L0-LD0(J)
DL1=L1-LD1(J)
DL2=L2-LD2(J)
STS=ST
IF(DB.GE.DSG(1,1)) GO TO 600
400 CONTINUE
450 CONTINUE
```

```

500 CONTINUE
560 REWIND(11)
    DO 550 ST = 1,STD
        WRITE(11,9050)((STR(ROW+(N+4)*(ST-1),COL),COL=1,20),ROW
*           =1,N+4)
550 CONTINUE
    LOCK(11)
    PRINT/,"JOB COMPLETED !"
    STOP
600 IF(DL0.LT.0.OR.DL1.LT.0.OR.DL2.LT.0) GO TO 400
650 IF(DB.GT.DSG(1+(K+8)*(STS-1),1)) GO TO 700
    IF(STS.EQ.1) GO TO 700
    STS=STS-1
    GO TO 650
700 DO 700 CL = 1,20
    IF(DB.EQ.DSG(1+(K+8)*(STS-1),CL)) GO TO 800
    IF(DB.LT.DSG(1+(K+8)*(STS-1),CL)) GO TO 400
750 CONTINUE
    STS=STS+1
    IF(DB.EQ.DSG(1+(K+8)*(STS-1),1)) GO TO 700
    GO TO 400
800 IF(DL0.EQ.DSG(K+3+(K+8)*(STS-1),CL).AND.DL1.EQ.DSG(K+4+
*   (K+8)*(STS-1),CL.AND.DL2.EQ.DSG(K+5+(K+8)*(STS-1),CL
*   )) GO TO 850
    GO TO 750
850 STR(J+4+(N+4)*(ST-1),COL)=1
    DO 900 JJ = 1,N
        IF(STR(JJ+4+(N+4)*(STS-1),CL).EQ.999) GO TO 900
        IF(STR(JJ+4+(N+4)*(ST-1),COL).EQ.999) GO TO 950
        STR(JJ+4+(N+4)*(ST-1),COL)=STR(JJ+4+(N+4)*(ST-1),COL)+
*   STR(JJ+4+(N+2)*(STS-1),CL)
900 CONTINUE
    GO TO 450
950 STR(JJ+4+(N+4)*(ST-1),COL)=STR(JJ+4+(N+4)*(STS-1),CL)
    GO TO 900
1000 IF(L0.EQ.LD0(J).AND.L1.EQ.LD1(J).AND.L2.EQ.LD2(J))
*   GO TO 1050
    GO TO 300
1050 STR(J+4+(N+4)*(ST-1),COL)=1
    GO TO 450
2000 FORMAT(I1,X,I2,16X,I3)
2500 FORMAT(20(X,F5.1))
3000 FORMAT(20(X,F5.2))
3500 FORMAT(20(X,A5))
8050 FORMAT(4(X,I3))
9050 FORMAT(20(X,I3))
    END

```



APPENDIX-D: ( PROGRAM "PRT1" )

```
FILE 7(KIND=DISK,TITLE="PARAMS",FILETYPE=7)
FILE 8(KIND=DISK,TITLE="DESIGN",FILETYPE=7)
FILE 9(KIND=PRINTER)
      INTEGER V,STD,ST,ROW,COL
      REAL HD(17),DSG(1700,20)
      READ(7,500)K,V,STD
      WRITE(9,600)K,V
      HD(1)=" BI"
      HD(2)=" ROI"
      HD(3)=" RI"
      HD(K+3)="LD0I"
      HD(K+4)="LD1I"
      HD(K+5)="LD2I"
      HD(K+6)="TRCI"
      HD(K+7)="DSNI"
      HD(K+8)="ADMI"
      DO 100 ST = 1,STD
      READ(8,700,END=105)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20
*                          ),ROW+1,K+5)
      READ(8,750,END=105)(DSG(K+6+(K+8)*(ST-1),COL),COL=1,20)
      READ(8,800,END=105)((DSG(ROW+(K+8)*(ST-1),COL),COL=1,20
*                          ),ROW=K+7,K+8)
100 CONTINUE
105 IF(DSG(1+(K+8)*(STD-1),1).EQ.0) STD=STD-1
      MROW=K+6
      DO 120 ST = 1,STD
      IF(ST.EQ.STD) MROW=1
      DO 120 ROW = MROW,K+6
      DO 120 COL = 1,20
      IF(ROW.EQ.K+6) GO TO 130
      IF(DSG(ROW+(K+8)*(ST-1),COL).LE.0) DSG(ROW+(K+8)*(ST-1)
*      ,COL)=999999.
120 CONTINUE
      GO TO 150
130 IF(DSG(ROW+(K+8)*(ST-1),COL).LE.0) DSG(ROW+(K+8)*(ST-1)
*      ,COL)=999999.
      GO TO 120
150 DO 300 ST = 1,STD
      IF(DSG(1+(K+8)*(ST-1),1).EQ.999999..OR.DSG(1+(K+8)*(ST-
*      1),1).EQ.0) GO TO 320
      DO 200 ROW = 1,K+8
      IF(ROW.LE.2.OR.ROW.EQ.4.OR.ROW.EQ.K+3.OR.ROW.EQ.K+6.OR.
*      ROW.EQ.K+7) WRITE(9,900)
      IF(ROW.LE.3.OR.ROW.EQ.K+3.OR.ROW.EQ.K+4.OR.ROW.EQ.K+5)
*      WRITE(9,1000)(HD(ROW),(DSG(ROW+(K+8)*(ST-1),COL),COL
*      =1,20,HD(ROW))
      IF(ROW.GE.4.AND.ROW.LE.K+2) WRITE(9,1100)(ROW-3,(DSG(RO
*      W+(K+8)*(ST-1),COL),COL=1,20),ROW-3)
```

```

IF(ROW.EQ.K+6) WRITE(9,1200)(HD(ROW),(DSG(ROW+(K+8)*(ST
*   -1),COL).COL=1,20),HD(ROW))
IF(ROW.GT.K+6) WRITE(9,1300,END=320)(HD(ROW),(DSG(ROW+(
*   K+8)*(ST-1),COL),COL=1,20),HD(ROW))
200 CONTINUE
WRITE(9,900)
WRITE(9,1400)
IF(DSG(1+(K+8)*(ST-1),20).EQ.999999.) GO TO 320
PRINT/,"DESIGN IN PROCESS B=",INT(DSG(1+(K+8)*(ST-1),20
*   ))
300 CONTINUE
320 PRINT/,"JOB COMPLETED !"
350 STOP
500 FORMAT(I1,X,I2,16X,I3)
600 FORMAT(10X,"K = ",I2/10X,"-----"/10X,"V = ",I2/10X,"--
*   ----"/10X,"A L L D E S I G N S"/10X,21("-"))
700 FORMAT(20(X,F5.1))
750 FORMAT(20(X,F5.2))
800 FORMAT(20(X,A5))
900 FORMAT(X,"+---+",20("-----+"),"----+")
1000 FORMAT(X,"I",A4,20(I5,"I"),A4)
1100 FORMAT(X,"IM",I2,"I",20(I5,"I"),"M",I2,"I")
1200 FORMAT(X,"I",A4,20(F5.2,"I"),A4)
1300 FORMAT(X,"I",A4,20(A5,"I"),A4)
1400 FORMAT(/)
END

```

APPENDIX-E: ( PROGRAM "PRT2" )

```

FILE 7(KIND=DISK,TITLE="PARAMS",FILETYPE=7)
FILE 8(KIND=DISK,TITLE="OPTIMAL",FILETYPE=7)
FILE 9(KIND=PRINTER
INTEGER V,STO,ST,ROW,COL
REAL HD(16),OPT(320,20)
READ(7,500)K,V,STO
WRITE(9,600)K,V
HD(1)=" BI"
HD(2)=" ROI"
HD(3)=" RI"
HD(K+3)="LD0I"
HD(K+4)="LD1I"
HD(K+5)="LD2I"
HD(K+6)="TRCI"
HD(K+7)="DSNI"
DO 100 ST = 1,STO
READ(8,700,END=105)((OPT(ROW+(K+7)*(ST-1),COL),COL=1,20
*   ),ROW=1,K+ 5)
READ(8,750,END=105)(OPT(K+6+(K+7)*(ST-1),COL),COL=1,20
READ(8,800,END=105)(OPT(K+7+(K+7)*(ST-1),COL),COL=1,20)

```

```

100 CONTINUE
105 IF(OPT(1+(K+7)*(STO-1),1).EQ.0) STO=STO-1
    DO 120 ST = 1,STO
    DO 120 ROW = 1,K+6
    DO 120 COL = 1,20
    IF(ROW.EQ.K+6) GO TO 130
    IF(OPT(ROW+(K+7)*(ST-1),COL).GE.0) OPT(ROW+(K+7)*(ST-1)
*   ,COL)=999999.
120 CONTINUE
    GO TO 150
130 IF(OPT(ROW+(K+7)*(ST-1),COL).LE.0) OPT(ROW+(K+7)*(ST-1)
*   ,COL)=999999.
    GO TO 120
150 DO 300 ST = 1,STO
    IF(OPT(1+(K+7)*(ST-1),1).EQ.999999.) GO TO 320
    WRITE(9,900)
    WRITE(9,650)
    DO 200 ROW = 1,K+7
    IF(ROW.LE.2.OR.ROW.EQ.4.OR.ROW.EQ.K+3.OR.ROW.EQ.K+6.OR.
*   ROW.EQ.K+7) WRITE(9,900)
    IF(ROW.LE.3.OR.ROW.EQ.K+3.OR.ROW.EQ.K+4.OR.ROW.EQ.K+5)
*   WRITE(9,1000)(HD(ROW),(OPT(ROW+(K+7)*(ST-1),COL),COL=1
*   ,20),HD(ROW))
    IF(ROW.GE.4.AND.ROW.LE.K+2)WRITE(9,1100)(ROW-3,(OPT(ROW
*   +(K+7)*(ST-1),COL),COL=1,20),ROW-3)
    IF(ROW.EQ.K+6) WRITE(9,1200)(HD(ROW),(OPT(ROW+(K+7)*(ST
*   -1),COL),COL=1,20),HD(ROW))
    IF(ROW.GT.K+6)WRITE(9,1300,END=320)(HD(ROW),(OPT(ROW+(K
*   +7)*(ST-1),COL),COL=1,20),HD(ROW))
200 CONTINUE
    WRITE(9,900)
    WRITE(9,1400)
    IF(OPT(1+(K+7)*(ST-1),17).EQ.999999.) GO TO 320
    PRINT/,"DESIGN IN PROCESS B=",INT(OPT(1+(K+7)*(ST-1),17
*   ))
300 CONTINUE
320 PRINT/,"JOB COMPLETED !"
350 STOP
500 FORMAT(I1,X,I2,20X,I2)
600 FORMAT(10X,"K = ",I2/10X,"-----"/10X,"V = "I2/10X,"---
*   ---"/10X,"O P T I M A L D E S I G N S"/10X,29("-"))
650 FORMAT(X,"IDSGI",5(" ALL I BIB I BTIBIMBTIBI"),"DSGI")
700 FORMAT(20(X,F5.1))
750 FORMAT(20(X,F5.2))
800 FORMAT(20(X,A5))
900 FORMAT(X,"+----+",20("-----+"),"----+")
1000 FORMAT(X,"I",A4,20(I5,"I"),A4)
1100 FORMAT(X,"IM",I2,"I",20(I5,"I"),"M",I2,"I")
1200 FORMAT(X,"I",A4,20(F5.2,"I"),A4)
1300 FORMAT(X,"I",A4,20(A5,"I"),A4)
1400 FOEMAT(/)
    END

```

APPENDIX-F: ( PROGRAM "PRT3" )

```
FILE 7(KIND=DISK,TITLE="PARAMS",FILETYPE=7)
FILE 8(KIND=DISK,TITLE="GD",FILETYPE=7)
FILE 9(KIND=DISK,TITLE="STRC",FILETYPE=7)
FILE 10(KIND=PRINTER)
      INTEGER V,STD,BK(100),LD0(100),LD1(100),LD2(100),ROW,CO
*      L,ST,STR(1000,20)
      READ(7,8000)K,V,STD
      DO 100 J = 1,100
      READ(8,8050,END=150)BK(J),LD0(J),LD1(J),LD2(J)
      N=N+1
100 CONTINUE
150 DO 200 ST = 1,STD
      READ(9,9050)((STR(ROW+(N+4)*(ST-1),COL),COL=1,20),ROW=1
*                ,N+4)
200 CONTINUE
      HD(1)=" B "
      HD(2)=" LD0"
      HD(3)=" LD1"
      HD(4)=" LD2"
      DO 250 ST = 1,STD
      DO 250 ROW = 1,N+4
      DO 250 COL = 1,20
      IF(STR(ROW+(N+4)*(ST-1),COL).EQ.999)STR(ROW+(N+4)*(ST-1
*                ),COL)=9999
250 CONTINUE
      WRITE(10,9100)K,V
      WRITE(10,9150)
      WRITE(10,9160)
      WRITE(10,9150)
      DO 300 J = 1,N
      WRITE(10,9200)J,BK(J),LD0(J),LD1(J),LD2(J)
300 CONTINUE
      WRITE(10,9150)
      WRITE(10,9250)
      DO 350 ST = 1,STD
      IF(STR(1+(N+4)*(ST-1),1),1).EQ.9999) GO TO 360
      WRITE(10,9300)
      DO 400 ROW = 1,4
      WRITE(10,9350)(HD(ROW,(STR(ROW+(N+4)*(ST-1),COL),COL=1,
*                20),HD(ROW))
      IF(ROW.EQ.1.OR.ROW.EQ.4) WRITE(10,9300)
400 CONTINUE
      DO 450 ROW = 1,N
      WRITE(10,9360)(STR(ROW+4+(N+4)*(ST-1),COL),COL=1,20),RO
*                W)
450 CONTINUE
      WRITE(10,9300)
      WRITE(10,9400)
```

```

      IF(STR(1+(N+4)*(ST-1),20).EQ.9999) GO TO 360
      PRINT/,"DESIGN IN PROCESS B=",STR(1+(N+4)*(ST-1),20)
350  CONTINUE
360  PRINT/,"JOB COMPLETED !"
      STOP
8000 FORMAT( (I1,X,I2,16X,I3) )
8050 FORMAT(4(X,I3) )
9050 FORMAT(20(X,I3) )
9100 FORMAT(10X,"K = ",I2/10X,"-----"/10X,"V = ",I2/10X,"--
*      ----"/10X,"GENERATOR DESIGNS"/)
9150 FORMAT(10X,"+-----+",4("----+"))
9160 FORMAT(10X,"IDSGNI B ILDOILD1ILD2I")
9200 FORMAT(10X,"ID",I3,"I",4(I3,"I"))
9250 FORMAT(/10X,"S T R U C T U R E"/10X,17("-")/)
9300 FORMAT(X,"+-----+",20("----+"),"-----+")
9350 FORMAT(X,"I",A4,"I",20(I3,"I"),A4,"I")
9360 FORMAT(X,"ID",I3,"I",20(I3,"I"),"D",I3,"I")
9400 FORMAT(/)
      END

```

**APPENDIX-G: ALL BINARY E MB GD TIB DESIGNS**  
 FOR  $k = 4$ ,  $v = 4$ ,  $b \leq 10$ .

K = 4

V = 4

A L L D E S I G N S

B	1	2	2	3	3	4	4	4	B
RO	0	0	4	0	4	0	12	4	RO
R	1	2	1	3	2	4	1	3	R
M 1	0	0	0	0	0	0	0	0	M 1
M 2	0	0	2	0	2	0	0	2	M 2
M 3	0	0	0	0	0	0	4	0	M 3
LD0	0	0	2	0	2	0	3	2	LD0
LD1	1	2	1	3	2	4	0	3	LD1
LD2	1	2	0	3	1	4	0	2	LD2
TRC	*****	*****	6.00	*****	3.67	*****	5.33	3.07	TRC
DSN	BTIB	BTIB	MBTIB	BTIB	MBTIB	BTIB	BTIB	MBTIB	DSN

B	4	4	4	5	5	5	5	5	B
RO	8	8	4	0	4	8	8	4	RO
R	2	2	3	5	4	3	3	4	R
M 1	0	2	4	0	0	0	3	4	M 1
M 2	4	0	0	0	2	4	1	0	M 2
M 3	0	2	0	0	0	0	1	0	M 3
LD0	4	3	3	0	2	4	4	3	LD0
LD1	0	1	2	5	4	1	1	3	LD1
LD2	1	1	2	5	3	2	2	3	LD2
TRC	2.83	3.05	2.42	*****	2.79	2.13	2.13	2.13	TRC
DSN	MBTIB	BTIB	BTIB	BTIB	MBTIB	MBTIB	MBTIB	BIB	DSN

B	6	6	6	6	6	6	6	6	B
RO R	0 6	4 5	16 2	8 4	12 3	8 4	12 3	8 4	RO R
M 1 M 2 M 3	0 0 0	0 2 0	0 2 4	0 4 0	0 6 0	2 0 2	2 2 2	3 1 1	M 1 M 2 M 3
LDO LD1 LD2	0 6 6	2 5 4	5 1 0	4 2 3	6 1 1	3 3 3	5 2 1	4 2 3	LDO LD1 LD2
TRC	*****	2.62	2.74	1.82	1.87	2.13	1.97	1.82	TRC
DSN	BTIB	MBTIB	MBTIB	MBTIB	BTIB	BTIB	MBTIB	MBTIB	DSN

B	6	6	7	7	7	7	7	7	B
RO R	4 5	8 4	0 7	4 6	8 5	12 4	8 5	12 4	RO R
M 1 M 2 M 3	4 0 0	4 2 0	0 0 0	0 2 0	0 4 0	0 6 0	3 1 1	3 3 1	M 1 M 2 M 3
LDO LD1 LD2	3 4 4	5 3 2	0 7 7	2 6 5	4 3 4	6 2 2	4 3 4	6 2 2	LDO LD1 LD2
TRC	1.96	1.64	*****	2.52	1.64	1.52	1.64	1.52	TRC
DSN	BTIB	MBTIB	BTIB	MBTIB	MBTIB	BTIB	MBTIB	BTIB	DSN

B	7	7	8	8	8	8	8	8	B
RO R	4 6	8 5	0 8	4 7	8 6	20 3	12 5	16 4	RO R
M 1 M 2 M 3	4 0 0	4 2 0	0 0 0	0 2 0	0 4 0	0 4 4	0 6 0	0 8 0	M 1 M 2 M 3
LDO LD1 LD2	3 5 5	5 4 3	0 8 8	2 7 6	4 4 5	7 0 1	6 3 3	8 2 1	LDO LD1 LD2
TRC	1.86	1.46	*****	2.44	1.53	1.82	1.33	1.40	TRC
DSN	BTIB	MBTIB	BTIB	MBTIB	MBTIB	MBTIB	BTIB	MBTIB	DSN

B	8	8	8	8	8	8	8	8	B
RO R	16 4	8 6	12 5	4 7	8 6	12 5	12 5	8 6	RO R
M 1 M 2 M 3	2 4 2	3 1 1	3 3 1	4 0 0	4 2 0	4 4 0	6 0 2	8 0 0	M 1 M 2 M 3
LD0 LD1 LD2	7 1 2	4 4 5	6 3 3	3 6 6	5 5 4	7 2 3	6 3 3	6 4 4	LD0 LD1 LD2
TRC	1.45	1.53	1.33	1.78	1.34	1.25	1.33	1.21	TRC
DSN	MBTIB	MBTIB	BTIB	BTIB	MBTIB	MBTIB	BTIB	BTIB	DSN

B	9	9	9	9	9	9	9	9	B
RO R	0 9	4 8	8 7	16 5	8 7	16 5	4 8	8 7	RO R
M 1 M 2 M 3	0 0 0	0 2 0	0 4 0	0 8 0	3 1 1	3 5 1	4 0 0	4 2 0	M 1 M 2 M 3
LD0 LD1 LD2	0 9 9	2 8 7	4 5 6	8 3 2	4 5 6	8 3 2	3 7 7	5 6 5	LD0 LD1 LD2
TRC	*****	2.38	1.45	1.19	1.45	1.19	1.72	1.26	TRC
DSN	BTIB	MBTIB	MBTIB	MBTIB	MBTIB	MBTIB	BTIB	MBTIB	DSN

B	9	9	9	10	10	10	10	10	B
RO R	12 6	12 6	8 7	0 10	4 9	8 8	24 4	16 6	RO R
M 1 M 2 M 3	4 4 0	7 1 1	8 0 0	0 0 0	0 2 0	0 4 0	0 6 4	0 8 0	M 1 M 2 M 3
LD0 LD1 LD2	7 3 4	7 3 4	6 5 5	0 10 10	2 9 8	4 6 7	9 1 1	8 4 3	LD0 LD1 LD2
TRC	1.13	1.13	1.13	*****	2.34	1.39	1.37	1.06	TRC
DSN	MBTIB	MBTIB	BTIB	BTIB	MBTIB	MBTIB	BTIB	MBTIB	DSN



B	10	10	10	10	10	10	10	10	B
RO R	20 5	20 5	8 8	16 6	4 9	8 8	12 7	16 6	RO R
M 1 M 2 M 3	0 10 0	2 6 2	3 1 1	3 5 1	4 0 0	4 2 0	4 4 0	4 6 0	M 1 M 2 M 3
LD0 LD1 LD2	10 1 2	9 2 2	4 6 7	8 4 3	3 8 8	5 7 6	7 4 5	9 3 3	LD0 LD1 LD2
TRC	1.12	1.15	1.39	1.06	1.68	1.20	1.04	1.02	TRC
DSN	MBTIB	BTIB	MBTIB	MBTIB	BTIB	MBTIB	MBTIB	BTIB	DSN

B	10	10	10	10	*****	*****	*****	*****	B
RO R	16 6	12 7	8 8	12 7	***** *****	***** *****	***** *****	***** *****	RO R
M 1 M 2 M 3	6 2 2	7 1 1	8 0 0	8 2 0	***** ***** *****	***** ***** *****	***** ***** *****	***** ***** *****	M 1 M 2 M 3
LD0 LD1 LD2	8 4 3	7 4 5	6 6 6	8 5 4	***** ***** *****	***** ***** *****	***** ***** *****	***** ***** *****	LD0 LD1 LD2
TRC	1.06	1.04	1.07	0.97	*****	*****	*****	*****	TRC
DSN	MBTIB	MBTIB	BIB	MBTIB					DSN

APPENDIX-H: OPTIMAL E MB GD TIB DESIGNS  
FOR  $k = 4, v = 4, b \leq 10$ .

K = 4

V = 4

OPTIMAL DESIGNS

DSG	ALL	BIB	BTIB	MBTIB	ALL	BIB	BTIB	MBTIB	DSG
B	1	*****	1	*****	2	*****	2	2	B
RO	0	*****	0	*****	4	*****	0	4	RO
R	1	*****	1	*****	1	*****	2	1	R
M 1	0	*****	0	*****	0	*****	0	0	M 1
M 2	0	*****	0	*****	2	*****	0	2	M 2
M 3	0	*****	0	*****	0	*****	0	0	M 3
LD0	0	*****	0	*****	2	*****	0	2	LD0
LD1	1	*****	1	*****	1	*****	2	1	LD1
LD2	1	*****	1	*****	0	*****	2	0	LD2
TRC	*****	*****	*****	*****	6.00	*****	*****	6.00	TRC
DSN	BTIB		BTIB		MBTIB		BTIB	MBTIB	DSN

DSG	ALL	BIB	BTIB	MBTIB	ALL	BIB	BTIB	MBTIB	DSG
B	3	*****	3	3	4	*****	4	4	B
RO	4	*****	0	4	4	*****	4	8	RO
R	2	*****	3	2	3	*****	3	2	R
M 1	0	*****	0	0	4	*****	4	0	M 1
M 2	2	*****	0	2	0	*****	0	4	M 2
M 3	0	*****	0	0	0	*****	0	0	M 3
LD0	2	*****	0	2	3	*****	3	4	LD0
LD1	2	*****	3	2	2	*****	2	0	LD1
LD2	1	*****	3	1	2	*****	2	1	LD2
TRC	3.67	*****	*****	3.67	2.42	*****	2.42	2.83	TRC
DSN	MBTIB		BTIB	MBTIB	BTIB		BTIB	MBTIB	DSN

DSG	ALL	BIB	BTIB	MBTIB	ALL	BIB	BTIB	MBTIB	DSG
B	5	5	5	5	6	*****	6	6	B
RO	8	4	0	8	8	*****	12	8	RO
R	3	4	5	3	4	*****	3	4	R
M 1	0	4	0	0	4	*****	0	4	M 1
M 2	4	0	0	4	2	*****	6	2	M 2
M 3	0	0	0	0	0	*****	0	0	M 3
LDO	4	3	0	4	5	*****	6	5	LDO
LD1	1	3	5	1	3	*****	1	3	LD1
LD2	2	3	5	2	2	*****	1	2	LD2
TRC	2.13	2.13	*****	2.13	1.64	*****	1.87	1.64	TRC
DSN	MBTIB	BIB	BTIB	MBTIB	MBTIB		BTIB	MBTIB	DSN

DSG	ALL	BIB	BTIB	MBTIB	ALL	BIB	BTIB	MBTIB	DSG
B	7	*****	7	7	8	*****	8	8	B
RO	8	*****	12	8	8	*****	8	12	RO
R	5	*****	4	5	6	*****	6	5	R
M 1	4	*****	0	4	8	*****	8	4	M 1
M 2	2	*****	6	2	0	*****	0	4	M 2
M 3	0	*****	0	0	0	*****	0	0	M 3
LDO	5	*****	6	5	6	*****	6	7	LDO
LD1	4	*****	2	4	4	*****	4	2	LD1
LD2	3	*****	2	3	4	*****	4	3	LD2
TRC	1.46	*****	1.52	1.46	1.21	*****	1.21	1.25	TRC
DSN	MBTIB		BTIB	MBTIB	BTIB		BTIB	MBTIB	DSN

DSG	ALL	BIB	BTIB	MBTIB	ALL	BIB	BTIB	MBTIB	DSG
B	9	*****	9	9	10	10	10	10	B
RO	12	*****	8	12	12	8	16	12	RO
R	6	*****	7	6	7	8	6	7	R
M 1	4	*****	8	4	8	8	4	8	M 1
M 2	4	*****	0	4	2	0	6	2	M 2
M 3	0	*****	0	0	0	0	0	0	M 3
LD0	7	*****	6	7	8	6	9	8	LD0
LD1	3	*****	5	3	5	6	3	5	LD1
LD2	4	*****	5	4	4	6	3	4	LD2
TRC	1.13	*****	1.13	1.13	0.97	1.07	1.02	0.97	TRC
DSN	MBTIB		BTIB	MBTIB	MBTIB	BIB	BTIB	MBTIB	DSN

**APPENDIX-I: STRUCTURE OF E MB GD TIB DESIGNS**  
**FOR  $k = 4, v = 4, b \leq 10.$**

K = 4

V = 4

GENERATOR DESIGNS

DSGN	B	LD0	LD1	LD2
D 1	1	0	1	1
D 2	2	2	1	0
D 3	4	4	0	1
D 4	4	3	2	2

S T R U C T U R E

B	1	2	2	3	3	4	4	4	4	4	4	5	B
LD0	0	0	2	0	2	0	3	2	4	3	3	0	LD0
LD1	1	2	1	3	2	4	0	3	0	1	2	5	LD1
LD2	1	2	0	3	1	4	0	2	1	1	2	5	LD2
D 1	1	2	***	3	1	4	***	2	***	***	***	5	D 1
D 2	***	***	1	***	1	***	***	1	***	***	***	***	D 2
D 3	***	***	***	***	***	***	***	***	1	***	***	***	D 3
D 4	***	***	***	***	***	***	***	***	***	***	1	***	D 4

B	5	5	5	5	6	6	6	6	6	6	6	6	B
LD0	2	4	4	3	0	2	5	4	6	3	5	4	LD0
LD1	4	1	1	3	6	5	1	2	1	3	2	2	LD1
LD2	3	2	2	3	6	4	0	3	1	3	1	3	LD2
D 1	3	1	1	1	6	4	***	2	***	***	***	2	D 1
D 2	1	***	***	***	***	1	***	***	1	***	***	***	D 2
D 3	***	1	1	***	***	***	***	1	1	***	***	1	D 3
D 4	***	***	***	1	***	***	***	***	***	***	***	***	D 4

B	6	6	7	7	7	7	7	7	7	7	8	8	B
LD0	3	5	0	2	4	6	4	6	3	5	0	2	LD0
LD1	4	3	7	6	3	2	3	2	5	4	8	7	LD1
LD2	4	2	7	5	4	2	4	2	5	3	8	6	LD2
D 1	2	***	7	5	3	1	3	1	3	1	8	6	D 1
D 2	***	1	***	1	***	1	***	1	***	1	***	1	D 2
D 3	***	***	***	***	1	1	1	1	***	***	***	***	D 3
D 4	1	1	***	***	***	***	***	***	1	1	***	***	D 4

B	8	8	8	8	8	8	8	8	8	8	8	8	B
LD0	4	7	6	8	7	4	6	3	5	7	6	6	LD0
LD1	4	0	3	2	1	4	3	6	5	2	3	4	LD1
LD2	5	1	3	1	2	5	3	6	4	3	3	4	LD2
D 1	4	***	2	***	***	4	2	4	2	***	2	***	D 1
D 2	***	***	1	2	***	***	1	***	1	***	1	***	D 2
D 3	1	***	1	1	***	1	1	***	***	1	1	***	D 3
D 4	***	***	***	***	***	***	***	1	1	1	***	2	D 4

B	9	9	9	9	9	9	9	9	9	9	9	10	B
LD0	0	2	4	8	4	8	3	5	7	7	6	0	LD0
LD1	9	8	5	3	5	3	7	6	3	3	5	10	LD1
LD2	9	7	6	2	6	2	7	5	4	4	5	10	LD2
D 1	9	7	5	1	5	1	5	3	1	1	1	10	D 1
D 2	***	1	***	2	***	2	***	1	***	***	***	***	D 2
D 3	***	***	1	1	1	1	***	***	1	1	***	***	D 3
D 4	***	***	***	***	***	***	1	1	1	1	2	***	D 4

B	10	10	10	10	10	10	10	10	10	10	10	10	B
LD0	2	4	9	8	10	9	4	8	3	5	7	9	LD0
LD1	9	6	1	4	1	2	6	4	8	7	4	3	LD1
LD2	8	7	1	3	2	2	7	3	8	6	5	3	LD2
D 1	8	6	***	2	***	***	6	2	6	4	2	***	D 1
D 2	1	***	***	2	1	***	***	2	***	1	***	1	D 2
D 3	***	1	***	1	2	***	1	1	***	***	1	1	D 3
D 4	***	***	***	***	***	***	***	***	1	1	1	1	D 4

B	10	10	10	10	***	***	***	***	***	***	***	***	B
LD0	8	7	6	8	***	***	***	***	***	***	***	***	LD0
LD1	4	4	6	5	***	***	***	***	***	***	***	***	LD1
LD2	3	5	6	4	***	***	***	***	***	***	***	***	LD2
D 1	2	2	2	***	***	***	***	***	***	***	***	***	D 1
D 2	2	***	***	1	***	***	***	***	***	***	***	***	D 2
D 3	1	1	***	***	***	***	***	***	***	***	***	***	D 3
D 4	***	1	2	2	***	***	***	***	***	***	***	***	D 4

NOTE:The designs which are not equivalent to any combination of generator designs are actually inadmissible with respect to some design, in the table, with the same number of blocks.

V I T A

The author was born in India in 1951 and latter emigrated with his family to Pakistan. He received the Master's Degree in Statistics in 1972 from University of Karachi and joined the same University as a Lecturer. In 1974 he joined a newly growing University in a remote area of the country. He was married in 1977 and has a son and two daughters. In 1983, he got a chance to join Middle East Technical University as a Ph.D. student and subsequently as a staff member also.

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