A MIXED INTEGER SECOND ORDER CONE PROGRAMMING REFORMULATION FOR A CONGESTED LOCATION AND CAPACITY ALLOCATION PROBLEM ON A SUPPLY CHAIN NETWORK

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ABSTRACT

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Supply chain network design involves location decisions for production facilities and distribution centers. We consider a make-to-order supply chain environment where distribution centers serve as crossdocking terminals. Long waiting times may occur at a cross-docking terminal, unless sufficient handling capacity is installed. In this study, we deal with a facility location problem with congestion effects at distribution centers. Along with location decisions, we make capacity allocation (service rate) and demand allocation decisions so that the total cost, including facility opening, transportation and congestion costs, is minimized.

Response time to customer orders is a critical performance measure for a supply chain network. The decisions like where the plants and distribution centers are located affect the response time of the system. Response time is more sensitive to these decisions in a make-to-order business environment. In a distribution network where distribution centers function as cross-docking terminals, capacity or the service rate decisions also affect the response time performance.

This study is closely related to a recent work Vidyarthi et al. (2009) which models distribution centers as M/G/1 queuing systems. They use the average waiting time formula of M/G/1 queuing model. Thus, the average waiting time at a distribution center is a nonlinear function of the demand rate allocated to and the service rate available at the distribution center. The authors Vidyarthi et al. (2009) propose a linear approximation approach and a Lagrangian based heuristic for the problem.

Different than the solution approach proposed in Vidyarthi et al. (2009), we propose a closed form formulation for the problem. In particular, we show that the waiting time function derived from M/G/1 queuing model can be represented via second order conic inequalities. Then, the problem becomes a mixed integer second order cone programming problem which can be solved by using commercial branch-and-bound software such as IBM ILOG CPLEX. Our computational tests show that proposed

reformulation can be solved in reasonable CPU times for practical size instances.

Keywords: Congestion, Cross docking, Second order conic programming

TEDARİK ZİNCİRİ AĞ TASARIMINDA BİR SIKIŞIK YER BELİRLEME VE KAPASİTE ATAMA PROBLEMİNİN KARIŞIK TAMSAYILI İKİNCİ DERECE KONİK PROGRAMLAMA İLE YENİDEN FORMÜLASYONU

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Tedarik zinciri ağ tasarım problemleri üretim ve dağıtım merkezlerinin yerlerinin belirlenmesi kararlarını içerir. Bu tezde siparişe üretim yapan bir tedarik zinciri ve çapraz yükleme işlevi gören dağıtım merkezlerinin bulunduğu bir problemi çözmeyi amaçlıyoruz. Yeterli yükleme indirme kapasitesi olmayan bir dağıtım merkezinde uzun bekleme süreleri oluşabilir. Bu çalışmada ele alınan yer belirleme problemi dağıtım merkezlerindeki bekleme sürelerini de dikkate alıyor. Çalışmada yer belirleme kararları kapasite ve talep atama kararları ile birlikte verirken tesis açma, taşıma ve dağıtım merkezlerindeki sıkışıklık maliyetlerinin toplamı minimize edilmeye çalışılıyor. Tedarik zinciri ağlarında siparişe yanıt süresi de önemli bir performans ölçüsüdür. Siparişe yanıt süresi üretim ve dağıtım tesisleri yer belirleme kararlarından etkilenir. Siperişe üretim yapan sistemlerde yanıt süresi bu kararlardan daha çok etkilenir. Ele alınan tipte dağıtım ağlarında dağıtım merkezinin kapasitesi ve işleme hızı da yanıt süresini etkiler. Bu tezde ele alınan problem Vidyarthi ve arkadaşları(2009) tarafından yapılan çalışmaya oldukça yakındır. O çalışmada dağıtım merkezleri M/G/1 kuyruk sistemleri olarak modellenmiştir. Yani, bir dağıtım merkezinde siparişlerin ortalama bekleme süreleri merkezin işleme hızı ve merkeze atanan talebin doğrusal olmayan bir fonksiyonu olarak modellenmektedir. Vidyarthi ve arkadaşları(2009) bu problemedoğrusal yaklaşıklama ve Lagrange temelli sezgisel algoritmalar önermişlerdir. Bu tezde Vidyarthi ve arkadaşlarından (2009) farklı olarak probleme kesin çözüm öneren bir formülasyon önerilmektedir. Dağıtım merkezlerinde M/G/1 kuyruk modelinin getirdiği toplam bekleme süresi fonksiyonunun ikinci derece konik programlama kısıtlarıyla ifade edilebildiği gösterilmiştir. Böylece çözülen problemin karışık tamsayılı ikinci derece konik programlama problemi olarak modellenebildiği ve IBM ILOG CPLEX gibi ticari dal-sınır yazılım paketleriyle çözülebilir olduğu gösterilmiştir. Yapılan hesaplamalı deneylerde gerçekçi boyutlarda problem örneklerinin makul sürelerde çözülebildiği gösterilmiştir.

Anahtar Kelimeler: Sıkışıklık, Çapraz yü kleme, İkinci derece konik programlama

This thesis is dedicated to my parents and who have supported me all the way since the beginning of my studies. Also, this thesis is dedicated to my Iranian friends who has been a great source of motivation and inspiration. Finally, this thesis is dedicated to all those who believe in the richness of learning.

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CHAPTER 1

INTRODUCTION

In this thesis, we consider a supply chain network design problem with congested distribution centers. We consider plant and distribution center location decisions along with capacity and demand allocations.

A supply chain includes all flows and transformations from the initial raw materials to the purchase of finished-items by the users. Each node of a supply chain network perform some activities such as manufacturing, product assembly or sales. These activities, however, necessitate logistical support, e.g., storage of intermediate or end goods, consolidation of orders for each consumer, and transportation.

Make-to-order (MTO) supply chain system is a business strategy which is applied in the cases of high product variety, variable customer demand, perishable products or obsolescence. Regarding to the strategic importance of response time in global business environment, MTO is a production process in which manufacturing resumes only after a customer's order is received. Due to extensive customization and competition, many firms adopt an MTO strategy to offer wide range of variety contrary to make-to-stock (MTS) supply chains in which the customer orders are met from stocks of finished products. The disadvantage associated with holding inventory of finished products may outweigh the advantage, particularly when we deal with products which become obsolete as technology advances or fashion changes. In addition, for many reasons, product and technology life cycles are getting shorter. Competitive market require more frequent product changes or innovation and consumers demand a greater variety of products than ever before.

Two critical decisions in supply chain network design are the location decisions for production plants and distribution centers. Transportation and inventory decisions are periodic and short term decisions which can be changed in response to external changes like demand variability. On the other hand, location decisions have long-term effects. Inefficient location choices made for plants and distribution centers can cause considerable surplus costs even the other factors like transportation or quantities of production are optimal. However, these long term decisions are subject to demand uncertainty at the time these decisions must be made. The collection of uncertainty in demand or capacity decisions and demand allocations, if not made carefully, may cause congestion in system or shortage in inventory which in turn make location decisions more critical.

The other challenge in supply chain management is the response time. Sule (2009) has performed a survey which supports the hypothesis that among the quality parameters of logistics the time is the most important one, even overtaking the price. A critical step to achieve targeted response time is the design of distribution network. The distribution network design addresses where to locate distribution centers, how much capacity to install and which demand points to be served from each distribution center. Distribution centers (DC) play major role in distribution networks. In fact, DC is a specific type of a warehouse. Frazelle (2002) refers DCs as distribution warehouses and defines them as facilities that accumulate and consolidate products from various manufacturing plants within a single firm, or from several firms, for combined shipment (economies of scale) to common customers. They

perform valuable functions which support the movement of materials. Storing goods (temporarily or longer), processing products, de-aggregating vehicle loads, creating SKU assortments, and assembling shipments are all activities commonly performed in these facilities. The main classification of DCs are: make-bulk/break-bulk consolidation terminal, a cross-docking center, a transshipment node, an assembly facility, a product fulfillment center or a returned goods depot.

Cross-docking (CD) is a strategy which appeared to cut the time items spend in the supply chain and reduce transportation cost. CD is a logistics technique applied in the retail and trucking industries to quickly consolidate shipments from separate sources and realize economies of scale in outbound transportation. There are three methods of CD (Burt, 2000):

- Manufacturing cross-docking: finished goods transfered from production line to a waiting truck or items produced are staged for later loading, are the categories.
- Distribution center cross-docking: consolidate inbound products from different suppliers which can be delivered when the last inbound shipment is received
- Terminal cross-docking: Products from DCs are dispatched to a break-bulk terminal for shipment of mixed loads to customers.

CD essentially eliminates the expensive inventory-holding costs of a warehouse, while still allowing short and temporary holding for consolidation and shipping functions. The idea is to transfer shipments directly from incoming (large scale) truck trailers to outgoing (small scale) truck trailers, without storage in between. Figure 1.1 shows schematic of CD centers (Yang et al. (2010)). With the process of moving shipments from the receiving dock (strip door) to the shipping dock (stack door), goods typically spend less than 24 hours in a cross-dock, sometimes even less than an hour.

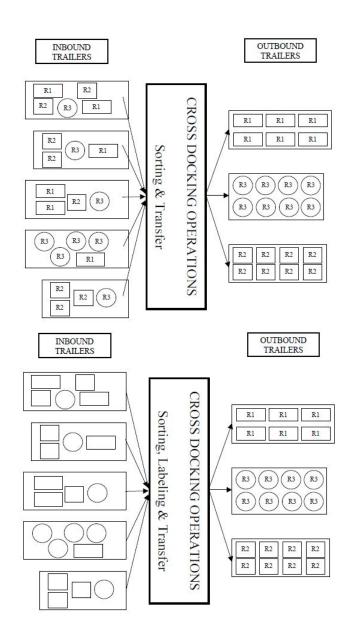


Figure 1.1: Cross docking center

The main advantages of utilizing CD are listed below:

- Elimination of activities associated with storage of products, such as incoming inspection, putaway, storage, pick-location replenishment, and order picking.
- Faster product flow and improved customer service.
- Reduced product handling.
- Cuts in inventory.
- Lower costs due to elimination of the above-mentioned activities.

But, in practice, both plant and CD centers face congestion for various reasons. Congestion in a distribution center may have several reasons. Below we list some of them which are observed especially in CD centers:

- Interference among forklifts: When a forklift makes a delivery to a stack door it must turn and maneuver its way in. Since loads are frequently bulky and hard to manipulate or carry, and there is usually freight sitting in the center of the dock, the forklift blocks each other trying to pass by that stack door. This phenomenon, usually referred as so-called interference, is most noticeable on docks that are operating close to capacity.
- **Dragline congestion:** a worker interacts with the dragline by pulling empty carts off the line and placing full carts on the line. Depending on the number of full and empty carts passing his door, he has to wait during either of these operations.
- **Congested floor space:** Sometimes workers, because of shipment consolidation cannot load a shipment directly into a stack door, but must park it temporarily on the floor nearby. Undoubtedly the existence of congestion in system increases response time which results higher inventory cost and lower customer attraction. Innumerous papers have been published in hierarchical congestion, hub congestion, emergency service congestion, distribution network congestion aiming to control congestion effect. In this thesis we consider the effect (cost) of distribution network congestion in MTO supply chain.

The congestion in a CD center can be controlled by either changing rate of service (capacity decision) or the amount of demand allocated. In fact, in congestible systems, there is the tradeoff between rate of service and demand allocation.

1.1 Contribution

In this thesis, we design a location and allocation model for supply chain network design. The decisions to be made are locations of plants, locations and capacities of CD centers and the assignment of customer demand to plants and CD centers. The objective is to minimize the total cost, which includes facility opening costs, transportation costs, and congestion costs. We assume that a customer's demand can be assigned to multiple plants and a single CD center. The CD centers are modeled as spatially distributed queues with Poisson arrivals. We assume the CD centers function as M/G/1 servers (to capture the dynamics of the response time). The model is a mixed integer nonlinear programming problem and presence of the congestion function makes it hard to solve.

We show that the problem can be reformulated as a Mixed Integer Second Order Cone Programming problem (MISOCP). This, in contrary to previous studies which propose linear approximation and heuristic algorithms (Vidyarthi et al. (2009) and Huang et al. (2002)), allows us to solve the problem using a commercial software which can solve subproblems using Second Order Conic Programming (SOCP) algorithm. To the best of our knowledge, this is the first exact approach for the plants and distribution center location with congested distribution centers.

The paper is organized as follows. In chapter 2, we present literature review. In Chapter 3, we present the mathematical model for the problem and analyze M/G/1 waiting time function and its characteristics. Then we develop a SOCP reformulation for the congestion (total waiting time) function. Finally, in Chapter 4, we give the computational results.

CHAPTER 2

LITERATURE REVIEW

In this thesis, we study a supply chain network design problem with facilities in high level, CD centers in middle level and customers in downstream. The objective is to minimize the total cost including fixed cost, transportation cost and congestion cost by finding the best locations for facilities and CD centers with appropriate handling capacities and assigning customer demand to facilities and CD centers. In this section, we give literature review, we consider five major groups of work: supply chain network design, DC location problem, structure of CD center, congestion (waiting time) and SOCP.

2.1 Supply Chain Design and Hierarchical System Design

The supply chain design problems have received significant attention from different aspects. The range of literature in this area is widespread but because location decisions play a critical role in the strategic design of a supply chain network most of the studies pertain location and allocation decisions. Melo et al. (2009) provide literature review of facility location models associated with supply chain management.

Olivares Benitez et al. (2010) introduce a bi-objective optimization problem for two echelon supply chain system where cost and time related objectives are considered. The most important feature added to the problem is to consider transportation mode as a decision to be made. Each mode represents a specific type of service with different costs. In our problem, transportation mode is not a decision and moreover, we consider waiting time cost.

One important aspect of supply chain design problems is the existence of different layers or echelons in the system, In other words, there exists a hierarchy in the system therefore we can say every multi layer supply chain is a hierarchical system. In our study, we have two interacting layers, facilities and CD centers. Each has a specified role in the system. Sahin and Sural (2007) classify the hierarchical facility location problems according to the features of the systems studied. They group them according to the flow pattern considered, service availability at each level, and spatial configuration of services. Moreover they investigate the applications, MIP models, and solution methods presented for the problem. Jayaraman et al. (2003) propose an integer programming model that would solve a comprehensive hierarchical problem to locate service facilities. The objective of the model is maximizing demand coverage while number of facilities is given and allocates different levels of service to the open facilities and also discuss some of the contributions to the current state-of-the-art in design of distribution systems.

Similar to above models, we define a hierarchical system which consists of facilities as upstream level

and CD centers as downstream level. We assign each customer to multiple plants, but all demand goes to customers through a single CD center. We, like most of the other studies, consider probabilistic demand in system which is more realistic. Demand is assumed to follow Poisson distribution.

2.2 Distribution center location problem

DC's are the foundation of a supply network. They, as middle layer of supply chain, has considerable effect on total logistic costs and timely service. Therefore DC related problems especially location problems have received special interest. A handful of publications address models and solution approaches for DC location problems. Higginson and Bookbinder (2005) explain the DC applications and different roles they can play in a supply chain.

Nozick and Turnquist (2001) integrated inventory and location decisions for DCs. They developed a model to find optimal location of DCs and to determine the minimum inventory level necessary to ensure a specified stockout probability for a given product. Their target is to retain service level by stocking optimal amount of safety stock in the best location.

Klose and Drexl (2005) illustrated different models for single stage capacitated and uncapacitated facility location problems (CFLP and UFLP) and then extended to two echelons cases (plant and depot) where fixed costs of both levels and transportation cost are considered. In addition, they discussed dynamic model of UFLP in which a given planning horizon is divided into two periods and all the fixed costs changed accordingly. Two echelons (plants and DC) case is the scenario that we consider in this thesis. The facilities have known capacity and CD centers have known service rates.

Our model deals with distribution network design problem for a two stage single product supply chain model similar to what Sourirajan et al. (2007) present. Their model seeks for DC locations such that lead time, including make up, replenishment and congestion times, is minimized and risk pooling benefits are maximized. They propose a Lagrangian heuristic. For congestion in DC, they assume M/M/1 queuing system whereas we consider M/G/1 system.

Aghezzaf (2005) presents two MIP models, one for the deterministic case and another for the robust optimization case. He considers strategic capacity planning and warehouse location problem. He integrates the issue of capacity expansion with distribution location. To incorporate demand uncertainty in the capacity expansion and warehouse location plans he utilized the concept of robust optimization then proposed a Lagrangian decomposition method. Similarly, our problem considers integration of capacity in CD centers locations.

In most of the aforementioned work, DCs behave as warehouse in which commodities are temporarily stored. The companies prefer to hold products as close as possible to customers to be able to satisfy demand faster and remain in competition. The DCs are also considered as assembly facilities to cover more customers taste. We consider DC for cross docking. In CD centers items are sorted and sent but not stored so there is no inventory cost considerations.

2.3 CD center Location Problem

A DC can be called a warehouse, a fulfillment center, a cross-dock facility, a bulk break center, and a package handling center. The name by which the distribution center is known is commonly based on the purpose of the operation. In this research area, existing work usually focus on two aspects: operational issues and the importance of the CD technique.

Galbreth et al. (2008) describe a multi-echelon supply chain in which both direct shipments and CD centers are available to move products from the manufacturer to customer locations. In their model, products can be shipped by truckload from a single supplier to a CD center. They propose a model for a single supplier and multiple CD centers supply chain. To show total costs saved by having the CD centers they compare the total costs of two supply chains, one with CD facilities and one without. The supply chain structure in this paper is similar to our design. Differently, we assume stochastic demand and in addition they consider fixed and transportation costs only whereas we also consider the congestion effect.

CD are sometimes called just in time distribution because most shipments spend less than 24 hours in CD centers. Napolitano (2000) and Bookbinder and Locke (1986), Shuib and Fatthi (2012), Vogt (2010) considered CD centers in supply chain management and propose some methods to design them optimally to improve operation. Yang et al. (2010) investigate the effects of various factors on the operation of CD centers. Computer simulation was done to monitor the travel and congestion time. The most comprehensive paper about congestion in CD center is Bartholdi and Gue (1999). They propose different congestion function for different parts of CD centers.

Despite the appealing affect of CD centers in supply chain network, there are few literature which consider CD centers in this area. These papers only cover CD centers location whereas our study consider location and capacity of them simultaneously and balance flows among them by adding congestion cost.

2.4 Waiting Time and Congestion

The waiting time is a result of congestion in a system. It has negative effect on customer's utility and hence, on a company's demand. There are different classes of literature which consider congestion in different ways. For instance, in demand capturing problems, customers are distance and time (congestion) sensitive. The relevant motivations are health and emergency services, banking or ticket selling centers. Usually, waiting time costs are calculated either by adding service level or queuing theory functions.

The studies in the literature, which consider service level, usually use Hillier and Liberman (1986) probability function of number of customers in system. They give different elasticity coefficients to distance and time to define specific levels of service. Marianov and Serra (2002) incorporate service level constraints to their model. The problem is to trade off between investment, operating cost and service quality. They chose heuristic concentration method (HC) for large instances. Sliva and Serra

(2007) consider a new version of demand capturing problem which not only takes into account the effect of traveling time but the waiting time on the market share. They propose metaheuristics which offer accurate results within acceptable computing time. They solved problem by ant colony optimization approach and because of time limitation of the methodology they used concentration algorithm for larger problems.

Marianov (2003) formulate a model for locating multiple server, congestible facilities. They control congestion in system by demand equilibrium constraint which is a nonlinear function of demand rate. They utilize traditional Lagrangean relaxation plus an iterative procedure. In Marianov et al. (2005), the goal is to maximize the number of people who travel to centers and stay in line until inoculated. They consider M/M/s/K queuing system and use the same demand equilibrium constraint as Marianov (2003). Obviously, the usage of Hillier and Liberman (1986) probability function of number in system increase computation times dramatically.

Marianov et al. (2008) present a model to maximize market share captured by the entering companies. They assume facilities function as M/M/m/K and customers are sensitive to distance and waiting time. To find shares in market they use logit function of cost which consists of convex combination of travel time and waiting time. Because of the probabilistic expression that model the number of customers captured (Hillier and Liberman, 1986), most of the constraints are nonlinear and complicated. They suggest ad hoc heuristics to solve the problem. All of these studies model congestion in facility location problem.

Aboolian et al. (2008) present the problem of locating facilities and allocation of servers on a congested network in order to reduce the costs of fixed installation, variable server, travel time and waiting time in the facilities act as M/M/K server. They proposed two heuristic approaches, descent approach and simulated annealing. In our model, allocation of service units is based on the service rate μ whereas they assume they can control the number of servers.

Berman and Drezner (2006) investigate the problem of locating a given number of facilities which can serve no more than a prespecified number of users at the same time. The goal is to maximize the number of customers captured. Unlike the other papers which use continuous variables to show proportion of arrival demand to each facility, they suggest following new form of waiting time function:

$$w(x) = \frac{1}{\mu - \sum_{i=1}^{n} \lambda \exp(-\alpha d(x, i))} \le UB$$

which is convex. They reformulate problem as capacitated facility location problem without fixed charge and compare exact (Cplex), Ascent algorithm, Simulated annealing and Tabu search approaches.

Castillo et al. (2009) consider two capacity choice scenarios, choosing a service rate for the servers and choosing the number of servers for the optimal location of facilities (M/M/s) which influence both the travel time cost and the waiting time of customers. The policy which they followed is replacing congestion term in the objective function by a simpler one and then applying a Lagrangean heuristic. A unique feature of their model is to deal with the social optimum rather than a user equilibrium. In another scenario with multiple servers, they approximate the number of servers by a continuous variable.

Marianov and Serra (2011) propose a multi server model for fixed number of facilities. They use famous standard equations for an M/M/s/K queuing system (Hillier and Liberman, 1986). The objective is Min(Z(1), Z(2)) where Z(1) and Z(2) are travel and congestion cost respectively. The presented model is a combinatorial, nonlinear optimization problem. They suggest a metaheuristic, the Max-Min Ant system, to obtain an initial solution. Then they use tabu search to improve the initial solutions.

In hierarchical systems, congestion may occur in different levels. Therefore, in some models, service constraints are incorporated to keep service level at a desirable level. The goal of these models is to find the minimum number of servers and their location which will cover a given region with distance or acceptable waiting time. Marianov and Serra (2001) control service level by adding probabilistic constraints whereby the probability of a customer standing in a line with *b* other customer is at most α , finally they apply a bi-level heuristic approach.

Another important class of network which may encounter congestion is hub-and-spoke network. Elhedhli and Hu (2005) proposed a model for uncapacitated single assignment p-hub location problem. In their model instead of using waiting time expression, a nonlinear power-law cost function was incorporated. Linearization and then Lagrangean heuristic were applied. Marianov and Serra (2003) analyze the queue formed by airplanes waiting for landing. To control congestion they insert probabilistic equations which bound the probability of the event that more than b airplanes on queue. They give a two phase heuristic approach, first greedy adding heuristic to find the initial set of p locations and second a one-opt exchange heuristic.

The capacity investment problem is another research area which take into account congestion cost. Rajagopalana and Yub (2001) address internal congestion in factories. Machines were modeled as nodes which act as M/G/1 queues. They include a constraint to guarantee that lead time satisfy target service level with prespecified probability, similar to service level constraint proposed in above papers. Although the research topic is different, the idea in this paper is close to our study in the sense that both consider tradeoff between fixed cost and congestion cost.

One similar paper to our model is Huang et al. (2005). They model a specific type of distribution network in which flows between origin and destination node must pass through connections which are congestible. Although the framework is similar to the one we use (three echelons system with congestion in the middle echelon), their solution approach is different. They consider M/G/1 server and use mean service time and second moment of service in congestion terms and set them as variable to find capacity in each connection. The resulting model is not convex or concave and they apply heuristics (Outer approximation and Lagrangean approaches). Differently, we reformulate the congestion terms and write them in convex form and then cast them as SOCP inequalities.

Model in Vidyarthi et al. (2009) is similar to our work. The objective of this paper is to model MTO and ATO supply chain with congestible DCs. Our work differ in two aspects. First one is the supply chain definition: they consider designing an MTO system, where manufacturing plants produce the wide range of disassembled items then ship them to DCs and after receiving orders the items are assembled according to customers expectations. In our model, orders are received in facilities and finished products are sent to customers through CD centers. Second one is the solution approach: in MTO case they suggest linearization method for M/G/1 term in the objective function. This is an ap-

proximation approach and may require a large number of linear inequalities to add. For ATO case they give a Lagrangian heuristic approach.

Despite the difference in basic definitions, our formulation is roughly the same. We show that the congestion function of M/G/1 model used by Vidyarthi et al. (2009), can be represented by SOCP inequalities. Hence, we show that the problem can be solved to optimum using MISOCP solvers instead of using approximation and heuristic methods. In their work, for MTO system the computational results reveal that the cutting plane algorithm provides optimal solution up to moderate instances in reasonable cpu time and for ATO system, the heuristic solution is on average within 6% of its optimal.

2.5 Second order cone programming

SOCP is a relatively new area in optimization and classified in convex optimization problem. An extensive review on SOCP is by Alizadeh and Goldfarb (2003). They present an overview of the SOCP problem and show SOCP form for LP, QP, quadratically constrained QP, and other classes of optimization problems. Lobo et al. (1998) recast different families of problem as SOCP and describe an efficient primal dual interior-point method for solving SOCP.

Gürel (2011) considered a multi-commodity network flow problem. The problem involves tradeoff between the total congestion and the capacity expansion costs on a given network. He estimated congestion on an arc by a convex increasing power function of the flow on it, then formulated problem as MISOCP problem and solved using Cplex. Atamtürk et al. (2012) studied joint facility location and multi-commodity inventory management problems with stochastic demand in uncapacitated and capacitated facilities case. They show how to formulate these problems as MISOCP. Valid inequalities were added to strengthen the model and improve the computational results.

Günlük and Linderoth (2008) describe the convex hull of mixed integer set. They show that for many classes of problems, the convex hull can be expressed via conic quadratic constraints, and can be solved via SOCP. They illustrated their approach on quadratic facility location and network design with congestion. They also show that congestion function of M/M/1 queuing model can be represented via SOCP inequalities.

In this thesis, we show that a congestion function based on M/G/1 model can be represented as SOCP inequalities and practical size problems can be solved in reasonable CPU time

2.6 Summary

As observed, the congestion phenomenon was studied in different contexts. Some papers consider congestion as part of service level, some present it as cost term in the objective function. Congestion is usually modeled as a nonlinear function and usually hard to deal with in mathematical programming models.

Studies in the literature often focus on how to model congestion. They apply heuristic algorithms to

solve the problem, in particular, for M/G/1 case to the our best knowledge no exact solution is proposed so far. Contrary to previous studies, we focus on reformulating the waiting time expression for an M/G/1 model and present a new formulation. Gürel (2011), Günlük and Linderoth (2008) and Atamtürk et al. (2012) consider SOCP reformulation in different area.

CHAPTER 3

PROBLEM DEFINITION

In this thesis, we study a plant and DC location problem in a MTO supply chain system. Along with location decisions for plants and DCs (CD centers), we consider capacity and demand allocation decisions to minimize the overall cost of the system. We consider a system where the demand location and demand rates are given. We need to find where to place plants and DCs and the capacity level to install at each DC. The objective to minimize is the sum of costs of opening plants, capacity of DCs, transportation and the congestion.

In our model, we suppose following supply processes: Plants receive orders and after the production, consolidate with other orders and transmit them, in large cargoes, to DCs for sorting and distribution.

We assume demand is concentrated at cities and we consider each city as a demand point. City k has a demand rate λ_k and the demand follows a Poisson distribution. There are two different sets of potential locations for plants and DCs. Furthermore, we assume a supply chain system which behaves like a referral hierarchy system in which a user cannot go to higher level server unless a low level server refers them to it (Narula, 1985), in other words, customers cannot receive their orders directly from plants.

While making location decisions we also consider capacity levels and demand allocation decisions. We decide which customer (demand point) will be served from which plants through which DC. We model the system in the following way: for the plants, requests of service are the union of all the requests for orders of all customers in its assignment set. Hence, demand in a plant can be viewed as a stochastic process, equal to the sum of several Poisson processes. By superposition property, this stochastic process known to be a Poisson process. Our problem involves the decision of which customer will be served by which plants and DCs. Therefore, we include a decision variable z_{ijk} denoting the proportion of customer k (λ_k) assigned to plant i and distributed from CD center j. Thereby, total customer demand assigned to plant i is:

$$\sum_{j}\sum_{k}\lambda_{k}z_{ijk}$$

In order to determine the input intensity of the DCs (second echelon) and also flow distribution between the first and the second echelon, we first recall the equivalence property for M/M/1 and M/M/mqueuing systems (Larson and Odoni, 1981). According to this property, if the system has an infinite (or large enough) queue capacity, and an arrival process of intensity of $\sum_j \sum_k \lambda_k z_{ijk}$, under steady state conditions, the departure process is also a Poisson process with the same intensity. Since, the input rates to DCs are sum of several Poisson departure rates from plants, according to Poisson superposition property (only some of the event being counted where this selection is made at random), we can conclude that the arrival rate to the DC *j* is also a Poisson process. Consequently, for DC *j* the arrival rate is:

$$\sum_{i}\sum_{k}\lambda_{k}z_{ijk}$$

The customers' orders arriving at the DCs are met on a First Come First Serve (FCFS) basis. We assume that each DC operates as a single flexible capacity server with infinite buffers to accommodate customer orders waiting for service. Under all these mentioned assumptions the DCs are modeled as a M/G/1 with service rates proportional to their capacity levels. We define discrete levels for capacity decisions with different costs.

3.1 Mathematical model

In this section, we give the mathematical formulation of the problem. We start with the notation used in the remaining part of the thesis. Initially, we write the nonlinear congestion function (F(z, y)) in the general form to make model more understandable, and later we present the explicit form of it.

Indices and parameters:

i	: index for plants $i = 1, 2, \ldots, I$
j	: index for DCs $j = 1, 2, \dots, J$
k	: index for customers $k = 1, 2, \dots, K$
F_i	: fixed cost of opening plant at location <i>i</i>
f_{js}	: fixed cost of opening DC j and acquiring capacity level s (\$/period)
C _{jk}	: unit transportation cost of serving customer k from DC j (\$/unit)
λ_k	: mean demand rate for the product from customer k (units/period)
C_{ij}	: unit cost of sending the product from plant <i>i</i> to DC <i>j</i> (/unit)
t	: mean response time cost per unit time per customer (\$/period/customer)
μ_s	: mean service rate, if capacity level s is allocated to a DC
μ^F	: mean service rate of plant
$G_j(\vec{Z}_j, \vec{Y}_j)$: congestion function of DC <i>j</i>

Decision variables:

Z _{i jk}	= fraction of demand k produced in plant i and distributed from DC j
y_{js}	= 1 if DC j is opened and capacity level s is acquired, 0 otherwise
x_i	= 1 if plant i is opened, 0 otherwise
W _{jk}	= 1 if customer k is assigned to DC j , 0 otherwise
$ec{Z}_j$	$: \{z_{ijk}: \forall i, k\}$
$ec{Y}_j$	$: \{y_{js}: \forall s\}$

$$\mathbf{Min} : \sum_{i} F_{i} x_{i} + \sum_{j} \sum_{s} f_{js} y_{js} + \sum_{j} \sum_{k} c_{jk} \lambda_{k} w_{jk}$$
$$+ \sum_{i} \sum_{j} \sum_{k} C_{ij} \lambda_{k} z_{ijk} + t \sum_{j} G_{j}(\vec{Z}_{j}, \vec{Y}_{j})$$

s.t.
$$\sum_{i} z_{ijk} = w_{jk}$$
 $\forall j, k$ (3.1)

$$\sum_{i} \sum_{k} \lambda_{k} z_{ijk} \leq \sum_{s} \mu_{s} y_{js} \qquad \qquad \forall j \qquad (3.2)$$

$$\sum_{s} y_{js} \le 1 \qquad \qquad \forall j \tag{3.3}$$

$$\sum_{j} w_{jk} = 1 \qquad \qquad \forall k \tag{3.4}$$

$$\sum_{j} \sum_{k} \lambda_{k} z_{ijk} \le \mu^{F} x_{i} \qquad \qquad \forall i \qquad (3.5)$$

$$w_{jk}, x_i, y_{js} \in \{0, 1\}$$
 $z_{ijk} \in [0, 1]$

The first and the second terms in objective function are the fixed costs of opening plants and DCs, respectively. The third term is the transportation cost from DCs to customers and the fourth term is the transportation costs from plants to DCs. The last term represents the congestion cost in the system. The *t* value can vary from customer to customer, but we assume that, it is the same across customers.

Constraint (3.1) ensures that demand of customer k is produced by the plants and transported through a selected DC. Constraint set (3.2) retains the steady state for DCs. In stochastic systems, when the system is in the steady state, then the probabilities that various states will be repeated will remain constant. Constraints (3.3) ensure that for a DC at most one capacity level is selected. Constraints (3.4) guarantee that every customer is assigned to a DC. Constraints (3.5) give the capacity limitation for each opened plant. To reduce the congestion and transportation costs we would open more plants which imposes more fixed cost. The model without $\sum_j G_j(\vec{Z}_j, \vec{Y}_j)$ is a simple hierarchical location problem which is a linear mixed integer programming model. We next discuss the congestion cost $G_j(\vec{Z}_j, \vec{Y}_j)$ that we use in our model.

3.2 The congestion function $G_j(\vec{Z}_j, \vec{Y}_j)$

 $G_j(\vec{Z}_j, \vec{Y}_j)$ is a nonlinear function representing the expected total waiting time at DC j, assuming the service times at each DC follow a general distribution M/G/1. The following notations are used for each DC: $\mu_j = \sum_{s=1}^{S} \mu_s y_{js}$ is the mean service rate and $\sigma_j^2 = \sum_{s=1}^{S} \sigma_{js}^2 y_{js}$ is the variance of the service rate. Assume X is independent and identically distributed (i.i.d) random variables denoting service time of a customer, then $E[X] = \tau_j = 1/\mu_j$ represents mean service time. The $E[X^2]$ is the second moment of service time, ρ_j is utilization value ($\rho_j = \lambda_j/\mu_j$) and CV_j^2 is squared coefficient of variation of service time ($CV_j^2 = \sigma_j^2/\tau_j^2$).

Under steady state and FCFS conditions, the average waiting time (including the service time) at a DC *j* is given by:

$$E[w_j(M/G/1)] = \frac{\Lambda_j E[x^2]}{2(1 - \Lambda_j E[x])} + E[x]$$
(3.6)

Where $\Lambda_j = \sum_i \sum_k \lambda_k z_{ijk}$. One possible way to present waiting time in M/G/1 queue system is the above expression. In the following proposition we show the above formula is convex.

Proposition 1 $E[w_i(M/G/1)]$ is convex in z if steady state condition in constraint (3.2) holds.

Proof. The denominator is always positive if (3.2) holds, then the denominator is concave. Numerator is convex and from result in Bector (1968), we can conclude that function F is convex. But in our model capacity level is a decision to be made, i.e, $E[X^2]$ and E[X] are decision variables. As a result, the function is not convex anymore. To use the waiting time formula in (3.6), we utilize the Pollaczek - Khintchine (PK) formula:

$$E[w_j(M/G/1)] = \frac{\rho + \Lambda \mu \operatorname{Var}(S)}{2(\mu - \Lambda)} + \mu^{-1}$$

Now by using $CV_i^2 (= \sigma_i^2 \mu_i^2)$ (Vidyarthi et al. (2009)), we come up with the following formula:

$$E[w_j(M/G/1)] = \left(\frac{1+CV_j^2}{2}\right)\frac{\tau_j\rho_j}{1-\rho_j} + \tau_j = \left(\frac{1+CV_j^2}{2}\right)\frac{\Lambda_j}{\mu_j(\mu_j - \Lambda_j)} + \frac{1}{\mu_j}$$

To obtain the expected total waiting time in the entire system, $G_j(\vec{Z}_j, \vec{Y}_j)$, $E[w_j(M/G/1)]$ is multiplied by the total demand rate from plants arriving to DC *j* and all resulting terms are summed together:

$$G_{j}(\vec{Z}_{j},\vec{Y}_{j}) = 1/2 \sum_{j} \left[\left(1 + \sum_{s} CV_{js}^{2} y_{js} \right) \frac{\sum_{i} \sum_{k} \lambda_{k} z_{ijk}}{\sum_{s} \mu_{s} y_{js} - \sum_{i} \sum_{k} \lambda_{k} z_{ijk}} + \left(1 - \sum_{s} CV_{js}^{2} y_{js} \right) \frac{\sum_{i} \sum_{k} \lambda_{k} z_{ijk}}{\sum_{s} \mu_{s} y_{js}} \right]$$
(3.7)

In (3.7), we assume, there are different squared coefficient of variation of service time, CV, for every pair of (j,s), i.e. CV_{js} . The M/G/1 expected waiting term is still not convex because of variable y in the denominator. We introduce the decision variable z_{ijks} which is the fraction of demand k satisfied by plant i and distributed through DC j with capacity s. The congestion function in (3.7) becomes:

$$F(\vec{Z}_j) = 1/2 \sum_j \sum_s \left[(1 + CV_{js}^2) \frac{\sum_i \sum_k \lambda_k z_{ijks}}{\mu_s - \sum_i \sum_k \lambda_k z_{ijks}} + (1 - CV_{js}^2) \frac{\sum_i \sum_k \lambda_k z_{ijks}}{\mu_s} \right]$$
(3.8)

In the next proposition we prove (3.8) is a convex function and then mathematical model will be presented to include the new decision variables and the new objective function.

Proposition 2 Function $F(\vec{Z}_j)$ given as (3.8), is a convex function when $\mu_s > \sum_i \sum_k \lambda_k z_{ijks}$.

Proof. By introducing the new variable z_{ijks} in place of z_{ijk} and adding:

$$z_{ijks} \le y_{js} \quad \forall i, j, k, s$$

to the constraint sets, we don't need variable y_{js} anymore because if DC j is not opened then, $y_{js} = 0$, according to (3.2), and no flow passes through that center i.e. $\sum_i \sum_k \lambda_k z_{ijks} = 0$. Thus average waiting

time for that center will be 0. On the other hand, due to adding capacity index to variables z_{ijk} , each term in (3.8) represents congestion term in DC *j* with capacity setting *s*. As a result, nonconvex fractional term in (3.7) turn to general following form:

$$\frac{x}{c-x}$$
 (c is constant)

Where $x = \sum_i \sum_k \lambda_k z_{ijks}$. When c > x, the second derivative is positive and the function is convex. Now the model with new variables and constraints is as follows:

$$\mathbf{Min} : \sum_{i} F_{i}x_{i} + \sum_{j} \sum_{s} f_{js}y_{js} + \sum_{j} \sum_{k} c_{jk}\lambda_{k}w_{jk} + \sum_{i} \sum_{j} \sum_{k} \sum_{s} C_{ij}\lambda_{k}z_{ijks}$$
$$+ t/2 \sum_{j} \sum_{s} \left[(1 + CV_{js}^{2}) \frac{\sum_{i} \sum_{k} \lambda_{k}z_{ijks}}{\mu_{s} - \sum_{i} \sum_{k} \lambda_{k}z_{ijks}} + (1 - CV_{js}^{2}) \frac{\sum_{i} \sum_{k} \lambda_{k}z_{ijks}}{\mu_{s}} \right]$$

s.t.
$$\sum_{i} \sum_{s} z_{ijks} = w_{jk} \qquad \forall j,k \qquad (3.9)$$

 $\sum_{j} y_{js} \le 1 \qquad \qquad \forall j \qquad (3.10)$

$$\sum_{i} \sum_{s} \sum_{s} \lambda_{k} z_{ijks} \leq \sum_{s} \mu_{s} y_{js} \qquad \forall j \qquad (3.11)$$

$$z_{ijks} \le y_{js} \qquad \qquad \forall i, j, k, s \qquad (3.12)$$

$$\sum_{j} w_{jk} \le 1 \qquad \qquad \forall k \tag{3.13}$$

$$\sum_{j} \sum_{k} \sum_{s} \lambda_{k} z_{ijks} \le \mu^{F} x_{i} \qquad \qquad \forall i \qquad (3.14)$$

$$w_{jk}, x_i, y_{js} \in \{0, 1\}$$
 $z_{ijks} \in [0, 1]$

3.3 SOCP representation of model

Recent developments in SOCP and the available commercial software allow us to solve a variety of convex problems by using SOCP inequalities.

In the most simple form, a SOCP is a convex optimization problem of the form:

Min
$$f^T x$$

s.t $||A_i x + b_i|| \le c_i^T x + d_i$ $i = 1, ..., N$

Where $x \in \mathbb{R}^n$ is a variable of the problem, $f \in \mathbb{R}^n$ are scalars and $A_i \in \mathbb{R}^{(n_i-1)\times n}$, $b_i \in \mathbb{R}^{(n_i-1)}$, $c_i \in \mathbb{R}^n$ and $d_i \in \mathbb{R}$ are parameters. The norm in constraint is the standard Euclidean norm which is called SOCP constraint. Generally, the definition of standard second-order (convex) cone of dimension k is:

$$\xi_k = \left\{ \begin{bmatrix} u \\ t \end{bmatrix} | u \in \mathbb{R}^{k-1}, t \in \mathbb{R}, ||u|| \le t \right\}$$

A n-dimensional convex set C is SOCP representable if, possibly after introducing auxiliary variables, it can be represented by a number of SOCP constraints and also function f is SOCP representable if its epigraph{ $(x, t)|f(x) \le t$ } has a SOCP representation. In other words, if in the problem:

$$\begin{array}{ll} Min & f(x) \\ s.t & x \in C \end{array}$$

f and C are SOC representable, then convex optimization can be reformulated as an SOCP and solved via algorithm available in most commercial optimization software (more information in Alizadeh and Goldfarb (2003), Tal and Nemirovski (2001) and Lobo et al. (1998)). In our model, we have linear constraints which are a special case of SOCP (Alizadeh and Goldfarb, 2003). But, f(x) consists congestion cost terms.

In the next proposition, we explain how we can cast expected waiting time for M/G/1 system as SOCP.

Proposition 3 Expected total waiting time of M/G/1 queue as given in (3.8), is SOCP representable.

Proof. (3.8) includes linear and nonlinear terms. To prove SOCP representability we will consider the nonlinear parts. To this end, we define an auxiliary variable S_{js} where:

$$S_{js} \ge \frac{\sum_{i} \sum_{k} \lambda_{k} z_{ijks}}{\mu_{s} - \sum_{i} \sum_{k} \lambda_{k} z_{ijks}} \qquad \forall j, s$$
(3.15)

then the nonlinear term in the objective function will be replaced by S_{js} . The new total expected waiting time for entire system will be:

$$1/2 \sum_{j} \sum_{s} \left[(1 + CV_{js}^2) S_{js} + (1 - CV_{js}^2) \frac{\sum_{i} \sum_{k} \lambda_k z_{ijks}}{\mu_s} \right]$$

For simplicity, define $R_{js} = \sum_i \sum_k \lambda_k z_{ijks}$. We will also drop indices of decision variables. We first multiply both sides in (3.15) with μ :

$$S\mu(\mu - R) \ge \mu R$$

We add R^2 to both side:

$$R^{2} + S\mu(\mu - R) \ge R^{2} + \mu R$$

$$R^{2} \le S\mu(\mu - R) - R(\mu - R)$$

$$R^{2} \le (S\mu - R)(\mu - R)$$
(3.16)

The constraint 3.16, is a **hyperbolic constraint**(hyperbolic constraints are the constraints which describe half a hyperboloid) of the form:

$$w^2 \le xy, \quad x \ge 0, \quad y \ge 0, \Leftrightarrow \left\| \begin{bmatrix} 2w\\ x-y \end{bmatrix} \right\| \le x+y$$
 (3.17)

or in the matrix form :

$$w^T w \le xy, \quad x \ge 0, \quad y \ge 0, \Leftrightarrow \left\| \begin{bmatrix} 2w \\ x-y \end{bmatrix} \right\| \le x+y$$

The hyperbolic constraint, is a variety of conic SOCP sets (Alizadeh and Goldfarb (2003)). The SOCP will be accepted for solution by the optimizers if it can be transformed to the following convex SOC constraint:

$$-c_0 x_0^2 + \sum_i c_i x_i^2 \le 0$$

Therefore we need to define new variables and replace hyperbolic constraints by them. the following procedure is done for all (j, s). Again for simplicity, we drop indices of variables:

$$P1 = S\mu - R - \mu + R$$
$$P2 = S\mu - R + \mu - R$$
$$P3 = 2R$$
$$P3^{2} + P1^{2} \le P2^{2}$$

Hence, (3.8) is SOCP representable.1

At the end, SOCP of our location and allocation model will be:

$$\mathbf{Min} : \sum_{i} Fx_{i} + \sum_{j} \sum_{s} f_{js}y_{js} + \sum_{j} \sum_{k} c_{jk}\lambda_{k}w_{jk} + \sum_{i} \sum_{j} \sum_{k} \sum_{s} C_{ij}\lambda_{k}z_{ijks}$$
$$+ t/2 \sum_{j} \sum_{s} \left[(1 + CV_{jk}^{2})S_{js} + (1 - CV_{jk}^{2})\frac{\sum_{i} \sum_{k} \lambda_{k}z_{ijks}}{\mu_{s}} \right]$$

s.t. $\left(\sum_{i}\sum_{k}\lambda_{k}z_{ijks}\right)^{2} \leq \left(S_{js}\mu_{s} - \sum_{i}\sum_{k}\lambda_{k}z_{ijks}\right)\left(\mu_{s} - \sum_{i}\sum_{k}\lambda_{k}z_{ijks}\right)$ $\forall j, s$

(and) constraints (3.9)-(3.14)

 $w_{jk}, x_{ic}, y_{js} \in \{0, 1\}$ $z_{ijks} \in [0, 1]$ S_{js} Free

3.4 Summary

In this section we first presented a mathematical model for the problem. The model was a MINLP. We showed that it can represented as MISOCP.

What we did is, to represent each waiting time expression in the objective function by a set of new decision variables, and linear and SOCP inequalities. This help to solve the problem to exact optimum as the problem size permits via commercial MISOCP solvers such as Cplex. This allows to employ strong B&B, preprocessing and cut generation features of such solvers in solving our problem.

CHAPTER 4

COMPUTATIONAL STUDY

In this section, we report our computational results for the proposed MISOCP reformulation. We tested MISOCP formulation over several instances generated by using SGB128 data (City Distance Datasets, www.sc.fsu.edu) which describes 128 cities in North America. Appendix A.1 gives the map of the cities in the dataset. All instances were coded in MATLAB version R2012a to generate LP files and were solved by using ILOG CPLEX 12.5 on 8 GB RAM, 2.66 GHz computer.

We assume that the demand of each city is concentrated in the city center and hence consider each city as an individual customer. Demand rate (λ_k) at a city k is proportional to the population and obtained by dividing the population of each city by 1000 (The name and population of cities are given in Appendix A). We generate 50, 100 and 128 cities instances.

We also assume that plants and DCs are located close to city centers, so out of selected cities, We randomly choose two sets of candidate locations, for plants (N) and DCs (M). We assume the capacity of plants and DCs are both proportional to total demand,

Total demand rate =
$$\sum_{k} \lambda_{k}$$

The capacity levels of possible plants are given, and we assume the fixed cost of opening a plant is a function of its capacity. The capacity and the fixed cost of plants are obtained by:

$$\mu^F = U(0.1, 0.5) \times \text{total demand rate}$$
(4.1)

$$F = \beta \times U(1000, 2000) \times \sqrt{\mu^F} \tag{4.2}$$

Where β is a constant multiplier and it is used to explore the sensitivity of solutions to different levels of fixed cost. In the model, we assume the candidate locations for DCs, are different from the candidate location of plants, and as mentioned earlier, we choose them randomly. The capacity of DCs is a decision variable (y_{js}), and we assume, the available capacities are discrete and proportional to total population. We define three levels for capacities, according to the number of demand point. The capacities of DCs and the corresponding fixed cost are generated by the formulas (4.3) and (4.4), respectively,

$$\mu_s = \alpha_s \times \text{total demand} \tag{4.3}$$

$$f = \beta \times 100 \times \sqrt{\mu_s} \tag{4.4}$$

 α is dependent to number of customers where $\alpha_s = 0.15, 0.2, 0.45$ for $K = 50; \alpha_s = 0.1, 0.2, 0.3$ for K = 100 and $\alpha_s = 0.1, 0.15, 0.2, 0.3, 0.45$ for K = 128. To calculate transportation costs (c_{jk}, C_{ij}) , we use the distance matrix in SGB128 dataset. We divide distances between plants and DCs, and between DCs and customers by .01 and .04 respectively, to differentiate the types of transportation in network. These types of difference between transportation mode is consistent to our motivation, where DCs act

CD operation in supply chain network (one of the main incentive to use CD centers in supply chain is to reduce the transportation costs by altering vehicles type). We use a multiplier in calculation of transportation cost, δ , to explore the model sensitivity to this cost. To show the effects of the fixed and transportation costs, we set them in two levels (low and high), and to explore the effect of congestion cost accurately, we set it in three levels (low, mid and high).

The response waiting cost can be expressed in different cost functions such as piecewise linear function or exponential cost function. However, in practice, determining the value of average response time is difficult. In this computational study, we generate values of t by using the following formula:

$$t = \theta \times \left(\frac{\sum_{i} \sum_{j} \lambda_k c_{ij}}{I \times J}\right)$$
(4.5)

where θ is the response time cost coefficient. The high θ can be interpreted to indicator of a situation which losing customers due to high expected waiting time is costly. The numerator of fraction of expression in (4.5) is total transportation cost in system. For M/G/1 expected waiting time, the coefficient of variation (*CV*) is set to 1.5.

All experimental factors and their levels are listed in Table 4.1.

Table 4.1: Experimental factors

Parameters	Levels
K	50,100,128
N	10,20
M	10,20
δ	1,10
β	1,10
θ	.01,5,10

We have $3 \times 2 \times 2 \times 2 \times 3 = 144$ experimental settings and for each setting we have three replications. All results are given in Appendix B. We set two hours time limit for all runs.

Tables 4.2 - 4.4, show the summary of all experimental results. The column *NP* is the number of times which optimal solution was found for each instance in three replications. The column Gap(%) is the gap between the best integer objective and the objective of the best remaining node. The columns, FC(%), TC(%) and CC(%) represent the proportion of the fixed cost, transportation cost and congestion cost in total cost respectively.

When all costs are low, the total cost almost consists of fixed and transportation costs and congestion cost is negligible. As the θ increases, the congestion cost's proportion increases, and it is what we expect. The greatest value of congestion's proportion is when the congestion cost is mid or high and the other costs are low.

The results show that, the CPU time and gap are affected by the size of problem, the congestion cost and fixed cost. As the number of customers increase average CPU time and gap increase. Maximum number of non optimal solution is observed when fixed cost (β) and congestion cost (θ) are both high. One possible reason is, due to high congestion cost, we need to open more DC or increase the capacities to reduce total response time, but on the other hand, the fixed cost is also high. In most instances which the congestion cost is low, we obtain optimal solutions or very small gap, but when congestion cost is high, CPU is more likely to hit time limit with a significant gap. When the congestion cost is low and fixed cost is high, less plants and DCs are opened because waiting cost is negligible, and conversely when fixed costs are low and the other is high, it is profitable to open more plants or DCs with higher capacity (Table 4.8).

Table 4.2: K = 50

N	М	θ	δ	β	NP	Gap(%)	CPU(s)	FC(%)	TC(%)	CC(%)
			T	Low	3	0	70.67	64	35	0.6
		-	Low	High	3	0	77.5	94	5	0.1
		Low	x x · 1	Low	3	0	17.25	16	84	0.2
			High	High	3	0	43.84	64	36	0.1
			-	Low	3	0	141.09	53	31	16.4
	10		Low	High	3	0	557.16	90	5	4.6
	10	Mid	x x · 1	Low	3	0	16.74	15	78	7.2
			High	High	3	0	64.87	63	33	3.9
			T	Low	3	0	191.35	44	27	29.3
		11.1	Low	High	3	0	3868.08	88	5	7.1
		High	11.1	Low	3	0	27.02	17	71	11.1
10			High	High	3	0	146.09	60	34	6.1
10			T	Low	3	0	169.73	67	33	0.2
		T	Low	High	3	0	80.5	95	5	0.1
		Low	II: .1	Low	3	0	26.13	17	83	0.1
			High	High	3	0	141.93	66	34	0.1
			Larr	Low	3	0	459.98	55	28	16.4
	20	Mid	Low	High	2	2	7204.3	90	6	4.5
	20	wiid	II: -h	Low	3	0	24.52	16	78	5.8
			High	High	3	0	208.2	65	32	2.7
			Low	Low	3	0	358.54	48	26	26.2
		Uich	Low	High	0	5	7204.97	89	5	5.9
		High High	Uigh	Low	3	0	47.97	19	68	12.4
			High	High	3	0	219.81	62	32	6.3
			Low	Low	3	0	180.87	64	36	0.4
		Low	LOW	High	3	0	153.43	94	6	0.1
		LOW	High	Low	3	0	54.65	15	85	0.1
			Ingn	High	3	0	236.22	64	36	0.1
			Low	Low	3	0	412.81	54	30	16.8
	10	Mid	LOW	High	0	0	5337.46	90	5	4.5
	10	wita	High	Low	3	0	74.74	19	74	6.7
				High	3	0	1455.43	63	32	4.2
			Low	Low	3	0	2486.78	42	28	29.4
		High		High	3	3	7203.33	87	6	7.2
			High	Low	3	0	50.17	17	71	11.6
20			ingn	High	3	0	597.52	59	34	6.8
			Low	Low	3	0	676.78	70	30	0.3
		Low		High	3	0	1927.38	96	4	0.1
		W	High	Low	3	0	44.66	19	81	0.1
			- ingi	High	3	0	475.51	70	30	0.1
			Low	Low	3	0	2820.14	56	27	17.3
	20	Mid		High	0	3	7206.36	91	5	4.6
	20	1,110	High	Low	3	0	126.14	18	74	8.0
			- ngn	High	3	0	2525.98	66	30	4.1
			Low	Low	3	6	7205.87	47	26	26.7
		High		High	0	8	7204.3	87	6	6.2
			High	Low	3	0	121.12	26	61	12.9
				High	3	0	2885.66	64	28	7.4

Table 4.3: K = 100

N	М	θ	δ	β	NP	Gap(%)	CPU(s)	FC(%)	TC(%)	CC(%)		
			T	Low	3	0	178.51	54	46	0.3		
		Ŧ	Low	High	3	0	5567.21	92	8	0.1		
		Low	TT' 1	Low	3	0	35.83	11	89	0.2		
			High	High	3	0	147.95	54	46	0.1		
			Ŧ	Low	3	0	593.74	47	38	14.6		
	10	201	Low	High	0	0	4944.73	89	7	3.7		
	10	Mid	TT: 1	Low	3	0	120.12	11	83	6.0		
			High	High	3	0	351.14	55	42	3.3		
			T	Low	3	0	1625.61	39	36	25.6		
		TT: -h	Low	High	0	4	7200.85	86	8	5.7		
		High	II: -h	Low	3	0	105.29	16	77	7.1		
10			High	High	3	0	862.95	51	44	5.0		
10			т.	Low	3	0	398.47	57	43	0.4		
		T	Low	High	3	0	940.11	93	7	0.1		
		Low	II: -1-	Low	3	0	58.08	16	84	0.1		
			High	High	3	0	645.74	57	43	0.1		
			Low	Low	3	0	1562.6	50	35	14.4		
	20	ма	Low	High	0	10	7200.62	90	6	3.2		
	20	Mid	High	Low	3	0	105.89	12	83	4.4		
			High	High	3	0	360.96	57	40	2.3		
				Low	Low	1	0	3834.46	44	35	21.7	
		Uich	High	High	0	25	7203.83	84	9	7.0		
		підії	High	Low	3	0	193.38	18	74	7.9		
			Ingn	High	3	0	1295.7	54	41	5.0		
			Low	Low	3	0	709.51	56	43	0.3		
		Low	LOW	High	3	0	1692.83	93	7	0.1		
		LOW	High	Low	3	0	106	11	88	0.1		
			Ingn	High	3	0	554.27	56	44	0.1		
			Low	Low	3	0	7202.35	47	36	16.3		
	10	Mid	LOW	High	3	4	7200.2	88	7	5.0		
	10	with	High	Low	3	0	540.86	15	80	5.1		
			Ingi	High	3	0	3867.23	56	40	3.8		
			Low	Low	3	13	7203.41	38	36	26.2		
		High	LOW	High	0	6	7202.41	85	8	6.5		
		111511	High	Low	3	0	465.73	13	77	9.5		
20				High	3	0	2126.68	52	42	6.6		
			Low	Low	3	0	2733.62	61	39	0.2		
		Low	W	High	3	0	2959.67	94	6	0.1		
		2011	High	Low	3	0	145.02	19	81	0.0		
				High	3	0	4072.03	61	39	0.0		
			Low	Low	1	7	7202.97	50	35	15.4		
	20	Mid	2011	High	0	5	7202.86	89	7	4.5		
	20	1,110	High	Low	3	0	340.55	16	79	5.5		
				High	3	5	7203.69	58	39	3.4		
			Low	Low	0	16	7203.42	46	35	19.2		
		High		High	0	8	7204.44	87	7	5.9		
		-	High	High	High	Low	3	0	347.09	21	71	8.3
			Ingn	High	3	2	7201.33	56	37	7.1		

Table 4.4: K = 128

N	М	θ	δ	β	NP	Gap(%)	CPU(s)	FC(%)	TC(%)	CC(%)	
			T area	Low	3	0	308.94	51	49	0.3	
		T	Low	High	1	0	231.16	91	9	0.1	
		Low	TT: 1	Low	3	0	73.21	13	87	0.1	
			High	High	3	0	326.14	51	49	0.1	
			•	Low	0	0	655.47	50	42	8.3	
	10	201	Low	High	0	2	7204.14	90	8	2.2	
	10	Mid		Low	3	0	124.61	11	87	1.9	
			High	High	0	0	546.71	54	45	1.3	
			•	Low	0	0	913.87	40	41	19.0	
		*** 1	Low	High	0	0	2186.99	86	9	5.8	
		High		Low	3	0	60.92	12	83	5.5	
			High	High	1	0	601.93	49	47	3.8	
10			-	Low	3	0	1014.6	55	44	0.1	
		_	Low	High	3	0	7127.78	92	8	0.1	
		Low		Low	3	0	159.03	15	85	0.0	
			High	High	3	0	1078.65	55	45	0.0	
				Low	3	0	7100.18	52	40	7.3	
			Low	High	0	3	7202.46	89	8	2.7	
	20	Mid		Low	3	0	247.35	12	86	1.7	
			High	High	3	0	2273.22	55	43	1.5	
				Low	0	0	4423.13	47	38	14.9	
		High —	Low	High	0	25	7202.69	85	9	6.1	
			-	Low	3	0	412.93	14	81	4.3	
			High	High	2	0	2967.23	54	43	2.8	
					Low	3	0	4065.93	54	46	0.2
			Low	High	2	19	7201.97	92	8	0.1	
		Low		Low	3	0	427.99	11	89	0.1	
			High	High	3	0	1893.62	54	46	0.0	
				Low	0	7	7203.44	50	40	9.8	
			Low	High	0	11	7203.44	87	9	4.2	
	10	Mid		Low	3	0	1452.1	14	84	1.9	
			High	High	0	4	7203.02	54	44	1.9	
				Low	0	17	7201.85	42	41	17.4	
			Low	High	0	47	7201.85	58	8	34.3	
		High		Low	3	0	687.04	14	82	4.1	
			High	High	1	6	7202.28	52	45	3.0	
20				Low	0	4	7202.20	58	42	0.3	
			Low	High	0	1	7202.43	93	7	0.1	
		Low		Low	3	0	319.71	21	79	0.0	
			High	High	2	0	5219.28	58	42	0.0	
				Low	0	9	7213.88	51	39	9.1	
			Low	High	0	7	7213.88	87	10	3.4	
	20	Mid		Low	3	0	885.12	15	83	2.2	
			High	High	0	5	7207.81	56	43	1.4	
				Low	0	26	7207.81	40	43 39	20.7	
			Low		3	0		20	59 74	6.0	
		High		High Low	3 0	10	702.43	55	40	5.3	
			High				7202.08				
				High	0	26	7200.68	86	9	4.7	

The last row in Table 4.4 is result of a instance. We didn't get any integer solution within two hours for two other replications.

Table 4.5 shows the effect of number of candidate plant locations, DCs locations and customers on CPU time and gap. It gives statistic of the model's performance for all sets of plants, DCs and customers. The column *Opt* indicates how much percent of instances solved to optimum. The column *Gap* shows the average gap of non optimal solutions. As the size of problems get larger, as we expect, the number of optimal solutions decrease and CPU time increases. The model could solve around 87% of instances for k = 50, 73% for k = 100 and 64% for k = 128. Totally, 75% of instances were solved optimally. Regarding to two hours time limit, the results indicate that the model is able to solve the practical size instances of this problem.

N	М	Κ	Opt(%)	CPU(s)	Gap(%)
10	10	50	100	333.03	-
		100	83	936.86	3
		128	47	581.78	6
10	20	50	88	194.14	3
		100	88	1108.88	14
		128	72	2153.94	9
20	10	50	92	1019.65	2
		100	92	1259.80	7
		128	53	2034.85	10
20	20	50	83	1216.37	7
		100	69	1927.15	9
		128	30	1253.88	9

Table 4.5: The model performance

In the remaining parts of this section, we explore the effects of costs on solutions. To observe the sensitivity of solutions to transportation cost and fixed cost, we show the relation between this cost and the number of DCs to be opened in Table 4.6. The data are for N = 10, M = 10 and $\theta = .01$ instances from appendix B. The column *Op.DCs* indicates the average number of opened DCs for different number of customers. When the transportation cost is high ($\delta = 10$), on average, more DCs were opened to reduce transportation cost. When fixed cost is high ($\beta = 10$), on average, less DCs were opened. Next, we will explain the number of DCs are not just affected by these two parameters, but also it is affected by waiting time cost.

Table 4.6: Transportation cost and fixed cost versus CD centers

K	δ	β	Op. DCs		
	Low	Low	5		
50	Low	High	3		
30	High	Low	6		
	High	High	5		
	Low	Low	6		
100	Low	High	4		
100	High	Low	7		
	High	High	6		
	Low	Low	6		
Continued on next page					

K	δ	β	Op. DCs
	Low	High	3
	High	Low	7
	High	High	6

Table 4.6: Transportation cost and fixed cost versus CD centers

The waiting time cost or congestion cost actually is penalty of not fulfilling customer commitments in lead time and the value we assign to it. It can present customers delay sensitivity. In the case of high sensitivity, waiting time imposes too much cost and take grater percentage of total cost. In Table 4.7, we can observe different levels of congestion cost (different level of θ) and their waiting time and cost. The data are for N = 10, M = 10, K = 100. The *CC* column is the average congestion cost percentage in the total cost, and E(W) column represents the average total expected waiting time in the system. Table 4.7 shows that, as the congestion cost increases, the proportion of congestion cost in total cost increases. To mitigate the congestion cost, we open more DCs or increase capacities, consequently the total waiting time (E(W)) decreases. In Table 4.7, when the congestion cost is very low ($\theta = .01$), the E(W) is very high. The reason is, because of low congestion rate, we can open less DCs to reduce fixed cost which in turn, increases the waiting time in system.

Table 4.7: Effect of waiting time cost on total expected waiting time

θ	CC	E(W)
.01	.1	183.25
5	3.8	17.31
10	7.2	14.76

Another observation in experimental results, is the relation between congestion cost and capacity decision in DCs. To show this relation accurately, we solved extra instances with more levels. We tabulate the solutions for M = 10, N = 10, K = 50 with different levels of congestion cost in Table 4.8. In this table, NOP is the number of opened plants, NODC is the number of opened DCs and DC capacity column is the capacities of the opened DCs. The numbers in parenthesis indicate the capacities, where 1 corresponds to lowest and 3 corresponds to highest capacity. As θ increases, the DCs capacity increases. In fact, we can say, because the cost of upgrading the capacity is lower than opening new facility, the number of opened DCs doesn't change considerably as the congestion cost increases, but the capacities increases.

Table 4.8: Effect of waiting time cost on chosen capacity levels

θ	NOP	NODC	DCs capacity
0.01	9	7	1(2)4(2)5(1)7(2)8(2)9(2)10(3)
0.1	9	6	1(2)4(2)5(2)6(3)8(3)9(3)
1	9	7	1(3)4(3)5(2)7(3)8(3)9(3)10(3)
10	9	8	1(3)2(3)4(3)5(3)7(3)8(3)9(3)10(3)

Note that, because transportation and fixed cost of plants are constant, the number of opened plants remained unchanged. The capacity decision in our model has direct and indirect effects on total cost. According to (4.2), the fixed cost of establishing CD centers depends on capacities. On the other hand,

the capacity decision is a tool to control congestion cost (3.8). In fact, we can say capacity increment is a cheap alternative for establishing a new DC and it is what we observed in 4.8, where by increasing congestion cost, only highest level capacities were opened.

In the model, we assumed capacity levels are proportional to entire demand. We set the different sets of discrete capacities in accordance with the number of demand points. The selection of the levels of capacities is very important. The experimental results reveal that, any change in the levels of capacities cause considerable changes in the total expected waiting time in the system. To realize the effect of capacity levels, we did extra runs for M = 10, N = 10, K = 50. We defined three sets of levels, tight, moderate, loose (all other setting parameters remain unchanged). Table 4.9 shows the results. The values in the E(W) depends on the fixed cost of opening a DC and the cost of congestion. For instance, when fixed cost is low and the levels of capacities are tight, we can open more DCs to reduce E(W).

θ 0.01 1	E(W) 184.92 30.86
0.00	
1	30.86
200	13.36
0.01	131.11
1	10.29
200	5.89
0.01	27.97
1	8.24
200	5.39
	0.01 1 200 0.01 1

Table 4.9: Effect of capacity in E(W)

In the next chapter, we present the conclusion and bibliography.

CHAPTER 5

CONCLUSION AND FUTURE STUDY

In this thesis, we studied a congestible supply chain network design problem. We consider make-toorder supply chain which consists of plants, distribution centers and customers. Distribution centers act as Cross docking terminals. The resulting problem was MINLP and non convex. We showed the waiting time function derived from M/G/1 queuing model can be represented via second order conic inequalities. We proposed a closed form formulation for the problem. Then, the problem became a MISOCP which can be solved by using commercial branch-and-bound software such as IBM ILOG CPLEX.

The experiments showed that the model can solve practical size problems in reasonable CPU times. We implemented the model on 128 cities of U.S, and it could find optimal solution in 75% of instances. In the comparison to the approaches in the literature, this exact approach is easier to employ and there is no need to use any heuristics. The proposed SOCP formulation in this thesis, is not restricted to supply chain design problems, it is applicable in all M/G/1 queuing systems.

One possible extension of our model is considering congestion in both echelons, plants and DCs, simultaneously. It means that every customer may have to wait to get service in plants as may in DC. In fact, because establishing a new plant or even developing existing ones (to increase capacity) is very costly, possibility of congestion in higher level in supply chain increases.

In most studies, it is assumed that plants act as M/M/1 or M/M/s queuing systems. We can add plant capacity decision problem in our model in the cast of expected waiting time to control congestion in upstream level and then, similarly like DCs, utilize SOCP reformulation to solve new problems. But the key point is, now we have two sets of conic quadratic constraints, average waiting time function of M/M/1 and M/G/1 for plants and DCs respectively, and in this case Cplex or other softwares can solve at most moderate size problems in reasonable CPU time because as seen in computational results section the conic constraint have most effect in CPU time. The research is needed, first to model the problem as SOCP and second figure out whether we are able to solve large problems by adding some valid constraints or heuristics approach should also be applied.

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APPENDIX A

SGB128 DATA

Name	Population	Name	Population
Youngstown, OH	115436	Springfield, OH	72563
Yankton, SD	12011	Springfield, MO	133116
Yakima, WA	49826	Springfield, MA	152319
Worcester, MA	161799	Springfield, IL	100054
Wisconsin Dells, WI	2521	Spokane, WA	171300
Winston-Salem, NC	131885	South Bend, IN	109727
Winnipeg, MB	564473	Sioux Falls, SD	81343
Winchester, VA	20217	Sioux City, IA	82003
Wilmington, NC	139238	Shreveport, LA	205820
Wilmington, DE	70195	Sherman, TX	30413
Williston, ND	13336	Sheridan, WY	15146
Williamsport, PA	33401	Seminole, OK	8590
Williamson, WV	5219	Selma, AL	26684
Wichita Falls, TX	94201	Sedalia, MO	20927
Wichita, KS	279835	Seattle, WA	493846
Wheeling, WV	43070	Scranton, PA	88117
West Palm Beach, FL	63305	Scottsbluff, NB	14156
Wenatchee, WA	17257	Schenectady, NY	67972
Weed, CA	2879	Savannah, GA	141634
Waycross, GA	19371	Sault Sainte Marie, MI	14448
Wausau, WI	32426	Sarasota, FL	48868
Waukegan, IL	67653	Santa Rosa, CA	83320
Watertown, SD	15649	Santa Fe, NM	48953
Watertown, NY	27861	Santa Barbara, CA	74414
Waterloo, IA	75985	Santa Ana, CA	204023
Waterbury, CT	103266	San Jose, CA	629546
Washington, DC	638432	San Francisco, CA	678974
Warren, PA	12146	Sandusky, OH	31360
Walla Walla, WA	25618	San Diego, CA	875538
Waco, TX	101261	San Bernardino, CA	118794
Vincennes, IN	20857	San Antonio, TX	786023
Victoria, TX	50695	San Angelo, TX	73240
Vicksburg, MS	25434	Salt Lake City, UT	163697
Vancouver, BC	414281	Salisbury, MD	16429
		Continued	on next page

Table A.1: Name and population of cities

Name	Population	Name	Population
Valley City, ND	7774	Salinas, CA	80479
Valdosta, GA	37596	Salina, KS	41843
Utica, NY	75632	Salida, CO	44870
Uniontown, PA	14510	Salem, OR	89233
Tyler, TX	70508	Saint Paul, MN	270230
Twin Falls, ID	26209	Saint Louis, MO	453085
Tuscaloosa, AL	75211	Saint Joseph, MO	76691
Tupelo, MS	23905	Saint Joseph, MI	9622
Tulsa, OK	360919	Saint Johnsbury, VT	7150
Tucson, AZ	330537	Saint Cloud, MN	42566
Trinidad, CO	9663	Saint Augustine, FL	11985
Trenton, NJ	92124	Saginaw, MI	77508
Traverse City, MI	15516	Sacramento, CA	275741
Toronto, ON	599217	Rutland, VT	18436
Topeka, KS	115266	Roswell, NM	39676
Toledo, OH	354635	Rocky Mount, NC	41283
Texarkana, TX	31271	Rock Springs, WY	19458
Terre Haute, IN	61125	Rockford, IL	139712
Tampa, FL	271523	Rochester, NY	241741
Tallahassee, FL	81548	Rochester, MN	57890
Tacoma, WA	158501	Roanoke, VA	100220
Syracuse, NY	170105	Richmond, VA	219214
Swainsboro, GA	7602	Richmond, IN	41349
Sumter, SC	24890	Richfield, UT	5482
Stroudsburg, PA	5148	Rhinelander, WI	7873
Stockton, CA	149779	Reno, NV	100756
Stevens Point, WI	22970	Regina, SA	162613
Steubenville, OH	26400	Red Bluff, CA	9490
Sterling, CO	11385	Reading, PA	78686
Staunton, VA	21857	Ravenna, OH	11987

 Table A.1: Name and population of cities (continued)

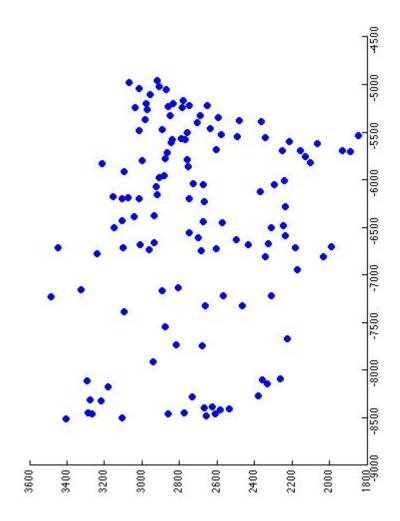


Figure A.1: Map of cities

APPENDIX B

COMPUTATIONAL RESULTS

The computational results of experiments are reported here. The meaning of abbreviations used in tables are as follows.

- **Rep** : Number of replications
- n : Number of candidate location for facilities
- $\bullet~m~$: Number of candidate location for CD centers
- k : Number of customers
- θ : Multiplier of waiting time cost
- δ : Multiplier of transportation cost
- β : Multiplier of fixed cost
- $\bullet~FC~:$ Fixed Cost
- CC : Congestion Cost
- \bullet **E**(**W**) : Total waiting time cost
- NOP : Number of opened Plants
- NODC : Number of opened CD centers

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B.1:
Table

-		10%	Gan(%)	R Gan(%		В	<i>δ B</i>	A S B
120,317	<u> </u>	<u> </u>	<u> </u>	0 1	1 0 1	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	
108,414	10		0 10				1 1 0	0.01 1 1 0
115,642	=		0 11				1 1 0	0.01 1 1 0
808,642	~		8 0		0	0	1 10 0	0.01 1 10 0
747,010			0		0	0	1 10 0	0.01 1 10 0
		0	0	10 0			1 10	0.01 1 10
496,763			0			1 0	10 1 0	0.01 10 1 0
426,618			0			1 0	1 10 1 0	0.01 10 1 0
469,785			0			1 0	10 1 0	0.01 10 1 0
1,196,805	,		0 1,		0	10 0	10 10 0	0.01 10 10 0
1,077,378	1,		0 1,		0	10 0	10 10 0	0.01 10 10 0
1,148,802	1,		0 1,		0	10 0	0.01 10 10 0	0.01 10 10 0
157,667	1:		0 1:				1 1 0	5 1 1 0
143,375	1		0 1				1 1 0	5 1 1 0
151,509	-		0 1				1 1 0	5 1 1 0 1
888,558	8		0		0	0	1 10 0	5 1 10 0
822,369	~		0		0	0	5 1 10 0	5 1 10 0
		0	0	10 0			1 10	5 1 10
		0	0	1 0	10 1 0	1	10 1	5 10 1
479,833	7		0			1 0	10 1 0	5 10 1 0
517,761	4,		0			1 0	10 1 0	5 10 1 0
1,279,432	1,	1	0 1,	1	0 1	10 0 1	10 10 0 1	5 10 10 0 1
1,158,091	1	[0 1	[0	10 0 1	10 10 0 1	5 10 10 0 1
		0	0	10 0		10	10 10	5 10 10
			_	_ ,	,		,	
· .						1 1 0	10 1 1 0	10 50 10 1 1 1 0 184,063

NODC	7	9	9	9	7	7	7	9	9	9	9	7	9	4	4	4	7	7	7	9	7	9	8	8	8	6
NOP	3	3	3	3	4	3	3	3	3	3	ю	3	3	3	3	ю	3	4	3	ю	3	3	4	4	4	4
E(w)	9.79	11.20	11.08	11.07	11.55	12.25	12.25	13.43	13.35	13.16	93.57	120.13	101.30	150.13	142.92	129.15	273.32	220.97	196.64	209.20	271.63	290.09	12.61	12.49	12.54	15.74
CC	52,785	64,138	59,544	61,679	66,656	67,357	69,927	80,023	74,784	76,216	069	768	750	1,126	929	696	2,038	1,430	1,464	1,552	1,754	2,158	35,772	33,251	34,706	46,937
TC	46,545	47,953	42,237	45,980	428,845	386,974	423,915	444,417	384,032	420,672	92,278	85,562	90,601	95,242	88,593	94,208	913,985	818,018	896,712	921,138	852,109	905,090	94,084	90,141	92,683	90,998
FC	78,299	796,530	739,791	772,317	105,100	75,002	78,299	796,530	739,791	772,317	108,061	105,186	107,312	1,066,868	1,033,864	1,059,428	108,879	138,229	108,125	1,077,246	1,051,849	1,069,731	114,556	111,016	113,764	1,117,236
CPU(s)	295.05	3,868.08	1,978.92	1,037.27	27.02	37.08	24.12	146.09	108.25	115.05	178.51	159.76	165.44	5,567.21	681.07	504.68	35.83	33.79	26.40	147.95	157.17	120.92	593.74	669.34	922.48	4944.73
obj	177,629	908,621	841,572	879,977	600,601	529,333	572,141	1,320,970	1,198,607	1,269,205	201,029	191,516	198,663	1,163,236	1,123,386	1,154,605	1,024,902	957,677	1,006,301	1,999,936	1,905,712	1,976,979	244,413	234,408	241,153	1,255,171
$\operatorname{Gap}(\%)$	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
β	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	1	-	-	10	10	10	-	-	1	10
δ	1	-	1	1	10	10	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	1	-	-	1
θ	10	10	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5	5	5
К	50	50	50	50	50	50	50	50	50	50	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100
Σ	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	ω	-	2	3	-	7	ω	-	7	ω	-	7	ω	-	7	с	-	7	ω	-	7	б	-	7	ю	1

Ŋ																										
NODC	9	9	7	7	7	9	9	9	٢	7	7	7	7	9	8	8	7	7	7	7	9	9	9	3	3	3
NOP	4	4	4	4	4	4	4	4	3	3	3	3	ю	3	5	5	4	3	3	3	3	3	3	3	3	ю
E(w)	15.64	15.86	19.74	19.51	19.61	21.38	21.41	21.22	12.85	12.86	12.94	12.93	14.25	15.10	14.11	14.15	15.44	18.01	17.04	17.38	109.35	117.87	108.72	168.50	154.68	160.87
CC	43,817	46,226	61,201	56,788	59,277	66,698	62,856	64,533	73,291	69,037	72,138	73,871	77,585	87,172	82,230	77,694	89,538	109,000	96,588	102,109	814	761	808	1,264	1,005	1,207
TC	87,142	91,201	838,533	801,532	825,678	843,733	805,193	832,386	102,188	95,870	99,625	104,712	103,077	102,801	887,752	828,306	900,414	950,526	895,529	935,650	116,563	109,195	113,173	121,865	115,573	120,727
FC	1,082,666	1,109,434	113,139	109,642	112,356	1,114,635	1,080,146	1,106,851	111,196	107,763	110,428	1,111,991	1,075,069	1,090,161	180,131	174,567	143,952	1,106,001	1,075,069	1,098,287	121,598	118,151	120,133	1,192,194	1,158,426	1,177,881
CPU(s)	6597.75	7203.03	120.12	81.88	68.45	351.14	230.24	830.3	1,625.61	1,355.54	1,533.29	7,200.85	7,202.46	7,203.69	105.29	100.31	100.07	862.95	424.25	683.46	308.94	219.15	298.82	231.16	487.19	372.20
ido	1,213,625	1,246,861	1,012,873	967,961	997,311	2,025,065	1,948,195	2,003,770	286,675	272,670	282,191	1,290,574	1,255,731	1,280,133	1,150,113	1,080,567	1,133,904	2,165,527	2,067,186	2,136,046	238,975	228,107	234,114	1,315,323	1,275,004	1,299,815
Gap(%)	0	2	0	0	0	0	0	0	0	0	0	4	S	б	0	0	0	0	0	0	0	0	0	0	0	0
β	10	10	-	-	-	10	10	10	-	1	-	10	10	10	1	1	1	10	10	10	1	1	-	10	10	10
δ	1	-	10	10	10	10	10	10	-	-	-	1	-		10	10	10	10	10	10	1	1	-	1	-	1
θ	5	5	5	5	5	5	5	5	10	10	10	10	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01
К	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	128	128	128	128	128	128
Σ	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	2	3	1	2	3	1	2	ю	1	2	3	1	5	ю	-	2	ю	-	2	3	1	2	3	1	2	ю

Table B.1: Experimental results (continued)

NODC	~	7	7	9	9	9	9	9	9	4	4	4	7	7	7	9	5	S	7	7	7	5	5	5	7	7
NOP	4	4	3	ω	3	3	4	4	4	4	4	4	4	4	4	4	4	4	3	з	3	3	3	ω	4	4
E(w)	165.09	217.54	285.61	297.12	291.49	434.58	7.72	7.71	7.71	10.17	12.33	10.2	8.07	8.1	8.1	10.48	10.37	10.07	10.25	10.20	10.15	13.40	13.39	13.38	12.53	12.74
СС	1,230	1,409	2,136	2,220	1,889	3,252	21,408	20,097	20,881	30,155	35,074	29,525	22,518	21,282	22,212	29,896	28,339	28,515	57,344	53,627	55,297	79,822	75,035	77,804	73,286	70,443
TC	1,113,882	1,039,647	1,118,951	1,161,108	1,089,788	1,125,786	107,368	102,090	104,001	112,663	105,349	114,941	1,060,597	1,009,126	1,024,484	1,065,737	1,022,171	1,038,119	123,321	115,047	120,125	117,290	108,570	113,474	1,093,359	1,009,234
FC	162,053	156,187	121,631	1,215,956	1,179,503	1,201,358	127,962	124,334	126,426	1,240,528	1,201,898	1,225,644	129,558	125,885	128,357	1,262,775	1,216,351	1,238,520	121,037	117,610	119,582	1,175,896	1,142,601	1,161,781	156,765	152,327
CPU(s)	73.21	68.33	70.90	326.14	193.61	304.01	655.47	793.47	1811.84	7204.14	7202.61	7202.77	124.61	105.63	144.63	546.71	752.83	480.14	913.87	703.13	764.11	2,186.99	2,858.31	1,653.05	60.92	59.89
óbj	1,277,165	1,197,243	1,242,718	2,379,284	2,271,180	2,330,396	256,738	246,520	251,308	1,383,345	1,342,321	1,370,110	1,212,673	1,156,293	1,175,053	2,358,408	2,266,861	2,305,154	301,702	286,284	295,004	1,373,008	1,326,206	1,353,059	1,323,410	1,232,004
$\operatorname{Gap}(\%)$	0	0	0	0	0	0	0	0	0	5	6	6	0	0	0	0	0	0	0	0	0	0	0	0	0	0
β	1	-	1	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	1	1	-	10	10	10	-	1
δ	10	10	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	1	1	-	-	-	-	10	10
θ	0.01	0.01	0.01	0.01	0.01	0.01	5	5	5	5	5	5	5	5	5	5	5	5	10	10	10	10	10	10	10	10
Х	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128
М	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	1	2	3	1	2	3	1	5	ε	1	5	3	1	7	ε	-	5	ю	1	5	3	-	5	ε	1	5

	E(1 2
	CC	75 175
(p	TC	15/ 881 1 050 070 75 75
lts (continued	FC	151 881
imental resul	CPU(s)	55 01
Table B.1: Experimental results (continued	obj	1 787 376
Tał	Jap(%)	0

NODC	7	S	S	S	4	4	4	ω	б	с	9	S	S	4	4	4	8	~	8	S	9	S	6	6	6	7
NOP	4	ω	ю	ю	ω	ю	ω	ю	ю	ю	ω	3	ю	б	3	ω	4	4	4	4	4	4	4	4	4	4
E(w)	13.08	15.13	15.05	15.42	31.32	29.10	59.86	124.74	121.72	121.99	38.19	33.65	39.50	98.16	122.47	167.35	9.13	9.1	9.12	13.37	11.59	13.47	10.6	11.42	10.85	11.83
CC	75,425	92,095	86,131	91,998	235	189	445	936	791	915	283	218	295	730	792	1,251	24,689	23,116	24,067	39,804	31,484	39,209	29,943	30,891	30,070	33,753
TC	1,052,070	1,143,434	1,058,454	1,101,022	38,752	33,798	37,966	41,542	34,994	39,813	381,206	334,403	375,002	389,959	338,941	378,543	42,930	37,063	40,926	49,179	36,716	43,193	403,764	343,928	384,755	408,470
FC	154,881	1,175,896	1,142,601	1,161,781	77,445	71,927	74,720	763,413	709,032	740,205	78,625	72,475	75,661	770,640	715,744	747,212	83,264	77,328	80,724	799,499	752,802	775,168	84,368	78,353	81,794	817,765
CPU(s)	55.91	601.93	661.54	310.01	169.73	130.60	199.71	80.50	42.70	193.60	26.13	25.12	29.55	141.93	150.82	127.42	459.98	290.22	210.91	7204.3	7206.17	7203.36	24.52	36.83	29.78	208.2
ido	1,282,376	2,411,425	2,287,186	2,354,800	116,432	105,914	113,131	805,891	744,817	780,933	460,114	407,095	450,959	1,161,329	1,055,477	1,127,006	150,884	137,507	145,717	888,482	821,002	857,570	518,075	453,172	496,619	1,259,988
Gap(%)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	2	2	-	0	0	0	0
β	1	10	10	10	-	-	-	10	10	10	-	1	-	10	10	10	1	-	-	10	10	10	-	-	-	10
δ	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	1	-	-	-	-	-	10	10	10	10
θ	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5	5	S	S	5	5	5	S	5
Х	128	128	128	128	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50
Я	10	10	10	10	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	ю	1	5	ю	1	2	б	1	2	б	1	2	ю	1	7	б	1	5	ю	1	7	б	1	7	б	-

NODC	9	9	6	6	6	7	5	9	7	7	7	5	S	9	5	5	5	4	4	4	7	7	7	5	5	5
NOP	4	4	3	3	3	3	3	3	4	3	4	3	3	3	3	3	3	3	3	ю	4	4	4	3	3	3
E(w)	12.37	14.53	8.62	8.52	8.51	9.80	13.39	11.21	11.62	10.72	11.72	13.40	13.40	11.63	97.16	71.58	80.82	140.88	110.81	135.71	62.74	100.60	78.51	218.03	205.89	138.17
CC	34,096	42,370	45,736	42,349	43,916	54,159	75,005	62,670	67,154	57,140	66,206	79,805	75,114	65,620	726	465	606	1,057	720	1,018	467	652	585	1,632	1,333	1,031
TC	353,083	390,801	45,674	39,517	44,476	49,407	36,514	41,753	369,092	344,445	360,356	406,896	341,334	395,485	80,161	73,882	79,920	80,833	74,877	82,023	741,901	680,463	738,986	799,687	738,786	800,244
FC	752,802	782,175	82,965	77,052	80,439	807,569	729,538	772,317	105,100	75,002	101,902	785,491	729,538	772,317	107,503	104,256	106,835	1,066,868	1,033,864	1,059,428	142,977	138,557	141,985	1,072,451	1,039,274	1,064,971
CPU(s)	320.58	658.39	358.54	581.43	405.57	7,204.97	7,203.25	7,209.09	47.97	54.63	49.70	219.81	221.91	327.51	398.47	311.61	322.55	940.11	2,463.40	2,398.56	58.08	69.31	61.84	645.74	434.40	490.87
obj	1,139,981	1,215,347	174,376	158,918	168,831	911,135	841,057	876,740	541,346	476,587	528,465	1,272,193	1,145,986	1,233,422	188,390	178,603	187,361	1,148,758	1,109,461	1,142,469	885,344	819,672	881,556	1,873,770	1,779,393	1,866,246
Gap(%)	0	0	0	0	0	S	4	S	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
β	10	10	1	-	-	10	10	10	-	-	1	10	10	10	-	-	-	10	10	10	1	-	-	10	10	10
δ	10	10	1	-	-	1	1	-	10	10	10	10	10	10	-	-	-	-	1	-	10	10	10	10	10	10
θ	5	5	10	10	10	10	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01
К	50	50	50	50	50	50	50	50	50	50	50	50	50	50	100	100	100	100	100	100	100	100	100	100	100	100
Я	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	7	æ	1	7	ε	1	7	ε	1	7	ε	-	7	ε	-	7	ε	1	2	ε	-	7	ω	-	7	3

NODC	8	8	8	7	7	7	10	10	11	7	7	7	10	10	10	7	7	9	6	8	6	9	9	9	S	4
NOP	4	4	4	4	4	4	4	4	4	4	4	4	ю	ю	e	3	3	ю	5	5	5	3	ю	ю	e	3
E(w)	11.8	11.93	11.79	13.71	14.34	17.38	14.25	14.69	13.37	15.42	15.31	16.6	10.64	10.64	10.75	15.64	14.97	17.80	13.65	13.98	13.68	16.65	16.85	16.68	30.80	37.60
CC	32,890	31,376	32,098	39,715	38,987	50,124	41,334	40,316	37,345	45,517	42,473	48,181	57,552	54,159	56,915	92,271	82,436	105,493	78,862	76,503	77,242	100,317	95,748	98,192	228	244
TC	80,835	75,245	80,487	80,922	78,327	86,277	779,865	722,834	774,375	789,964	735,297	781,454	92,226	86,137	91,202	114,969	103,644	117,310	732,794	673,875	731,472	816,490	752,417	813,909	97,344	91,060
FC	114,556	111,016	113,764	1,131,410	1,093,881	1,115,760	116,192	112,602	117,211	1,128,809	1,093,881	1,118,343	115,447	111,885	114,652	1,106,001	1,071,784	1,087,578	181,288	174,567	180,033	1,097,817	1,063,854	1,090,161	121,312	116,978
CPU(s)	1562.6	1429.2	1775.9	7200.62	7203.39	7203.89	105.89	132.3	100.29	360.96	501.78	992.04	3,834.46	4,762.76	4,413.18	7,203.83	7,203.72	7,203.14	193.38	239.34	298.59	1,295.70	1,637.57	1,035.60	1,014.60	1,559.84
ido	228,281	217,637	226,349	1,252,047	1,211,196	1,252,161	937,391	875,752	928,932	1,964,290	1,871,651	1,947,978	265,225	252,181	262,769	1,313,241	1,257,864	1,310,380	992,944	924,945	988,747	2,014,624	1,912,019	2,002,262	218,884	208,282
Gap(%)	0	0	0	10	10	10	0	0	0	0	0	0	0	0	0	25	24	~	0	0	0	0	0	0	0	0
β	1	1	1	10	10	10	-	-	1	10	10	10	-	1	1	10	10	10	1	-	1	10	10	10	-	1
δ	1	1	1	1	-		10	10	10	10	10	10		1		-	-		10	10	10	10	10	10		1
θ	5	5	5	5	5	5	S	S	5	5	S	S	10	10	10	10	10	10	10	10	10	10	10	10	0.01	0.01
К	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	128	128
Σ	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	1	2	ю	1	2	3	1	5	3	1	7	3	1	2	ε	-	2	ю	1	2	ю		2	3	1	2

NODC	4	ю	3	ю	7	9	9	5	4	4	7	7	7	5	4	4	11	11	11	7	7	7	9	9	8	5
NOP	3	3	3	3	4	4	4	3	3	3	4	4	4	4	4	4	4	4	4	4	4	4	3	3	3	3
E(w)	43.72	114.51	123.31	118.77	39.27	37.58	50.83	71.40	44.74	61.97	6.83	6.78	6.76	12.68	12.59	14.64	7.4	7.74	7.59	11.93	10.51	10.85	7.30	7.38	7.49	14.67
CC	324	859	802	891	291	244	379	533	291	461	18,226	16,978	17,552	37,716	36,128	44,131	19,847	19,629	19,781	33,763	27,607	29,737	39,840	38,031	39,562	88,361
TC	97,837	105,418	96,700	101,197	924,887	854,132	904,761	974,410	910,585	977,865	100,765	93,988	98,346	112,950	111,178	119,461	987,966	914,657	963,896	1,001,801	929,430	977,539	102,426	96,081	102,512	132,308
FC	118,737	1,190,856	1,157,127	1,176,559	161,020	155,561	157,999	1,211,357	1,169,790	1,187,388	129,917	126,233	128,357	1,247,044	1,201,898	1,219,207	134,638	130,469	132,202	1,270,474	1,237,968	1,260,173	125,778	122,211	127,417	1,224,524
CPU(s)	1,283.14	7,127.78	3,243.37	2,800.01	159.03	221.40	167.70	1,078.65	2,658.43	763.87	7100.18	2167.68	2838.34	7202.46	7204.03	7202.88	247.35	241.82	299.4	2273.22	3561.91	4500.58	4,423.13	7,202.57	7,202.38	7,202.69
óbj	216,898	1,297,133	1,254,629	1,278,647	1,086,198	1,009,937	1,063,139	2,186,300	2,080,666	2,165,714	248,909	237,199	244,256	1,397,710	1,349,203	1,382,799	1,142,451	1,064,755	1,115,879	2,306,037	2,195,005	2,267,448	268,043	256,323	269,491	1,445,193
Gap(%)	0	0	0	0	0	0	0	0	0	0	0	0	0	б	10	4	0	0	0	0	0	0	0	7	9	25
β	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	-	-	1	10	10	10	-	-	-	10
δ	-	-	1	1	10	10	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	-	-	-	-
θ	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5	5	5	5	5	5	5	5	5	5	5	10	10	10	10
Х	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128
М	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
Rep	ω	1	2	ю	-	7	ω	1	6	ω	1	6	ω	-	6	ω	-	6	ю	-	7	ω	1	6	ω	

NODC	5	5	7	7	7	5	4	5	4	4	4	ε	ю	ю	4	5	5	4	4	4	8	8	8	5	5	5
NOP	3	3	4	4	4	ю	ю	ю	3	ю	3	3	3	ю	ю	3	3	ю	ю	3	4	4	4	4	4	4
E(w)	18.99	15.23	8.58	8.89	8.93	10.76	11.30	10.01	63.00	66.04	57.48	123.55	130.99	121.80	64.05	82.56	92.21	146.35	111.76	202.27	9.53	9.47	9.52	13.44	13.36	13.48
CC	114,402	112,221	48,699	47,915	49,995	63,739	64,315	57,156	469	424	427	927	851	914	480	536	688	1,094	721	1,512	26,108	24,375	25,474	40,037	37,420	39,231
TC	97,519	120,100	929,583	857,883	910,258	982,595	920,810	978,836	41,496	35,082	38,907	44,371	37,391	41,632	412,697	347,263	384,886	412,675	349,206	385,306	45,926	39,673	43,679	44,197	38,840	41,463
FC	1,188,785	1,221,524	165,097	160,416	163,110	1,231,698	1,184,150	1,219,804	73,815	68,552	71,567	730,908	678,831	708,666	74,196	69,454	72,508	738,135	685,543	715,673	83,264	77,328	80,724	799,499	742,549	775,168
CPU(s)	7,204.77	7,202.03	412.93	603.74	340.85	2,967.23	2,949.04	3,305.04	180.87	324.42	264.33	153.43	287.51	104.91	54.65	53.90	56.16	236.22	359.77	252.55	412.81	1245.47	1101.43	5337.46	4714.72	4320.67
ido	1,400,705	1,453,845	1,143,379	1,066,214	1,123,363	2,278,032	2,169,275	2,255,796	115,780	104,058	110,901	776,206	717,073	751,212	487,373	417,253	458,082	1,151,904	1,035,470	1,102,491	155,299	141,376	149,877	883,732	818,809	855,863
Gap(%)	8	11	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
β	10	10			-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	-	-		10	10	10
δ	1		10	10	10	10	10	10	1	-	-		-	-	10	10	10	10	10	10	1	1	-	-		1
θ	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5	S	5	S	5
Х	128	128	128	128	128	128	128	128	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50
Σ	20	20	20	20	20	20	20	20	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	10	10	10	10	10	10	10	10	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	2	3	1	2	Э	-	2	3	-	2	3	-	2	ю	1	2	3	1	2	Э	1	5	3	1	2	ю

Table B.1: Experimental results (continued)

NODC	8	8	7	9	S	5	7	7	7	9	9	9	9	9	9	5	5	S	5	5	5	4	4	4	5	5
NOP	5	5	4	4	4	4	3	3	3	3	ю	3	4	4	4	3	3	3	3	3	3	3	3	3	3	3
E(w)	12.33	12.14	14.56	17.34	18.55	18.87	9.74	9.71	9.71	11.04	10.95	10.95	11.63	11.63	11.31	14.22	13.88	14.00	76.55	69.92	92.15	117.24	115.26	138.94	167.77	143.86
CC	36,089	33,200	42,974	52,844	54,209	57,494	53,725	50,367	52,252	62,988	58,680	60,905	67,162	63,194	63,407	85,630	78,330	82,054	572	452	687	879	749	1,042	1,256	933
TC	399,387	342,546	384,230	411,323	351,424	379,212	51,680	44,231	48,470	48,930	42,011	45,412	410,469	349,493	388,404	428,575	364,394	401,618	79,645	74,218	78,444	80,241	75,040	79,477	794,068	739,988
FC	101,872	94,256	79,654	801,588	739,008	771,471	77,508	71,981	75,146	764,025	709,590	740,778	99,664	92,558	96,628	752,986	699,337	730,074	102,994	99,810	102,279	1,021,787	990,163	1,014,650	102,994	99,810
CPU(s)	74.74	107.49	95.79	1455.43	2809.27	1084.43	2,486.78	2,660.11	1,393.23	7,203.33	7,200.88	7,200.13	50.17	72.81	59.73	597.52	758.62	481.03	709.51	912.42	2,305.21	1,692.83	7,210.32	7,201.27	106.00	137.75
óbj	537,348	470,002	506,858	1,265,755	1,144,641	1,208,177	182,913	166,579	175,868	875,943	810,281	847,095	577,295	505,245	548,440	1,267,191	1,142,061	1,213,746	183,211	174,480	181,410	1,102,907	1,065,953	1,095,169	898,318	840,731
Gap(%)	0	0	0	0	0	0	0	0	0	ю	ю	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0
β	1	1	-	10	10	10	-	-	1	10	10	10	-	-	1	10	10	10	1	1	-	10	10	10	1	1
δ	10	10	10	10	10	10	1	-	1	-	-	-	10	10	10	10	10	10	1	1	-	-	1	-	10	10
θ	5	5	5	5	5	5	10	10	10	10	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01
м	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	100	100	100	100	100	100	100	100
Σ	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	1	2	б	1	7	ю	1	7	Э	1	7	б	1	7	ю	1	7	ω	1	7	б	1	2	ю	1	2

NODC	S	4	5	S	7	7	7	S	9	9	8	9	9	9	9	5	7	7	~	7	9	9	9	9	7	5
NOP	3	ω	ω	ω	4	4	4	4	4	4	5	5	S	4	4	4	б	ω	ω	ω	б	б	4	4	S	ю
E(w)	144.31	284.34	165.48	248.72	13.53	13.35	13.7	20.3	16.54	19.22	16.44	18.95	18.88	23.62	23.54	24.18	12.85	13.60	13.16	13.96	15.26	19.25	16.16	16.01	14.04	21.04
CC	1,078	2,133	1,074	1,854	39,066	36,147	38,737	63,177	46,827	57,294	49,436	54,903	56,729	74,415	69,719	75,160	73,288	73,982	73,083	80,484	85,118	114,530	96,885	90,110	79,832	131,644
TC	782,832	799,866	739,773	784,166	87,124	86,357	86,578	90,148	90,042	88,891	778,132	732,053	757,373	790,291	751,112	780,209	99,651	93,758	100,406	102,856	112,022	102,795	788,699	736,048	759,464	836,663
FC	102,279	1,021,787	998,093	1,019,410	113,139	109,642	112,356	1,103,062	1,082,666	1,103,485	140,361	133,274	136,573	1,112,034	1,077,626	1,095,359	106,687	103,392	107,098	1,064,309	1,020,153	1,039,434	137,526	133,276	170,011	1,038,562
CPU(s)	121.34	554.27	797.46	686.92	7202.35	7202.6	7203.08	7200.2	7203.02	7200.2	540.86	354.95	417.07	3867.23	3108.6	3523	7,203.41	7,201.83	7,200.93	7,202.41	7,201.26	7,204.49	465.73	671.68	512.10	2,126.68
óbj	886,189	1,823,786	1,738,940	1,805,430	239,328	232,147	237,671	1,256,388	1,219,535	1,249,670	967,929	920,230	950,676	1,976,739	1,898,456	1,950,728	279,626	271,132	280,587	1,247,649	1,217,292	1,256,759	1,023,110	959,434	1,009,307	2,006,869
Gap(%)	0	0	0	0	0	3	4	4	4	4	0	0	0	0	0	0	13	14	13	9	8	8	0	0	0	0
β	1	10	10	10	-	1	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	-	-	-	10
δ	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	-	-	-	-	-	-	10	10	10	10
θ	0.01	0.01	0.01	0.01	5	5	5	S	5	S	5	5	S	S	S	5	10	10	10	10	10	10	10	10	10	10
К	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100
М	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	ε	1	5	ε	-	5	e	1	5	б	1	2	e	1	2	ю	-	5	б	-	2	б	1	7	ε	-

NODC	5	5	5	5	5	4	4	ю	5	5	5	5	S	S	9	9	9	ω	Э	4	~	9	9	5	9	5
NOP	ю	ю	ω	ω	ω	ю	3	ю	ω	ю	ю	ю	ω	ω	4	4	4	4	4	4	5	5	4	4	4	4
E(w)	20.93	21.27	68.97	53.65	66.67	211.98	181.75	123.53	101.70	41.57	69.97	89.00	172.18	69.97	8.8	8.32	7.99	18.46	20.37	12.59	7.99	9.6	11.79	14.57	12.7	14.37
CC	123,143	130,077	507	346	495	1,590	1,181	926	753	270	519	657	1,116	519	24,898	22,012	21,828	59,653	62,541	37,240	22,133	26,331	34,769	44,375	35,532	42,561
TC	775,937	827,946	97,742	90,208	94,349	100,972	93,064	103,489	976,882	901,763	943,280	976,986	902,693	943,280	103,093	95,666	98,473	122,996	98,839	105,361	970,575	911,099	949,468	1,009,154	939,139	956,902
FC	1,006,418	1,031,308	115,667	112,676	114,277	1,147,777	1,115,825	1,127,469	115,667	112,878	114,277	1,156,667	1,123,916	1,142,783	127,244	123,986	126,426	1,220,980	1,186,389	1,222,101	159,745	152,221	125,429	1,243,227	1,221,819	1,228,311
CPU(s)	1,839.05	2,264.87	4,065.93	2,789.22	6,317.96	7,201.97	7,202.83	7,203.33	427.99	271.47	268.46	1,893.62	2,379.95	3,278.36	7203.44	7201.89	7202.86	7202.99	7200.41	7203.38	1452.1	738.6	899.52	7203.02	7202.96	7201.8
obj	1,905,498	1,989,332	213,916	203,230	209,121	1,250,339	1,210,070	1,231,885	1,093,302	1,014,911	1,058,076	2,134,311	2,027,725	2,086,582	255,235	241,664	246,727	1,403,628	1,347,769	1,364,702	1,152,453	1,089,651	1,109,666	2,296,756	2,196,491	2,227,774
Gap(%)	0	0	0	0	0	19	-	0	0	0	0	0	0	0	7	e	9	11	11	10	0	0	0	4	4	4
β	10	10	-	-	-	10	10	10	-	-	1	10	10	10	-	-	-	10	10	10	1	-	1	10	10	10
δ	10	10	-	-	-	1	1	-	10	10	10	10	10	10	-	-	-	-	1	-	10	10	10	10	10	10
θ	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	S	S	5	5	5	5	5	5	5	S	S	5
К	100	100	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128
Σ	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10	10
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	2	ω	-	7	ω	-	2	ω	-	5	б	-	7	ω	-	7	ω	-	2	ю	-	7	ы	-	2	3

(۲)									1																	
NODC	9	5	9	4	9	4	9	9	9	4	5	4	5	4	5	æ	æ	3	9	9	٢	4	4	5	~	∞
NOP	3	e	ю	3	Э	e	4	4	4	ю	e	e	ю	e	e	e	e,	e.	3	ю	4	e	Э	3	4	4
E(w)	8.90	8.15	12.74	98.22	11.47	14.22	8.40	8.44	8.67	11.05	10.69	13.76	44.80	60.54	43.56	120.40	132.54	126.31	46.76	79.82	20.14	106.21	60.54	60.24	9.42	9.23
CC	50,114	43,164	72,490	681,815	63,491	87,001	47,639	45,092	48,435	66,516	59,288	82,710	330	389	321	903	862	947	342	510	145	796	389	446	25,733	23,570
TC	117,344	94,668	116,950	153,757	104,428	103,520	951,159	881,031	917,287	997,600	914,653	1,008,888	31,572	28,578	31,084	33,317	29,395	32,696	312,637	278,869	283,355	320,941	285,794	310,493	39,739	34,700
FC	119,981	115,380	117,369	1,154,198	1,159,284	1,153,868	157,022	152,574	155,134	1,167,888	1,145,788	1,150,325	74,272	68,429	72,010	730,908	678,831	708,666	74,862	69,525	96,204	736,809	684,312	720,110	83,264	77,328
CPU(s)	7,201.85	7,200.87	7,201.05	7,201.86	7,202.30	7,202.29	687.04	644.99	565.97	7,202.28	5,876.49	7,201.41	676.78	637.97	743.17	1,927.38	1,779.92	452.11	44.66	62.49	118.28	475.51	1,041.68	753.80	2820.14	3059.62
ĺdo	287,439	253,212	306,809	1,989,770	1,327,203	1,344,390	1,155,820	1,078,697	1,120,856	2,232,004	2,119,729	2,241,923	106,174	97,396	103,415	765,128	709,088	742,309	387,840	348,904	379,704	1,058,545	970,495	1,031,049	148,736	135,599
Gap(%)	17	6	22	47	9	9	0	0	0	9	0	12	0	0	0	0	0	0	0	0	0	0	0	0	0	0
β	-	-	-	10	10	10	-	-	-	10	10	10		-	-	10	10	10	-	-	-	10	10	10		-
δ	1	-		-	-	-	10	10	10	10	10	10	-	-	-	-	-	1	10	10	10	10	10	10	-	-
θ	10	10	10	10	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5
К	128	128	128	128	128	128	128	128	128	128	128	128	50	50	50	50	50	50	50	50	50	50	50	50	50	50
Σ	10	10	10	10	10	10	10	10	10	10	10	10	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	1	7	ω	-	7	ю	-	7	ω	-	5	ω	-	7	ю	-	2	ю	-	7	ю	-	7	ε	-	5

NODC	~	5	5	9	6	8	6	9	5	9	6	6	10	7	9	7	8	8	8	9	9	9	4	4	4	3
NOP	4	4	4	4	4	4	4	4	4	4	ю	3	3	ю	ю	ю	5	5	5	3	ю	ю	3	ω	e	3
E(w)	9.31	13.5	18.01	11.15	13.07	14.26	12.24	16.28	16.17	15.27	8.51	8.52	8.23	9.86	13.79	11.13	10.75	10.46	10.56	14.51	13.32	14.98	43.93	65.90	40.32	126.56
CC	24,716	40,250	52,999	31,116	38,547	40,325	34,908	49,625	46,832	44,952	44,953	42,346	41,947	54,597	77,705	61,409	60,903	55,385	58,161	87,680	74,551	87,834	327	424	299	949
TC	38,516	42,050	39,776	41,472	356,613	301,283	347,506	359,921	309,143	348,454	43,582	37,589	42,693	57,580	40,477	50,372	287,413	253,623	274,216	336,780	302,617	335,598	64,924	59,812	63,766	66,633
FC	80,724	799,499	742,549	785,872	83,987	76,974	81,794	805,400	742,549	780,890	79,716	74,031	78,356	775,064	709,590	747,785	125,132	116,210	121,320	764,025	709,590	737,081	102,363	99,019	101,648	1,013,619
CPU(s)	1262.47	7206.36	7203.83	7201.85	126.14	195.39	133.72	2525.98	3836.3	2031.59	7,205.87	7,205.30	7,202.55	7,204.30	7,203.78	7,204.94	121.12	157.53	171.84	2,885.66	3,584.33	7,203.86	2,733.62	3,246.74	3,088.90	2,959.67
obj	143,956	881,798	835,325	858,460	479,148	418,582	464,208	1,214,946	1,098,523	1,174,296	168,251	153,966	162,996	887,241	827,772	859,565	473,448	425,217	453,697	1,188,485	1,086,758	1,160,514	167,613	159,255	165,713	1,081,201
$\operatorname{Gap}(\%)$	0	ю	11	~	0	0	0	0	0	0	9	8	10	~	~	7	0	0	0	0	0	4	0	0	0	0
β	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	1		-	10	10	10	-	-	1	10
δ	1	1	-	-	10	10	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	1	-	1	1
θ	5	5	5	5	5	5	5	5	5	5	10	10	10	10	10	10	10	10	10	10	10	10	0.01	0.01	0.01	0.01
К	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	50	100	100	100	100
М	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	ω	-	7	ω	-	2	с	-	7	ω	-	7	ω	-	7	ω	-	7	ω	1	7	б	-	7	ю	1

NODC	б	Э	9	7	7	4	4	4	8	8	10	9	6	7	6	6	10	9	7	9	10	6	10	7	9	9
NOP	e	3	4	4	4	ю	ю	e	4	4	4	4	4	4	S	5	5	4	4	4	ю	e	ω	e	e	ю
E(w)	120.12	131.03	24.32	54.62	22.94	70.11	100.57	107.88	12.55	12.1	11.7	18.31	16.22	13.28	16.43	15.76	15.24	21.39	20.4	20.84	8.72	8.93	8.53	12.64	11.98	11.45
CC	781	983	177	350	165	521	650	803	35,279	31,954	31,588	56,092	44,105	37,277	49,183	43,986	43,856	66,562	58,928	63,283	46,460	45,133	44,010	69,990	68,068	70,342
TC	62,155	65,518	593,486	553,562	584,164	649,736	597,148	638,842	79,187	72,988	80,132	83,724	87,858	85,954	697,655	651,592	699,592	754,181	690,168	743,709	84,045	78,589	95,452	81,155	113,899	110,210
FC	982,251	1,006,542	136,784	132,834	136,643	1,021,803	990,181	1,014,668	114,296	111,016	116,062	1,117,236	1,111,271	1,123,509	140,919	136,562	141,351	1,114,635	1,091,361	1,106,851	112,078	107,130	111,293	1,041,595	1,052,151	1,050,111
CPU(s)	7,201.69	2,429.26	145.02	223.91	256.00	4,072.03	3,333.66	3,212.93	7202.97	7202.97	7203.53	7202.86	7204.72	7203.99	340.55	443.93	416.27	7203.69	7200.46	7201.83	7,203.42	7,203.47	7,204.64	7,204.44	7,202.19	7,202.40
óbj	1,045,187	1,073,043	730,447	686,746	720,972	1,672,060	1,587,979	1,654,313	228,763	215,958	227,782	1,257,052	1,243,234	1,246,740	887,757	832,140	884,799	1,935,378	1,840,457	1,913,842	242,583	230,852	250,755	1,192,740	1,234,117	1,230,663
Gap(%)	19	0	0	0	0	0	0	0	7	7	4	S	7	5	0	0	0	S	-	4	16	17	23	~	26	11
β	10	10	-	-	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10
δ	-	-	10	10	10	10	10	10							10	10	10	10	10	10						-
θ	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5	5	S	S	S	5	5	5	S	5	5	10	10	10	10	10	10
К	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100	100
Σ	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	2	3	1	2	3	-	2	3	1	2	3	-	2	ю	-	2	3	-	2	3	1	2	3	1	2	ю

(continued)
results
Experimental
B.1:
Table

NODC	8	8	8	4	S	9	S	5	9	4	5	4	6	6	6	5	9	4	9	9	7	5	4	4	10	10
NOP	5	5	5	ω	ю	ю	б	ю	ω	ω	ю	ю	S	5	S	ω	б	ω	4	4	4	4	4	4	s	5
E(w)	11.86	11.33	11.33	20.41	18.29	19.48	74.33	132.29	153.08	88.83	75.49	91.66	55.21	64.28	41.92	137.19	394.13	160.86	8.05	7.44	7.85	15.55	10.38	48.59	8.51	8.42
СС	68,491	61,056	63,307	129,717	106,829	118,268	555	860	1,147	999	489	687	407	409	308	1,027	2,557	1,206	22,585	19,202	21,231	49,005	29,049	160,840	23,441	21,720
TC	580,819	539,063	568,188	670,081	623,285	666,267	84,875	80,083	81,290	87,192	81,048	83,827	727,548	680,206	705,131	847,400	833,385	826,203	98,101	91,199	94,504	143,782	95,050	99,382	890,265	837,812
FC	172,470	167,649	171,790	1,028,928	1,006,799	1,039,824	116,275	113,355	115,589	1,150,611	1,126,985	1,136,802	192,492	186,584	190,471	1,159,826	1,132,230	1,136,802	127,962	124,334	128,003	1,253,560	1,205,384	1,215,664	161,023	156,172
CPU(s)	347.09	371.17	445.18	7,201.33	6,622.79	7,202.13	7,202.43	7,203.33	7,201.58	7,203.41	7,204.28	7,203.16	319.71	483.32	283.83	5,219.28	7,203.50	3,031.49	7213.88	7204.13	7209.84	7201.29	7201.27	7200.76	885.12	806.21
įdo	821,780	767,768	803,285	1,828,726	1,736,913	1,824,359	201,706	194,298	198,026	1,238,470	1,208,522	1,221,317	920,448	867,199	895,910	2,008,253	1,968,172	1,964,211	248,647	234,735	243,738	1,446,347	1,329,483	1,475,885	1,074,729	1,015,704
Gap(%)	0	0	0	2	0	10	4	6	~	1	-	-	0	0	0	0	6	0	6	9	7	7	5	10	0	0
β	1	1	1	10	10	10	1	1	1	10	10	10	1	1	1	10	10	10	1	1	1	10	10	10	1	1
δ	10	10	10	10	10	10	-	-	-	-	-	-	10	10	10	10	10	10	-	-	-	-	-	-	10	10
θ	10	10	10	10	10	10	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	0.01	5	5	5	5	5	5	S	5
х	100	100	100	100	100	100	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128
Я	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	1	2	3	-	5	ю	-	5	ю	-	2	3		2	ю	-	2	3	-	5	3	-	5	ю	-	2

NODC	10	S	9	9	11	11	6	7			6	6	6	7	9	7
NOP	S	4	4	4	ω	ω	ω	б			5	S	5	ω	б	3
E(w)	8.23	10.69	11.61	13.88	11.76	11.56	13.28	11.99			10.51	10.65	11.16	18.46	15.84	17.69
CC	22,168	31,301	31,545	40,494	64,920	59,852	73,966	64965.96			59,004	56,524	61,802	113,578	90,726	105,601
TC	877,118	968,646	889,440	931,272	124,072	129,699	134,150	125116			733,046	686,468	708,658	853,181	791,368	823,046
FC	159,853	1,251,814	1,223,515	1,247,204	125,265	122,005	120,898	1170320			196,536	190,972	193,586	1,186,301	1,153,568	1,176,154
CPU(s)	670.74	7207.81	7204.58	7201.57	7,202.41	7,203.36	7,204.95	7200.68			702.43	652.80	737.78	7,202.08	7,200.55	7,203.39
obj	1,059,139	2,251,762	2,144,500	2,218,970	314,257	311,556	329,014	1360401.96			988,586	933,964	964,046	2,153,060	2,035,663	2,104,800
Gap(%)	0	5	5	9	26	31	31	26		1	0	0	0	10	e	12
β	-	10	10	10	-	-	-	10	10	10	-	-	-	10	10	10
δ	10	10	10	10		-	-	-	-		10	10	10	10	10	10
θ	5	5	5	5	10	10	10	10	10	10	10	10	10	10	10	10
м	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128	128
Σ	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
z	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20	20
Rep	ε	1	5	ю	1	7	ю	-	7	ю	-	7	3	1	7	3

Table B.1: Experimental results (continued)