HYBRID METAHEURISTIC ALGORITHMS FOR SINGLE AND MULTI-OBJECTIVE 2D BIN PACKING PROBLEM

A THESIS SUBMITTED TO<br>THE GRADUATE SCHOOL OF NATURAL AND APPLIED SCIENCES OF<br>MIDDLE EAST TECHNICAL UNIVERSITY

BY

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## IN PARTIAL FULFILLMENT OF THE REQUIREMENTS FOR <br> THE DEGREE OF MASTER OF SCIENCE <br> IN <br> COMPUTER ENGINEERING

Approval of the thesis:

## HYBRID METAHEURISTIC ALGORITHMS FOR SINGLE AND MULTI-OBJECTIVE 2D BIN PACKING PROBLEM


#### Abstract

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#### Abstract

HYBRID METAHEURISTIC ALGORITHMS FOR SINGLE AND MULTI-OBJECTIVE 2D BIN PACKING PROBLEM


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January 2015, 86pages

2D Bin packing problem (2DBPP) is an NP-hard combinatorial optimization problem. Objects with different width and length sizes are packed in order to minimize the number of unit-capacity bins according to an objective function. Single or multiobjective versions of this well-known industrial engineering problem can be faced frequently in real life situations. There have been several heuristics proposed for the solution of 2DBPP until now where it is not possible to find the exact solutions for large problem instances. Next fit, First Fit, Best Fit, Unified Tabu Search and Genetic Algorithms are some of these algorithms. Recently, Memetic Algorithms have put themselves forward as a new area of research in evolutionary computation with their ability to combine different heuristics and local search mechanisms together for higher quality solutions. In this thesis, we propose a set of single and multiobjective memetic and genetic algorithms that make use of the state-of-the-art metaheuristics and local search techniques for the solution of 2DBPP. We analyze the optimization time and the resulting solution quality of the algorithms on a 2DBPP benchmark offline problem set with 500 instances. Through results of exhaustive experiments and with the aid of a novel visual analyzer developed in this study, we conclude that the proposed memetic and hybrid genetic algorithms are robust with their ability to obtain very high percentage of the optimal solutions for the given benchmark problem instances.

Keywords: 2D Bin packing, memetic, genetic, load balancing, multiobjective, heuristic

## ÖZ

# TEK VE ÇOK AMAÇLI İKİ BOYUTLU KUTU PAKETLEME PROBLEMİ İÇİN MELEZ METASEZGİSEL ALGORİTMALAR 

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Ocak 2015, 86sayfa

İki boyutlu kutu paketleme problemi çözümü polinom zamanlı olmayan (NP) kombinatoriyal bir problemdir. Farklı genişlik ve uzunluktaki nesnelerin en az kutu kaplayacak şekilde birim kapasiteli kutulara bir nesnel fonksiyonun en az sonuç üretecek şekilde yerleştirilmesidir. Çok bilinen bu endüstri probleminin çok amaçlı versiyonlarıyla gündelik hayatta sıkça karşılaşılmaktadır. Kesin çözüm bulmanın mümkün olmadığı büyük ölçekli 2 boyutlu kutu paketleme problemleri için şimdiye kadar birkaç sezgisel yöntem önerilmiştir. Sonraki Ekleme, İlk Ekleme, En İyi Ekleme, Birleşik Tabu Araması ve Genetik Algoritmalar bu algoritmaların en çok kullanılan birkaçıdır. Son zamanlarda, farklı sezgisel ve lokal aramaları birleştirerek daha yüksek kalitede sonuç üreten Memetik Algoritmalar evrimsel hesaplamalarda yeni bir araştırma alanı olmuştur. Bu tezde, 2 boyutlu kutu paketleme problemi için metasezgisel ve lokal aramayı kullanan tek ve çok amaçlı memetik ve genetik algoritmalardan oluşan bir çözüm kümesi sunduk. 500 problem içeren çevrimdışı bir problem kümesi ile algoritmalarımızın optimizasyon zamanlarını ve sonuç kalitelerini analiz ettik. Bu projede geliştirdiğimiz görsel analiz aracı ile de yapılan yoğun testler sonucunda, geliştirdiğimiz memetik ve genetik algoritmaların kullanılan problem kümeleri üzerinde yüksek oranda optimum sonuç ürettiği neticesine vardık.

Anahtar Kelimeler: 2B Kutu paketleme, memetik, genetik, yük dengeleme, çok amaç, sezgisel

To my family and people who are reading this page

## ACKNOWLEDGMENTS

Many thanks to my supervisor, Prof. Dr. Ahmet Coşar, for spending his months and effort to propose this study. It has been a great privilege to have been mentored by such an esteemed researcher.

Special thanks to Dr. Tansel Dökeroğlu who provided encouragement for this study and also for his valuable insight and his guidance about this study.

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## LIST OF ABBREVIATIONS

| 1DBPP | One-Dimensional Bin Packing Problem |
| :--- | :--- |
| 2DBPP | Two-Dimensional Bin Packing Problem |
| 3DBPP | Three-Dimensional Bin Packing Problem |
| AIS | Artificial Immune System |
| AM | Ahuja and Murty |
| AM-H | Ahuja and Murty Tree Improvement Heuristic |
| BFB | Best Fit Bin |
| BF | Best Fit |
| BFD | Best Fit Decreasing |
| BFDH | Best Fit Decreasing Height |
| BRKGA | Biased Random-Key Genetic Algorithm |
| CG | Center of Gravity |
| CVC | California Vehicle Code |
| DD | Due Dates |
| DFTRC | Distance to the Front-Top-Right Corner |
| DG | Directed Graph |
| DJD | Djang and Finch |
| EA | Evolutionary Algorithm |
| FBSP | Finite Best Strip Packing |
| FF | First Fit |
| FFD | First Fit Decreasing |
| FFDH | First Fit Decreasing Height |
| FFF | Finite First Fit |
| FG-O-2DBPP | Free Guillotine Oriented Two-Dimensional Bin Packing Prob- |
|  | lem |
| FNF | Finite Next Fit |
| GA | Genetic Algorithm |
| GLS | Genetic Local Search |
| GRASP | Greedy Randomized Adaptive Search Procedure |
| HEA | Hybrid Evolutionary Algorithm |
| LB | Load Balancing |
| LGF | Lowest Gap Fill |
| LGFi | Improved Lowest Gap Fill |
| LGFof | Oriented Improved Lowest Gap Fill |
| LP | Linear Programming |
| MA | Memetic Algorithm |
|  |  |


| MH-MOGA | Multi-Heuristic Multi-Objective Genetic Algorithm |
| :--- | :--- |
| MH-SOMA | Multi-Heuristic Single Objective Memetic Algorithm |
| MHNO-MOGA | Multi-Heuristic Non-Oriented Multi-Objective Genetic Algorithm |
| MHNO-SOMA | Multi-Heuristic Non-Oriented Single Objective Memetic Algo- <br> rithm |
| MHO-MOGA | Multi-Heuristic Oriented Multi-Objective Genetic Algorithm |
| MHO-SOMA | Multi-Heuristic Oriented Single Objective Memetic Algorithm |
| MOGA | Multi-Objective Genetic Algorithm |
| MXGA | Multi-Crossover Genetic Algorithm |
| NFD | Next Fit Decreasing |
| NFDH | Next Fit Decreasing Height |
| NO-2DBP | Non-Oriented Two-Dimensional Bin Packing |
| NO-MOGA | Non-Oriented Multi-Objective Genetic Algorithm |
| NO-SOMA | Non-Oriented Single Objective Memetic Algorithm |
| O-2DBP | Oriented Two-Dimensional Bin Packing |
| O-BBP | Oriented Bin Packing Problem |
| O-MOGA | Oriented Multi-Objective Genetic Algorithm |
| O-SOMA | Oriented Single Objective Memetic Algorithm |
| OnBP | Online Bin Packing |
| OffBP | Offline Bin Packing |
| PCA | Principle Component Analysis |
| PSO | Particle Swarm Optimization |
| QAP | Quadratic Assignment Problem |
| RSAM | Random Sampling Ahuja and Murty |
| SR | Swap and Rotate |
| SA | Simulated Annealing |
| SO-2DBPP | Single Objective Two Dimensional Bin Packing Problem |
| SOMA | Single Objective Memetic Algorithm |
| TS | Tabu Search |
| TSP | Travelling Sales Person |
| UTS | Unified Tabu Search |

## CHAPTER 1

## INTRODUCTION

Given a set of rectangular items and a 2D bin of fixed width and variable length, the two-dimensional bin packing problem (2DBPP) consists of orthogonally placing all the pieces within the bin, without overlapping, such that the overall number of the bins is minimised [1, 2, 3, 4, 5]. The 2DBPP is an intractable optimization problem and widely faced during the industrial manufacturing processes. Textile manufacturing and newspaper page design are some of these optimization areas.

Online and offline are two different categories of 2DBPP according to the availability, at the beginning, of information about input. Online bin packing (OnBP) means that objects arrive one by one and there is no way to know complete input sequence, so it must be inserted into a bin immediately without waiting other objects. Hard drive partitioning for online storage systems usually deals with these types of problems. Offline bin packing (OffBP) means that all of the objects are known before they are packed so they can be reordered according to insertion heuristics. Plane load planning is a typical OffBPP.

Orientation is another key aspect of the 2DBPP. Rotating objects in packing creates better results, but objects may be not rotatable. Textile industry can change the orientation of single color shirts by rotating the shirts while the process is in cutting phase, because there is no difference between rotation or not. But shipping industry consider the orientation of fragile items.

In this thesis we study mechanisms, which are inspired by natural evolution, known as Evolutionary Algorithms (EA). Reproduction, mutation, recombination and selec-
tion are key mechanisms of EA that are used to solve optimization problems. Bin Packing, Travelling Sales Person and Quadratic Assignment Problem [5, 6] are wellknown problems that are solved with EAs. Genetic Algorithm (GA) and Memetic Algorithm (MA) are the most well-known approaches of EAs. The objective of these optimization problems can be single or multi-objective. GA mimics the natural evolution process and has the ability to find optimal solution in a large search space. In nature, survival of the fittest is a rule allowing the fittest solutions in each iteration to converge a (near-) optimal solution in a short time. In GA, the fittest of individual is the solution of optimization problems:

- Parents mate and produce offsprings
- Some of individuals are mutated
- Best individuals are selected to survive to the next generation

MA is another growing are of EA. It mimics natural evolution process but it may differ from GA by performing individual learning which is also known as meme(s). As in GA, MA applies below list to find (near-) optimal solution of optimization problems:

- Parents mate and produce offsprings
- Some of individuals are mutated
- Individual learning is performed (local search)
- Best individuals are selected to survive

In this thesis, we propose four different techniques to solve two different 2 D offline BPPs. In each technique, we use some of well known heuristics: Finite Next Fit (FNF), Finite First Fit (FFF), Best Fit Decreasing Height (BFDH), Unified Tabu Search(UTS)Improved Left Gap Fill (LGFi) and oriented Improved Left Gap Fill (LGFof).

The objective of our first problem is minimizing the number of bins. Orientation is crucial part of bin packing problems so we developed two different multi-heuristic single objective memetic algorithms (MH-SOMA). Our first algorithm, MHO-SOMA, uses heuristics: FNF, FFF, BFDH, UTS and LGFof. It tries to minimize the number
of bins for orientated 2D offline BPP. Our second algorithm, MHNO-SOMA, uses heuristics: FNF, FFF, BFDH, UTS, LGFof and LGFi in order to minimize the number of bins for non-orientated 2D offline BPP.

Our second optimization problem tries to find a pareto-optimal solution for both the minimal number of bins and the most efficiently load-balanced placing of the rectangular items. We propose two different multi-heuristic multi-objective genetic algorithms (MH-MOGA). The first proposed multi-objective technique, MHO-MOGA, uses heuristics: FNF, FFF, BFDH, UTS and LGFof to solve oriented multi objective 2D offline BPPs. The second proposed multi-objective technique, MHNO-MOGA, uses heuristics: FNF, FFF, BFDH, UTS, LGFof and LGFi. It tries to find paretooptimal solution (minimal number of bins and most effective load-balancing of items) for non-oriented multi objective 2D offline BPPs.

Through exhaustive experiments with 500 benchmark problem instances, we analyze the performance of our novel algorithms.

## CHAPTER 2

## RELATED WORK

In this chapter, we give information about the well known metaheuristics/heuristics such as Next Fit, First Fit, Best Fit, Unified Tabu Search, Left Gap Fill, Genetic Algorithm and Memetic Algorithm which are used to solve 2DBPP.

Next Fit is a level-oriented packing heuristic. It keeps a current level and tries to pack item into current level, if item does not fit then it creates a new level as current level and insert the item into current level.

First Fit is another level-oriented packing heuristic. It keeps all levels in all bins and tries to pack an item into first available level, if the item does not fit into any level then it creates new level and insert the item into new level

Best Fit is another level-oriented packing heuristic. It keeps all levels in all bins. First it calculates remaining gaps of all level even if we pack the item to that level and then tries to pack the item into minimum remaining gap level, if the item does not fit into any level then it creates new level and insert the item into new level.

Unified Tabu Search is a variant of Tabu Search Algorithm. First it packs the items into bins and then tries to replace items in order to minimize the number of bins. While minimizing the number of bins, it keeps a list known as tabu list and keeps track of last replacement movements.

Left Gap Fill tries to pack items into left most position of bin by calculating the best item which can fit into current gap.

Genetic Algorithm is an evolutionary algorithm. It tries to exploit the global search
space by mating best solution of problem. It uses inheritance, mutation, selection and crossover of natural evolution.

Memetic Algorithm is another evolutionary algorithm. It combines global search with an individual learning method.

### 2.1 Next Fit

Meir and Moser introduced a new packing algorithm in 1968 [32]. This algorithm packs the largest square cube into the left most corner of the hyberbox, and then it picks the second largest cube. It checks the remaining part of ground level of the hypercube. If the remaining part is enough then it inserts the second largest square cube into the position which it can touch the first largest cube. If the second largest cube cannot fit then it creates a new level with height of first largest cube and tries to insert the second largest cube at that level. This new introduced algorithm is known as Next Fit Decreasing (NFD) algorithm.

Coffman et al. studied on the NFD in order to pack the rectangles into rectangles or strips. This new extended algorithm is called Next Fit Decreasing Height (NFDH) algorithm [28].

Berkey and Wang enhanced Coffman's approach and proposed a new algorithm known ad Finite Next Fit (FNF). This level-oriened algorithm packs the item into finite bin set in one phase [35].

### 2.2 First Fit

Johnson et al. examined the good placement heuristics. One of the heuristics is known as First Fit (FF) [40]. In FF, items are packed into the first bin which has available space for item. First Fit Decreasing (FFD) algorithm is another variant of FF. It first sorts the items in descending order and then applies FF.

Coffman et al. applied non-increasing height sorting to FFD for 2DBPP and introduces a new algorithm First Fit Decreasing Height (FFDH) [28].

Berkey and Wang enhanced Coffman's first fit approach and proposed a new algorithm Finite First Fit (FFF) [35]. This level-oriened algorithm packs the item into finite bin set in one phase.

### 2.3 Best Fit

Johnson et al. described another good placement heuristics Best Fit (BF) algorithm [40]. It is similar to FF algorithm but items are inserted by remaining space in the bin. It calculates all the remaining spaces and inserts the item into bin which has smallest remaining space. Best Fit Decreasing (BFD) algorithm is an extended version of BF. First it sorts the items and then items are inserted.

Berkey and Wang enhanced BF and proposed a new algorithm Finite Best Strip Packing (FBSP) [35]. FBSP is a two level algorithm. At first level, items are inserted according to BF . At second level, prepared strip result is divided into levels and levels are merged into bins by using remaining horizontal space.

Coffman and Shor [23] described a new method similar to FBSP algorithm. It first sorts the items according to decreasing height and then inserts them by using BFD algorithm. This algorithm is Best Fit Decreasing Height (BFDH).

### 2.4 Unified Tabu Search

In 1986, Glover developed a new algorithm which is known as Tabu Search (TS) [19]. Glover described the details of algorithm [20, 21]. TS is composed of stopping condition, keeping tabu list, neighborhood search and aspiration condition. Stopping condition checks the endpoint of search. Tabu list keeps the track of last known moves in order to not to stuck in local optima. Neighborhood search is minor changes (swaps) of items. Aspiration condition or fitness function calculates the value of current search.

Lodi et al. developed new heuristics for different 2DBPP and used these heuristics within TS [11]. Lodi et al. explained the detail of their Unified Tabu Search (UTS)
algorithm [12] .

### 2.5 Left Gap Fill

Burke et al. introduced a new best fit strategy in order to pack 2D rectangles into strip [29]. New heuristic is using dynamic selection of best fit strategy in order to choose next rectangle. It used two stages processing of rectangles. First stage is preprocessing which rotates items whose height is greater than width and then sorts the item by decreasing width. Second stage it inserts items by using leftmost, tallest neighbor or shortest neighbor heuristics.

Bennell et al. enhanced the Burke strip packing strategy and used it to solve two dimensional bin packing problem [9]. It used the same preprocessing strategy and packing best fit rectangle into bin by using waste space insertion. This algorithm is known as Best Fit Bin (BFB). Lee changed BFB algorithm to Lowest Gap Fill (LGF) algorithm by using left most position search instead of waste space insertion [31]. Wong and Lee improved LGF into Improved Lowest Gap Fill (LGFi) by changing the shortest edge as remaining gap and also developed a new version of LGFi in order to pack oriented bin packing known as Oriented Improved Lowest Gap Fill (LGFof) [30].

### 2.6 Genetic Algorithms

Bremermann offers a mathematical model for solving problems by using gene (solutions), mutation and crossover [26]. This is the start point of Genetic Algorithms (GA). Holland extended GA and developed a framework for it [25]. This study made GA more applicable in computing enviroment. Goldberg introduced GA terminology and application areas [16]. Goldberg solves some problems in order to show the application of GA in a real problem. He also describes the implementation of the functionalities of genetic algorithm.

Smith tried to find the maximum number of rectangles that can be packed into a bin by using GA [14]. He used permutation of rectangles in order to encode the
chromosomes so the sequence of rectangles becomes important. The sequence gives the solution.

Hopper and Turton found solution to rotatable objects bin packing problems by use of BL and BLF heuristics [17]. He used order-based (permutation) encoding schema. Crossover and mutation generate valid chromosomes. There are two types of mutation operation. First one is order changing mutation and second one is rotation.

Kröger et al. creates a new vision of genetic algorithm encoding [13]. Directed Graph (DG) based encoding is used with top and right insertion technique. If a problem arises, then indices which shows priority are used to solve problem.

### 2.7 Recent Studies

Thapatsuwan et al. aim to find an optimal solution for multiple container packing problem [39]. Their study is composed of three different algorithms which are Artificial Immune System (AIS), GA and Particle Swarm Optimization (PSO). Result part of the study shows that according to found parameter sets AIS is better than GA and PSO but its run time is at least four times greater than GA.

Lim et al. encountered a practical problem California Vehicle Code (CVC) [10] . Containers are prepared by an oversea company and delivered in USA. Trucks has to obey axle weight constraints. A Greedy Randomized Adaptive Search Procedure (GRASP) is used in order to meet criteria.

Gonçalves and Resende presented a new way of 2DBP and 3DBP with Biased Random -Key Genetic Algorithm (BRKGA) [22]. He developed two different insertion heuristics which are based of Distance to the Front-Top-Right Corner (DFTRC) rule. DFTRC1 is calculating the distance according to the current presentation of box so it is available for oriented bin packing and DFTRC2 is calculating the distance for all the possible orientation of box so it is good to use it for orientation free problem.

Lopez-Camacho et al. suggested a heuristic and hyperheuristic method for packing problems [18]. Heuristics are FFD, Filler, BFD, Djang and Finch with initial fullness of DJD $\frac{1}{2}$, DJD $\frac{1}{3}$ and DJD $\frac{1}{2}$. Hyperheuristic finds the best combination of these
heuristic in order to solve one instance. GA is used to find hyperheuristics. Principle Component Analysis (PCA) is used to analyse BPPs.

Bennell et al. proposed a MultiCrossover Genetic Algorithm (MXGA) with BFB heuristic for NO-2DBPP with Due Dates (DD) [8]. BFB is a variant of BF heuristic and it has two phases. Phase one rotates the bins in order to maximize width and phase two sorts items with decreasing height. GA assigns the bin number of rectangles. Fitness function considers minimum number and minimum latency.

Fernández et al. considered rotation of items and Load Balancing (LB) of bins as multi-objective [4]. Multi-Objective Memetic Algorithm (MOMA) is developed. Insertion heuristic is made of two phase. First phase is using BLF heuristic and second phase is dividing bin into many bins in order to meet level criteria. Two types of local optimizers are used. First local optimizer tries to minimize the number of bins by trying the items from emptiest bin to insert to most occupied bin. Second local optimizer tries to gathers the items from most occupied zones of bins and insert them into new bins.

Blum and Schmid described an evolutionary algorithm (EA) for free guillotine oriented two-dimensional bin packing problem (FG-O-2DBPP) [3] . The proposed algorithm is based on LGFi heuristic. Crossover operator is a variant of uniform orderbased crossover.

Bansal and Khan used Round and Approximation Framework to pack 2D rectangles with and without rotation [46] . The proposed technique divides rectangles into five group: big, wide, long, medium and small. It creates containers with wide and long group of rectangles. First, it packs big, containers and medium by brute force approach. Then, it packs small rectangles by Next Fit.

Recent studies shows that bin packing problems do not have exact solutions. Recent optimization techniques such as GA, MA, PSO and AIS are tried to be applied for packing problems as soon as possible. Real life applications of BPP are not just about minimizing the number of bins so recent studies not only try to minimize the number of bins but also tries to keep balance of bins. Industrial engineering is curious about the solution of BPP.

### 2.8 Memetic Algorithms

Dawkins unfolded a new idea of gene-centered view of evolution. He coined a new phrase "meme" which means a behavior is passed from one individual to another individual. Moscato inspired by the book "The Selfish Gene" and developed a term Memetic Algorithm (MA) [37] . MA is also known as Hybrid Evolutionary Algorithms (HEA). It is a hybridization of population-based search and local search technique. In this study, he used MA as a combination of GA and Simulated Annealing (SA) for Travelling Sales Person (TSP) and some other problems. Moscato and Cotta described terms about MA and gave guidelines about the MA implementation [38] .

MA uses two different learning methods. First one is Lamarckian Learning. In Lamarckian Learning, individuals learn during their lifetime and keep the values of local search improvement. In selection mechanism, fitness of individuals are evaluated according to their local search improvements, so best improved fitnesses are transferred to next generation with their improvement information in genotype. Second one is Baldwinian Learning. In Baldwinian Learning, local search is applied to individual and fitness value of individual is calculated by the use of local learning, but local improvements are not part of genotype, so local search improvements are passed indirectly to next generation. Whitley et al. argued that Lamarckian Learning is faster than Baldwinian Learning but results of Lamarckian Learning can converge to local optima faster than Baldwinian Learning [15]. In 2005, Yao et al. showed that there is no significant difference between Lamarckian and Lamarckian Learnings [45] .

Development of MA can be divided into three generations. First generation describes the basis of Memetic Algorithm. It introduced Memetic Algorithm as marriage between a global search techique with local learning. Hyper-heuristic [24], Meta-Lamarckian MA [42] and Multi-meme [33] are classfied as second generation of MA. In Hyper-heuristic and Meta-Lamarckian MA, a reward mechanism is applied to individuals. Past values of memes are considered for reward value and individuals with higher reward values have more chances to survive and be replicated. In MultiMeme, memes are considered as a part of individual (genotype) so inheritance is used to passed memetic behavior of individuals to their offsprings. Third generation of MA
uses Co-evolution [27] and self-generation [34]. Rule-based local search is applied to heuristic so repeated patterns are captured and individuals are nourished by captured patterns.

Fischer and Merz developed a MA for cost-based communication problem. Recombination and mutation of population is proposed for global search and Ahuja and Murty (AM) tree improvement heuristic (AM-H) and the Random Sampling Ahuja and Murty (RSAM) algorithm are used for local search heuristics [44].

Merz and Freisleben proposed Genetic Local Search (GLS) algorithm for quadratic assignment problem (QAP). Proposed algorithm uses GA for global search and swapping items for local search [36].

Wolf and Merz described a new algorithm for supply chain problems. Linear programming (LP) is used as local search heuristic and mutation based filter checks offspring generation is applied for global search technique [43] .

## CHAPTER 3

## PROPOSED ALGORITHMS

2DBPP is an NP combinatorial problem. In daily life, humans encounter this problem directly or indirectly. Minimizing the textile usage for a manufacturer and planning the position of advertisements in a newspaper are well known problems. Some packing problems must be considered to solve with other objectives such as load balancing, due dates, etc

In this thesis, we tried to solve two different packing problems. For each problem, we developed two different algorithms.

Our first problem is minimizing the number of bins. We developed MHO-SOMA for oriented 2D offline BPP and MHNO-SOMA for non-oriented 2D offline BPP.

Minimizing number of bins and increasing load balance is our second problem. We proposed MHO-MOGA for oriented 2D offline BPP and MHNO-MOGA for nonoriented 2D offline BPP.

In first section, mathematical formulation of problems are described. In second section, we defined orientation in packing problems. In third section, we gave information about heuristics (FNF, FFF, BFDH, UTS, LGFof and LGFi) which are used in proposed algorithms. In fourth and fifth sections, we described details of our proposed algorithms (MHO-SOMA, MHNO-SOMA, MHO-MOGA and MHNO-MOGA).

### 3.1 Problem Definition

Load balancing for 2DBPP means that trying to balance total moments of rectangles which are on the left of centre of gravity (CG) of bins with total moments of rectangles which are on the right of CG of bins according to CG of bins. In our case, Euclidean Center of bin is considered as CG of bin. CG of bin and CG of rectangle is shown in Figure 3.1.


Figure 3.1: CG of bin and CG of rectangle

Two different load balancing of two different bins with same rectangle list is shown in Figure 3.2 and Bin2 has better load balance than Bin1.


Figure 3.2: Bin1 (left) and Bin2 (right), red dot for CG of bin, green dots for CG of rects and blue dot for CG of load

Definition of concept and mathematical formulation of single objective and multiobjective bin packing problem will be explained in two subsections.

### 3.1.1 Single Objective 2DBPP

Single objective two dimensional bin packing [39, 3] tries to minimize

$$
\begin{equation*}
C=\sum_{j=1}^{n} c_{j} \tag{3.1}
\end{equation*}
$$

subject to

$$
\begin{equation*}
x_{i}+\left(w_{i} w_{i}^{x}\right)+\left(h_{i} h_{i}^{x}\right) \leq x_{k}+\left(1-l e_{i k}\right) \quad, \forall i, k, i<k \tag{3.2}
\end{equation*}
$$

$$
\begin{equation*}
x_{k}+\left(w_{k} w_{k}^{x}\right)+\left(h_{k} h_{k}^{x}\right) \leq x_{i}+\left(1-r i_{i k}\right) \quad, \forall i, k, i<k \tag{3.3}
\end{equation*}
$$

$$
\begin{equation*}
y_{i}+\left(w_{i} w_{i}^{x}\right)+\left(h_{i} h_{i}^{x}\right) \leq y_{k}+\left(1-u n_{i k}\right) \quad, \forall i, k, i<k \tag{3.4}
\end{equation*}
$$

$$
\begin{equation*}
y_{k}+\left(w_{k} w_{k}^{x}\right)+\left(h_{k} h_{k}^{x}\right) \leq y_{i}+\left(1-a b_{i k}\right) \quad, \forall i, k, i<k \tag{3.5}
\end{equation*}
$$

$$
\begin{equation*}
l e_{i k}+r i_{i k}+u n_{i k}+a b_{i k} \leq p_{i j}+p_{k j}-1 \quad, \forall i, k, i<k \tag{3.6}
\end{equation*}
$$

$$
\begin{equation*}
\sum_{j=1}^{C} p_{i j}=1 \quad, \forall i \tag{3.7}
\end{equation*}
$$

$$
\begin{equation*}
\sum_{i=1}^{B} p_{i j} \leq M c_{j} \quad, \forall j \tag{3.8}
\end{equation*}
$$

$$
\begin{equation*}
x_{i}+\left(w_{i} w_{i}^{x}\right)+\left(h_{i} h_{i}^{x}\right) \leq W_{j}+\left(1-p_{i j}\right) M \quad, \forall i, j \tag{3.9}
\end{equation*}
$$

$$
\begin{equation*}
y_{i}+\left(w_{i} w_{i}^{y}\right)+\left(h_{i} h_{i}^{y}\right) \leq H_{j}+\left(1-p_{i j}\right) M \quad, \forall i, j \tag{3.10}
\end{equation*}
$$

$$
\begin{equation*}
w_{i}^{x}, w_{i}^{y}, h_{i}^{x}, h_{i}^{y}, l e_{i k}, r i_{i k}, a b_{i k}, u n_{i k}, p_{i j}, c_{j} \in 0,1 \quad, \forall i, k, i<k \tag{3.11}
\end{equation*}
$$

$$
\begin{equation*}
x_{i}, y_{i} \geq 0 \quad, \forall i \tag{3.12}
\end{equation*}
$$

Variables of formulation are described as
C total number of bin
$w_{i}, h_{i} \quad$ indicates width and height of rectangle i
$W_{j}, H_{j} \quad$ indicates width and height of bin j
$x_{i}, y_{i} \quad$ indicates left-bottom corner of rectangle i as coordinate
$w_{i}^{x}, w_{i}^{y} \quad$ binary variables indicate width of rectangle i is parallel to $X$ and $Y$ axis
$h_{i}^{x}, h_{i}^{y} \quad$ binary variables indicate height of rectangle i is parallel to X and Y axis
$l e_{i k} \quad$ a binary variable indicates rectangle $i$ is placed on the left side of rectangle $k$
$r i_{i k} \quad$ a binary variable indicates rectangle i is placed on the right side of rectangle $k$
$a b_{i k} \quad$ a binary variable indicates rectangle i is placed above rectangle k
$u n_{i k} \quad$ a binary variable indicates rectangle i is placed under rectangle k
$p_{i j} \quad$ a binary variable indicates; $p_{i j}=1$ if rectangle i is placed in bin j otherwise $p_{i j}=0$
$c_{j} \quad$ a binary variable indicate; $c_{j}=1$ if bin j is used otherwise $c_{j}=0$
M an arbitrarily large number used in Bin-M constraints

According to above definitions

- constraints (Eq. 3.2, 3.5) ensure that none of the rectangles overlaps each other
- constraint (Eq. 3.6) makes sure non-overlapping rectangles
- constraint (Eq. 3.7) guarantees that each rectangle is packed only in single bin
- constraint (Eq. 3.8) indicates that a bin is used iff a rectangle is packed into it
- constraints (Eq. 3.9-3.10) guarantee that all rectangles in the bin do not exceed the dimensions of the bin
- constraints (Eq. 3.11-3.12) describes the range of variables


### 3.1.2 Multi-Objective 2DBPP

Multi-objective two dimensional bin packing with load balancing [39, 3, 4] tries to minimize

$$
\begin{equation*}
(C / 2+L B / 2) \tag{3.13}
\end{equation*}
$$

when

$$
\begin{gather*}
C=\sum_{j=1}^{n} c_{j}  \tag{3.14}\\
L B=\sum_{j=1}^{C}\left|\sum_{i=1}^{B} p_{i j} d_{i} m_{i j} \sqrt{\left(x_{i j}+\left(w_{i j} / 2\right)-x_{C G}\right)^{2}+\left(y_{i j}+\left(h_{i j} / 2\right)-y_{C G}\right)^{2}}\right| \tag{3.15}
\end{gather*}
$$

subject to

$$
\begin{align*}
& x_{i}+\left(w_{i} w_{i}^{x}\right)+\left(h_{i} h_{i}^{x}\right) \leq x_{k}+\left(1-l e_{i k}\right) \quad, \forall i, k, i<k  \tag{3.16}\\
& x_{k}+\left(w_{k} w_{k}^{x}\right)+\left(h_{k} h_{k}^{x}\right) \leq x_{i}+\left(1-r i_{i k}\right) \quad, \forall i, k, i<k  \tag{3.17}\\
& y_{i}+\left(w_{i} w_{i}^{x}\right)+\left(h_{i} h_{i}^{x}\right) \leq y_{k}+\left(1-u n_{i k}\right) \quad, \forall i, k, i<k  \tag{3.18}\\
& y_{k}+\left(w_{k} w_{k}^{x}\right)+\left(h_{k} h_{k}^{x}\right) \leq y_{i}+\left(1-a b_{i k}\right) \quad, \forall i, k, i<k \tag{3.19}
\end{align*}
$$

$$
\begin{equation*}
l e_{i k}+r i_{i k}+u n_{i k}+a b_{i k} \leq p_{i j}+p_{k j}-1 \quad, \forall i, k, i<k \tag{3.20}
\end{equation*}
$$

$$
\begin{gather*}
\sum_{j=1}^{C} p_{i j}=1, \forall i \\
\sum_{i=1}^{B} p_{i j} \leq M c_{j}, \forall j \\
x_{i}+\left(w_{i} w_{i}^{x}\right)+\left(h_{i} h_{i}^{x}\right) \leq W_{j}+\left(1-p_{i j}\right) M \quad, \forall i, j \\
y_{i}+\left(w_{i} w_{i}^{y}\right)+\left(h_{i} h_{i}^{y}\right) \leq H_{j}+\left(1-p_{i j}\right) M \quad, \forall i, j \\
w_{i}^{x}, w_{i}^{y}, h_{i}^{x}, h_{i}^{y}, l e_{i k}, r i_{i k}, a b_{i k}, u n_{i k}, p_{i j}, c_{j} \in 0,1 \quad, \forall i, k, i<k \\
x_{i}, y_{i} \geq 0, \forall i \\
m_{i j} \in-1,1 \quad, \forall i \\
y_{C G} \geq(H / 2)
\end{gather*}
$$

Variables of formulation are described as

B total number of rectangle
C total number of bin
LB total sum of load balancing
$w_{i}, h_{i} \quad$ indicates width and height of rectangle i
$d_{i} \quad$ indicates weight of rectangle i
$W_{j}, H_{j} \quad$ indicates width and height of bin j
$x_{i}, y_{i} \quad$ indicates left-bottom corner of rectangle i as coordinate
$x_{C G} \quad$ indicates $x$ coordinate of center of gravity of bin which is equal to (W/2)
$x_{C G} \quad$ indicates y coordinate of center of gravity of bin which is equal to $(\mathrm{H} / 2)$
$w_{i}^{x}, w_{i}^{y} \quad$ binary variables indicate width of rectangle i is parallel to $X$ and $Y$ axis
$h_{i}^{x}, h_{i}^{y} \quad$ binary variables indicate height of rectangle i is parallel to $X$ and $Y$ axis
$l e_{i k} \quad$ a binary variable indicates rectangle $i$ is placed on the left side of rectangle k
$r i_{i k} \quad$ a binary variable indicates rectangle i is placed on the right side of rectangle k
$a b_{i k} \quad$ a binary variable indicates rectangle i is placed above rectangle k
$u n_{i k} \quad$ a binary variable indicates rectangle i is placed under rectangle k
$p_{i j} \quad$ a binary variable indicates; $p_{i j}=1$ if rectangle i is placed in bin j otherwise $p_{i j}=0$
$m_{i j} \quad$ a variable indicates; $m_{i j}=1$ if $\left(x_{i j}+\left(w_{i j} / 2\right)-x_{C G}\right) \geq 0$ otherwise $m_{i j}=-1$
$c_{j} \quad$ a binary variable indicates; $c_{j}=1$ if bin j is used otherwise $c_{j}=0$
M an arbitrarily large number used in Bin-M constraints

According to above definitions

- constraints (Eq. 3.16-3.19) ensure that none of the rectangles overlaps each other
- constraint (Eq. 3.20) makes sure non-overlapping rectangles
- constraint (Eq. 3.21) guarantees that each rectangle is packed only in single bin
- constraint (Eq. 3.22) indicates that a bin is used iff a rectangle is packed into it
- constraints (Eq. 3.23-3.24) guarantee that all rectangles in the bin do not exceed the dimensions of the bin
- constraints (Eq. 3.25 3.29) describes the range of variables


### 3.2 Orientation

Orientation means that an object can be rotated ( 90 angle) or not, while the object is being packed into bin. If an object of problem can be rotated while it is being packed into bin, then it is called as non-oriented or orientation-free. If an object of problem cannot be rotated, then it is called as oriented or orientation-fix. Therefore, if objects of problem can be rotated, then this problem is a non-oriented bin packing problem. Otherwise, the problem is an oriented bin packing problem. Textile industry can change the orientation of single color shirts by rotating the shirts while the process is in cutting phase, because there is no difference between rotation or not so it is a nonoriented bin packing problem. But shipping industry must consider the orientation of loads which are fragile, so it is an oriented packing bin problem.

### 3.3 State-of-the-art Algorithms

Insertion of items into bins is achieved by the use of heuristics. Each technique can result in a different arrangement of items so heuristic is one of the key part of BPP. Large numbers of methods have been proposed for bin packing problem. We choose six well known heuristics, which are FNF, FFF, BFDH, LGFi, LGFof and UTS, to pack items. In this study, FNF, FFF, BFDH, LGFi and LGFof heuristics are implemented and TSpack implementation of UTS is used.

### 3.3.1 Finite Next Fit

Berkey and Wang developed FNF algorithm [35]. The proposed FNF is described below in Algorithm 1 .

```
Algorithm 1 Finite Next Fit (FNF) algorithm
    function FNF(rectList, binWidht, binHeight)
        sort rectList by decreasing height
        currentBin \(\leftarrow\) CreateNewBin(binWidth, binHeight)
        currentLevel \(\leftarrow\) CreateNewLevel(currentBin)
        \(n \leftarrow\) rectList.length
        for \(i \leftarrow 1, n\) do
            if rect(i) fits into currentLevel then
                pack rect(i) into currentLevel
            else
                currentLevel \(\leftarrow\) CreateNewLevel(currentBin)
                if rect(i) fits into currentLevel then
                    pack rect(i) into currentLevel
                else
                currentBin \(\leftarrow\) CreateNewBin(binWidth, binHeight)
                    currentLevel \(\leftarrow\) CreateNewLevel(currentBin)
                end if
            end if
        end for
    end function
```

An example of proposed algorithm is shown in Table 3.1 and Figure 3.3
Table 3.1: Input List with Bin Width $=20$ and Bin Height $=20$

| Rect | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Height | 9 | 3 | 5 | 3 | 15 | 10 | 15 | 8 | 4 | 2 |
| Width | 10 | 8 | 6 | 10 | 2 | 6 | 4 | 3 | 6 | 12 |



Figure 3.3: Finite Next Fit Output

### 3.3.2 Finite First Fit

Berkey and Wang introduced FFF algorithm [35]. The proposed FFF is described below in Algorithm 2

```
Algorithm 2 Finite First Fit(FFF) algorithm
    function FFF(rectList, binWidht, binHeight)
        sort rectList by decreasing height
        \(n \leftarrow\) rectList.length
        for \(i \leftarrow 1, n\) do
            \(m \leftarrow\) bins.length
            for \(j \leftarrow 1, m\) do
                \(n \leftarrow\) bin(j).levels.length
            for \(k \leftarrow 1, n\) do
                    if rect(i) fits into bin(j).level(k) then
                                    pack rect(i) into bin(j).level(k)
                    end if
                end for
                end for
                if rect(i) does not fit into any bin then
                    create new bin
                    pack rect(i) into new bin
                end if
        end for
    end function
```

An example of proposed algorithm is shown in Table 3.1 and Figure 3.4


Figure 3.4: Finite First Fit Output

### 3.3.3 Best Fit Decreasing Height

Coffman and Shor described BFDH algorithm [23]. The proposed BFDH is described below in Algorithm 3 .

```
Algorithm 3 Best Fit Decreasing Height(BFDH) algorithm
    function BFDH (rectList, binWidht, binHeight)
        sort rectList by decreasing height
        \(n \leftarrow\) rectList.length
        for \(i \leftarrow 1, n\) do
            minRemainingS pace \(\leftarrow\) binWidth
            \(m \leftarrow\) bins.length
            for \(j \leftarrow 1, m\) do
            \(n \leftarrow\) bin(j).levels.length
            for \(k \leftarrow 1, n\) do
                if rect(i) fits into \(\operatorname{bin}(\mathrm{j})\).level( k ) then
                                    pack rect(i) into bin(j).level(k)
                                    if remainingS pace \(<\) minRemaningS pace then
                    minRemainingS pace \(\leftarrow\) remainingS pace
                end if
                end if
                end for
            end for
            if rect(i) fits fit into any bin then
                    packs the item into minRemainingSpace bin level
            else
                    create new bin
                    pack rect(i) into new bin
            end if
        end for
    end function
```

An example of proposed algorithm is shown in Table 3.1 and Figure 3.5:


Figure 3.5: Best Fit Decreasing Height Output

### 3.3.4 Unified Tabu Search

Lodi et al explained the detail of their UTS and implemented the algorithm [12]. UTS algorithm of Lodi is explained below in Algorithm 4

```
Algorithm 4 Unified Tabu Search(UTS) algorithm
    function TSpack(rectList, binWidht, binHeight)
        \(z * \leftarrow \mathrm{~A}(\mathrm{f} 1 ;::: ; \mathrm{ng})\)
        \(L \leftarrow\) alowerboundontheoptimalsolutionvalue
        if \(z *=L\) then
            stop
        end if
        initialize all tabu lists to empty
        pack each item into a separate bin
        \(z \leftarrow n \quad \triangleright\) Tabu Search solution value
        \(d \leftarrow 1 ;\)
        determine the target bin t
        while dotimelimitisnotreached
            diversity \(\leftarrow\) false
            \(k \leftarrow 1\)
            while dodiversify \(=\) falseandz* \(>L\)
                \(k i n \leftarrow l \operatorname{SEARCH}(t, k\), diversify,\(z)\)
                \(z \leftarrow \min (z, z *)\)
                if thenk <= kin
                    determine the new target bin \(t\)
                    end if
            end while
            if \(z *=L\) then
                stop
            elseDIVERSIFICATION \((d, z, t)\)
            end if
        end while
    end function
```

An example of proposed algorithm is shown in Table 3.1 and Figure 3.6:


Figure 3.6: Unified Tabu Search Output

### 3.3.5 Improved Left Gap Fill

Wong and Lee improved LGF into LGFi [30]. LGFi is used for non-oriented bin packing problem. The proposed LGFi is described below in Algorithm 5 .

```
Algorithm 5 Improved Left Gap Fill Non-Oriented(LGFi) algorithm
    function LGFI(rectList, binWidht, binHeight)
        rotate rects to make width greater than height
        sort rectList by decreasing height
        currentBin \(\leftarrow\) CreateNewBin(binWidth, binHeight)
        currentPoint \((x, y) \leftarrow 0,0\)
        while rectList.length \(!=0\) do
            \(r e c t \leftarrow\) GetNextRectangle(rectList)
            gapWidth \(\leftarrow\) binWidth - currentPoint X
            gapHeight \(\leftarrow\) binHeight - currentPoint \(Y\)
            gap \(\leftarrow \min (g a p W i d t h\), gapHeight)
            if \(g a p=0\) then
                currentBin \(\leftarrow\) CreateNewBin(binWidth, binHeight)
                currentPoint \((x, y) \leftarrow 0,0\)
                else if gap \(<\min (\) allRemainingRectHeight \()\) then
                    mark are as waste space
                    set currentPoint( \(\mathrm{x}, \mathrm{y}\) ) to lowest-leftmost point
            else if gap \(=\) rect.heightorgap \(=\) rect.width then
                    pack rect into currentPoint( \(\mathrm{x}, \mathrm{y}\) )
                set currentPoint( \(x, y\) ) to lowest-leftmost point
                    remove rect from rectList
            else if norect fitintogap then
                \(r e c t \leftarrow \operatorname{GetNextRectangle(rectList)~}\)
                pack rect into currentPoint( \(\mathrm{x}, \mathrm{y}\) )
                set currentPoint( \(\mathrm{x}, \mathrm{y}\) ) to lowest-leftmost point
                remove rect from rectList
                end if
        end while
    end function
```

An example of proposed algorithm is shown in Table 3.1 and Figure 3.7:


Figure 3.7: LGFi Output

### 3.3.6 Oriented Improved Left Gap Fill

Wong and Lee improved LGFi for O-BPP and developed LGFof [30]. The proposed LGFof is described below in Algorithm6.

```
Algorithm 6 Improved Left Gap Fill Oriented(LGFof) algorithm
    function LGFof(rectList, binWidht, binHeight)
        sort rectList by non-increasing area
        currentBin \(\leftarrow\) CreateNewBin(binWidth, binHeight)
        currentPoint \((x, y) \leftarrow 0,0\)
        while rectList.length \(!=0\) do
            \(r e c t \leftarrow\) GetNextRectangle(rectList)
            gapWidth \(\leftarrow\) binWidth - currentPoint X
            gapHeight \(\leftarrow\) binHeight - currentPoint \(Y\)
            gap \(\leftarrow \min (\) gapWidth, gapHeight)
            if \(g a p=0\) then
                currentBin \(\leftarrow\) CreateNewBin(binWidth, binHeight)
                currentPoint \((x, y) \leftarrow 0,0\)
            else if gap \(<\min (\) allRemainingRectHeight \()\) then
                    mark are as waste space
                    set currentPoint( \(\mathrm{x}, \mathrm{y}\) ) to lowest-leftmost point
            else if gap \(=\) rect.heightorgap \(=\) rect.width then
                pack rect into currentPoint( \(\mathrm{x}, \mathrm{y}\) )
                set currentPoint( \(\mathrm{x}, \mathrm{y}\) ) to lowest-leftmost point
                remove rect from rectList
            else if norect fitintogap then
                \(r e c t \leftarrow \operatorname{GetNextRectangle}(\) rectList \()\)
                pack rect into currentPoint( \(\mathrm{x}, \mathrm{y}\) )
                set currentPoint( \(\mathrm{x}, \mathrm{y}\) ) to lowest-leftmost point
                remove rect from rectList
            end if
        end while
    end function
```

An example of proposed algorithm is shown in Table 3.1 and Figure 3.8:


Figure 3.8: LGFof Output

### 3.3.7 Genetic Algorithm

GA mimics the behavior of the natural evolution process. GA has got the power of natural evolution process in order to find optimum solution for a large search space. In GAs, a chromosome means a possible solution for the given problem set. A list of chromosome generates the population of GA. Each chromosome consists of genes and has an, encoding schema. Heuristic can be part of a chromosome or genetic algorithm. Each item of the problem means a gene in the chromosome. Encoding keeps the identification of items. Heuristic is used to find the fitness value of the chromosome.

GA has three operators:

- Selection: survival of the individual
- Crossover: mating the parents in order to get offspring
- Mutation: random modification of individual

GA works as follows in Algorithm 7.

```
Algorithm 7 Genetic Algorithm (GA) algorithm
    Initialize population
    Calculate fitness of population
    while Best individual (chromosome) does not meet the criteria do
        Select parents
        Crossover operation in order to get offsprings
        Mutation operation over population
        Determine fitness of population
    end while
    Best individual is the solution of problem
```


### 3.3.8 Chromosome Structure

Chromosome is a candidate solution of BPP. For 2DBPP, rectangular items are the genes of a chromosome. In this study, chromosome is not only composed of genes also it has two heuristics to pack the genes into bins. We have used permutation encoding to keep the identification of rectangles. Permutation encoding is a form of keeping width-height of rectangular items and sequence between rectangles.

Chromosome Structure of 2DBPP Explanation of structure is in Figure 3.9

- W: width of rectangle
- H : height of rectangle:
- Heu1: heuristic one
- Heu2: heuristic two


Figure 3.9: Chromosome Structure

An example of chromosome is shown Figure 3.10:

| 1,2 | 2,5 | 4,3 | 3,1 | FFF | FFF |
| :--- | :--- | :--- | :--- | :--- | :--- |

Figure 3.10: Example Chromosome

### 3.3.9 Selection

We have used Elitism as the selection of parents in GA. Elitism means that best chromosomes in the population have the chance to mate in order to create new offsprings.

### 3.3.10 Crossover

Single point crossover is applied technique in our study. One point is chosen. From beginning of the chromosome to the crossover point is copied from first parent and the rest of the chromosome comes from the second parent. Our crossover point is half of the chromosome. We also get first heuristic from one parent and second heuristic from other parent.

An example of crossover is shown in Figure 3.11 and Figure 3.12:


Figure 3.11: Single Point Crossover Chromosome Structure

Parent 1


Offspring1
Figure 3.12: Single Point Crossover Example

### 3.3.11 Mutation

Mutation ensures the diversity of population. Invalid genes can be generated as a result of mutation operation. Mutation is only done in the concept of rectangular items and heuristics are not mutated. In this study, we do not want to create invalid genes so three types of mutation operations are implemented:

- Swap : Only swaps selected genes
- Rotate : Only rotate object, the value of width and height is swapped
- Swap and Rotate : Selected genes are swapped and rotated

An example of swap mutation is shown in Figure 3.13 .

## Chromosome



## Swapped Genes

Figure 3.13: Swap Mutation Example

An example of rotation mutation is shown in Figure 3.14.

## Chromosome



Selected Gene


Rotated Genes
Figure 3.14: Rotation Mutation Example

An example of swap-rotation mutation is shown in Figure 3.15:

## Chromosome



## Swapped-Rotated Genes

Figure 3.15: Swap-Rotation Mutation Example

### 3.3.12 Memetic Algorithm

MA is another growing are of EA. It also mimic natural evolution process but it may differ from GA by performing individual learning which is also known as meme(s). In this study, MA uses the same chromosome structure, selection, crossover and mutation operation of GA.

MA works as follows in Algorithm 8 .

```
Algorithm 8 Memetic Algorithm (MA) algorithm
    Initialize population
    Calculate fitness of population
    while Best individual (chromosome) does not meet the criteria do
        Select parents
        Crossover operation in order to get offsprings
        Mutation operation over population
        Individual learning is proceed (local search)
        Determine fitness of population
    end while
    Best individual is the solution of problem
```


### 3.4 Single Objective Proposed Algorithms

Bin packing problems mostly consider the number of bins so the goal of the algorithms is to pack items in minimum number of bins. We have developed a two MultiHeuristic Single Objective Memetic Algorithms for minimizing the number of bins by considering orientation is important or not. Our fitness functions are the result of applied heuristic in the term of number of bins. We explain two algorithms in two different subsections.

### 3.4.1 MHO-SOMA for O-2DBPP

Minimizing the number of bins is our goal but orientation of items must be conserved for some 2D offline BPP. Advertisement packing for a newspaper can be classified as one example of these problems. To keep orientation of items, we developed MultiHeuristic Oriented Single Objective Memetic Algorithm (MHO-SOMA). FNF, FFF, BFDH and LGFof are used as heuristics to pack items. First, genetic algorithm with heuristics is run by the use of swap mutation operation. Then, best individual of genetic algorithm is selected and UTS is applied to best individual as local search. Overall best result is the solution of oriented 2D offline BPP. MHO-SOMA is explained in Algorithm 9 and Algorithm 10 .

```
Algorithm 9 Multi-Heuristic Oriented Single Objective Memetic Algorithm Func-
tions
    function \(\operatorname{Initalize}(\) rectList \(, n, k)\)
        for \(i \leftarrow 1, n\) do
            copyList \(\leftarrow\) rectList
            for \(j \leftarrow 1, k\) do
                    ApplySwapMutation(copyList)
            end for
            heuristic \(\leftarrow\) СноозеOne(FNF, FFF, BFDH, LGFof)
            heuristic \(1 \leftarrow\) heuristic
            heuristic \(2 \leftarrow\) heuristic
            CreateСhromosome(copyList,heuristic1, heuristic2)
            PutPopulation(chromosome)
        end for
        return population
    end function
    function Fitness(population)
        for eachchromosomeinPopulation do
            if heuristic \(1=\) Heuristic 2 then
                    fitness \(\leftarrow\) Insert(heuristic1, rectList)
            else
                    resHeuristic \(1 \leftarrow\) Insert(heuristic 1, firstHalfRectList)
                    resHeuristic \(2 \leftarrow \operatorname{InsERT}(\) heuristic 2 , secpmdHalfRectList)
                    fitness \(\leftarrow\) resHeuristic \(1+\) reHeuristic 2
            end if
        end for
        SortChromosomesaccordingtotheirfitnessvalue
        return population
    end function
```

```
Algorithm 10 Multi-Heuristic Oriented Single Objective Memetic Algorithm (MHO-
SOMA) algorithm
    function MHO-SOMA(rectList, binW, binH, numGen, numPop, numIMut)
        population \(\leftarrow\) Initialize(rectList, numPop, numIMut)
        Fitness(population)
        for \(k \leftarrow 1\), numGen do
            parents \(\leftarrow\) ElitismSelectParents(population)
            off springs \(\leftarrow\) SinglePointCrossover(parents)
            SwapMutation(population)
            Fitness(population)
        end for
        rectSequence \(\leftarrow\) bestindividualrectanglesequence
        fitnessBest \(\leftarrow\) bestindividualfitness
        fitnessUTS \(\leftarrow \operatorname{Insert(UTS,~rectSequence)~}\)
        if fitnessUTS < fitnessBest then
            resultHeuristic \(\leftarrow U T S\)
            resultNumberOfBin \(\leftarrow\) fitnessUTS
        else
            resultHeuristic \(\leftarrow\) bestindividualheuristic
            resultNumberOfBin \(\leftarrow\) fitnessBest
        end if
        return resultHeuristic, resultNumberOf Bin
    end function
```


### 3.4.2 MHNO-SOMA for NO-2DBPP

Orientation of items is not considered for some 2D offline BPP. Textile usage of manufacturer can be classified as one example of these problems. We proposed MultiHeuristic Non-Oriented Single Objective Memetic Algorithm (MHNO-SOMA) for these problems. Rotation and swap-rotation mutations are used to change orientation of items. First we run genetic algorithm with heuristics: FNF, FFF, BFDH and LGFof by the use of rotation and swap-rotation mutations. Then, best chromosome
of genetic algorithm is selected to apply UTS and LGFi. Overall best result is the solution of non-oriented 2D offline BPP. MHNO-SOMA is explained in Algorithm 11 and Algorithm 12 .

```
Algorithm 11 Multi-Heuristic Non-Oriented Single Objective Memetic Algorithm
Functions
    function \(\operatorname{Intalize}(\) rect List, \(n, k\) )
        for \(i \leftarrow 1, n \mathbf{d o}\)
        copyList \(\leftarrow\) rectList
        for \(j \leftarrow 1, k\) do
            RotationMutation(copyList) or SwapRotationMutation(copyList)
        end for
        heuristic \(\leftarrow\) СноозеOne(FNF, FFF, BFDH, LGFof)
        heuristic \(1 \leftarrow\) heuristic
        heuristic \(2 \leftarrow\) heuristic
        CreateChromosome(copyList, heuristic 1 , heuristic 2 )
        PutPopulation(chromosome)
        end for
        return population
    end function
    function Fitness(population)
        for eachchromosomeinPopulation do
        if heuristic \(1=\) Heuristic 2 then
            fitness \(\leftarrow\) Insert(heuristic1, rectList)
        else
            resHeuristic \(1 \leftarrow\) Insert(heuristic 1 , firstHalfRectList)
            resHeuristic \(2 \leftarrow \operatorname{Insert}\) (heuristic 2 , secpmdHalfRectList)
            fitness \(\leftarrow\) resHeuristic \(1+\) reHeuristic 2
        end if
        end for
        SortChromosomesaccordingtotheirfitnessvalue
        return population
    end function
```

```
Algorithm 12 Multi-Heuristic Non-Oriented Single Objective Memetic Algorithm
(MHNO-SOMA) algorithm
    function MHNO-SOMA(rectList, binW, binH, numGen, numPop, numIMut)
        population \(\leftarrow \operatorname{Initialize}(r e c t L i s t\), numPop, numIMut)
        Fitness(population)
        for \(k \leftarrow 1\), numGen do
            parents \(\leftarrow\) ElitismSelectParents(population)
            off springs \(\leftarrow\) SinglePointCrossover(parents)
            RotateMutation(population) or SwapRotateMutation(population)
            Fitness(population)
        end for
        rectSequence \(\leftarrow\) bestindividualrectanglesequence
        fitnessBest \(\leftarrow\) bestindividualfitness
        fitnessUTS \(\leftarrow\) Insert(UTS, rectSequence)
        fitnessLGFi \(\leftarrow \operatorname{Insert}(\mathrm{LGFi}\), rectSequence)
```



```
        resultHeuristic \(\leftarrow\) resultNumberOfBinHeuristic
        return resultHeuristic, resultNumberOf Bin
    end function
```


### 3.5 Multi-Objective Proposed Algorithms

Load balancing of bins is an important issue so the goal of the algorithms is not only to pack items in minimum number of bins but also maximize the load balancing of bins. We have developed two Multi-Heuristic Multi-Objective Genetic Algorithms for minimizing number of bins and maximizing load balance.

Our fitness function is composed of two parts, fitness of number of bins and fitness of center of gravity. Number of bins is calculated directly by heuristic. Our algorithm is not directly consider load balancing so we add calculation of center of gravity to each bin. In order to calculate center of gravity, each bin's center point is selected as Center of Gravity (CG). When a rectangle box inserted, then the Euclidean distance of its center to bin's CG is calculated and multiplied by the weight of rectangle. This
calculated value is CG of rectangle. If sign of $x$ coordinate of rectangle is minus then CG of rectangle is subtracted from total CG of Bin. If sign of $x$ coordinate of rectangle is plus then CG of rectangle is added to total CG of Bin . When all rectangles are inserted to bins, then absolute values of total CG of bins are added. This calculated value is called as CG of Chromosome. Load balance of bin is explained in Equation 3.30 and load balance of chromosome is explained in Equation 3.31 .

$$
\begin{gather*}
L B_{\text {Bin }}=\sum_{j=1}^{\# \text { Rect }} d_{i j} m_{i j} \sqrt{\left(x_{i j}+\left(w_{i j} / 2\right)-x_{C G}\right)^{2}+\left(y_{i j}+\left(h_{i j} / 2\right)-y_{C G}\right)^{2}}  \tag{3.30}\\
L B_{\text {Chromosome }}=\sum_{i=1}^{\# B i n}\left|\sum_{j=1}^{\# \text { Rect }} d_{j} m_{i j} \sqrt{\left(x_{i j}+\left(w_{i j} / 2\right)-x_{C G}\right)^{2}+\left(y_{i j}+\left(y_{i j} / 2\right)-y_{C G}\right)^{2}}\right| \tag{3.31}
\end{gather*}
$$

Multi-objective 2D offline BPP can be classified into two categories according to rotation. We developed two algorithms: one for oriented and another for non-oriented problems. Two proposed multi-objective algorithms are explained in two different subsections.

### 3.5.1 MHO-MOGA for O-2DBPP with LB

Multi-Heuristic Oriented Multi-Objective Genetic Algorithm (MHO-MOGA) is proposed to solve oriented multi-objective 2D offline BPP. FNF, FFF, BFDH, LGFof and UTS are applied as heuristics on the base of genetic algorithm. Swap mutation operation is used to keep orientation of items. Best result of GA is the result of problem. MHO-MOGA is explained in Algorithm 13, Algorithm 14 and Algorithm 15 ,

```
Algorithm 13 Multi-Heuristic Oriented Multi-Objective Genetic Algorithm Func-
tions
    function Initalize(rectList, \(n, k\) )
        for \(i \leftarrow 1, n\) do
            copyList \(\leftarrow\) rectList
            for \(j \leftarrow 1, k\) do
                SwapMutation(copyList)
            end for
            heuristic \(\leftarrow\) СнооseOne(FNF, FFF, BFDH, LGFof, UTS)
            heuristic \(1 \leftarrow\) heuristic
            heuristic \(2 \leftarrow\) heuristic
            CreateChromosome(copyList, heuristic 1 , heuristic 2 )
            PutPopulation(chromosome)
        end for
        return population
    end function
```

```
Algorithm 14 Multi-Heuristic Oriented Multi-Objective Genetic Algorithm Func-
tions Part2
    function Fitness(population)
        \(\operatorname{maxBin} \leftarrow 0\)
        minBin \(\leftarrow\) MAXINTEGER
        \(\max C G \leftarrow 0\)
        \(\min C G \leftarrow M A X I N T E G E R\)
        for eachchromosomeinPopulation do
            if heuristic \(1=\) Heuristic 2 then
                chromosomeBin \(\leftarrow \operatorname{INSERT}(\) heuristic 1 , rectList)
            else
                resHeuristic \(1 \leftarrow\) Insert(heuristic 1, firstHalfRectList)
                resHeuristic \(2 \leftarrow\) Insert(heuristic2, secpmdHalfRectList)
                chromosomeBin \(\leftarrow\) resHeuristic \(1+\) reHeuristic 2
            end if
            if chromosomeBin < minBin then
                minBin \(\leftarrow\) chromosomeBin
            end if
            if chromosomeBin \(>\) maxBin then
                maxBin \(\leftarrow\) chromosomeBin
            end if
            chromosome \(C G \leftarrow \operatorname{Sum}(\) bin of chromosome)
        end for
        diffBin \(\leftarrow\) maxBin \(-\operatorname{minBin}\)
        diffCG \(\leftarrow \max C G-\min C G\)
        for eachchromosomeinPopulation do
            fitnessBin \(\leftarrow\) (chromosomeBin - minBin)/DiffBin
            Fitness \(C G \leftarrow(\) chromosomeCG \(-\operatorname{minCG}) / D i f f C G\)
            fitness \(\leftarrow(\) FitnessBin \(* 0.5)+(\) Fitness \(C G * 0.5)\)
        end for
        SortChromosomesaccordingtotheirfitnessvalue
        return population
    end function
```

```
Algorithm 15 Multi-Heuristic Oriented Multi-Objective Genetic Algorithm Oriented
(MHO-MOGA) algorithm
    function MHO-MOGA(rectList, binW, binH, numGen, numPop, numIMut)
        population \(\leftarrow\) Initialize(rectList, numPop, numIMut)
        Fitness(population)
        for \(k \leftarrow 1\), numGen do
            parents \(\leftarrow\) ElitismSelectParents(population)
            off springs \(\leftarrow\) SinglePointCrossover(parents)
            SwapMutation(population)
            Fitness(population)
        end for
        resultHeuristic \(\leftarrow\) bestindividualheuristic
        resultNumberOf Bin \(\leftarrow\) bestindividualfitness
        return resultHeuristic, resultNumberOf Bin
    end function
```


### 3.5.2 MHNO-MOGA for NO-2DBPP with LB

Sometimes, orientation is not part of problem so in order to solve such problems, we developed Multi-Heuristic Non-Oriented Multi-Objective Genetic Algorithm (MHNOMOGA). Each individual uses one or two of the heuristics: FNF, FFF, BFDH, UTS, LGFof and LFGi to pack rectangles into bins. At the beginning of genetic algorithm, each indivudual picks only one of the heuristics. In the next phases of GA, an individual can have two different heuristics. Each heuristic is applied to corresponding part of rectangle list. Rotation mutation and swap-rotation mutation are applied to change orientation of rectangles. Best result of GA gives the solution of problem. MHNO-MOGA is explained in Algorithm 16, Algorithm 17 and Algorithm 18

```
Algorithm 16 Multi-Heuristic Non-Oriented Multi-Objective Genetic Algorithm
Functions
    function \(\operatorname{Initalize}(\) rectList \(, n, k)\)
        for \(i \leftarrow 1, n\) do
            copyList \(\leftarrow\) rectList
            for \(j \leftarrow 1, k\) do
                    RotationMutation(S wapMutation) or SwapRotationMuta-
    tion(S wapMutation)
                end for
                heuristic \(\leftarrow\) ChоoseOne(FNF, FFF, BFDH, LGFof, UTS, LGFi)
                heuristic \(1 \leftarrow\) heuristic
                heuristic \(2 \leftarrow\) heuristic
                CreateChromosome(copyList, heuristic1, heuristic2)
                PutPopulation(chromosome)
        end for
        return population
    end function
```

```
Algorithm 17 Multi-Heuristic Non-Oriented Multi-Objective Genetic Algorithm
Functions Part2
    function Fitness(population)
    maxBin \(\leftarrow 0\)
        \(\operatorname{minBin} \leftarrow M A X I N T E G E R\)
        \(\max C G \leftarrow 0\)
        \(\min C G \leftarrow M A X I N T E G E R\)
        for eachchromosomeinPopulation do
            if heuristic \(1=\) Heuristic 2 then
                    chromosomeBin \(\leftarrow \operatorname{InSERT}(\) heuristic 1 , rectList)
        else
            resHeuristic \(1 \leftarrow\) Insert(heuristic 1, firstHalfRectList)
            resHeuristic \(2 \leftarrow \mathrm{Insert}\) (heuristic2, secpmdHalfRectList)
            chromosomeBin \(\leftarrow\) resHeuristic \(1+\) reHeuristic 2
        end if
            if chromosomeBin \(<\operatorname{minBin}\) then
                \(\operatorname{minBin} \leftarrow\) chromosomeBin
            end if
            if chromosome Bin \(>\operatorname{maxBin}\) then
                    maxBin \(\leftarrow\) chromosomeBin
            end if
            chromosome \(C G \leftarrow \operatorname{Sum}(\) bin of chromosome)
        end for
        diffBin \(\leftarrow\) maxBin - minBin
        diffCG \(\leftarrow \max C G-\min C G\)
        for eachchromosomeinPopulation do
            fitnessBin \(\leftarrow(\) chromosomeBin - minBin)/DiffBin
            FitnessCG \(\leftarrow(\) chromosomeCG - minCG)/DiffCG
            fitness \(\leftarrow(\) FitnessBin \(* 0.5)+(\) Fitness \(C G * 0.5)\)
        end for
        SortChromosomesaccordingtotheirfitnessvalue
        return population
    end function
```

```
Algorithm 18 Multi-Heuristic Non-Oriented Multi-Objective Genetic Algorithm
(MHNO-MOGA) algorithm
    function MHNO-MOGA(rectList, binW, binH, numGen, numPop, numIMut)
        population \(\leftarrow \operatorname{Intitialize}(\) rectList, numPop, numIMut)
        Fitness(population)
        for \(k \leftarrow 1\), numGen do
            parents \(\leftarrow\) ElitismSelectParents(population)
            off springs \(\leftarrow\) SinglePointCrossover(parents)
            RotationMutation(population) or SwapRotationMutation(population)
            Fitness(population)
        end for
        resultHeuristic \(\leftarrow\) bestindividualheuristic
        resultNumberOfBin \(\leftarrow\) bestindividualfitness
        return resultHeuristic, resultNumberOf Bin
    end function
```


## CHAPTER 4

## EXPERIMENTAL RESULTS

In this Chapter, we present the results of proposed algorithms. All the algorithms are implemented in C++ language and tests are executed under Windows7 OS. The computer, which is used for tests, has 2.10 GHz Intel Core 2 Duo CPU with usable 3 GB RAM.

### 4.1 Problem Sets

In this study, we tried to solve 2D offline BPPs. For test cases we used two well known problem instance sets. First one is Berkey-Wang classes [35]. They generated six classes. Their classses structures are as follows:

- Class 1 : $\mathrm{w}_{\mathrm{j}}$ and $\mathrm{h}_{\mathrm{j}}$ uniformly random in $[1,10], \mathrm{W}=\mathrm{H}=10$;
- Class 2 : $\mathrm{w}_{\mathrm{j}}$ and $\mathrm{h}_{\mathrm{j}}$ uniformly random in $[1,10], \mathrm{W}=\mathrm{H}=30$;
- Class $3: \mathrm{w}_{\mathrm{j}}$ and $\mathrm{h}_{\mathrm{j}}$ uniformly random in [1,35], $\mathrm{W}=\mathrm{H}=40$;
- Class 4 : $\mathrm{w}_{\mathrm{j}}$ and $\mathrm{h}_{\mathrm{j}}$ uniformly random in $[1,35], \mathrm{W}=\mathrm{H}=100$;
- Class 5 : $\mathrm{w}_{\mathrm{j}}$ and $\mathrm{h}_{\mathrm{j}}$ uniformly random in $[1,100], \mathrm{W}=\mathrm{H}=100$;
- Class $6: w_{j}$ and $h_{j}$ uniformly random in $[1,100], W=H=300$.

Second instance set is composed of four classes which are generated by Martello and Vigo [41]. Their proposed type structures is :

- Type $1: \mathrm{w}_{\mathrm{j}}$ uniformly random in $\left[\frac{2}{3} \mathrm{~W} ; \mathrm{W}\right], \mathrm{h}_{\mathrm{j}}$ uniformly random in $\left[1 ; \frac{1}{2} \mathrm{H}\right]$;
- Type $2: \mathrm{w}_{\mathrm{j}}$ uniformly random in $\left[1 ; \frac{1}{2} \mathrm{~W}\right], \mathrm{h}_{\mathrm{j}}$ uniformly random in $\left[\frac{2}{3} \mathrm{H} ; \mathrm{H}\right]$;
- Type $3: w_{j}$ uniformly random in $\left[\frac{1}{2} W ; W\right], h_{j}$ uniformly random in $\left[\frac{1}{2} H ; H\right]$;
- Type $4: \mathrm{w}_{\mathrm{j}}$ uniformly random in $\left[1 ; \frac{1}{2} \mathrm{~W}\right], \mathrm{h}_{\mathrm{j}}$ uniformly random in $\left[1 ; \frac{1}{2} \mathrm{H}\right]$.

Martello-Vigo class structures is using $\mathrm{W}=100$ and $\mathrm{H}=100$ and listed below:

- Class 7 : type 1 with probability $70 \%$, type $2,3,4$ with probability $10 \%$ each;
- Class 8 : type 2 with probability $70 \%$, type $1,3,4$ with probability $10 \%$ each;
- Class 9 : type 3 with probability $70 \%$, type $1,2,4$ with probability $10 \%$ each;
- Class 10 : type 4 with probability $70 \%$, type $1,2,3$ with probability $10 \%$ each. In this study, we have four different problems as listed below:
- Single Objective O-2DBPP
- Single Objective NO-2DBPP
- Multi-Objective O-2DBPP with LB
- Multi-Objective NO-2DBPP with LB

In order to solve these problems, we need to find the population size and generation number. We listed Berkey-Wang and Martello-Vigo instances and picked a random number between 1 and 500. The number is 245 and its properties are listed below :

- a Berkey-Wang instance
- its class number is 5
- has 100 rectangles
- its bin width and bin height are 100 unit

Our experimental setup is consists of UTS, LFGi and GA whose heuristics are FNF, FFF, BFDH and LGFof. First we apply GA to the problem and then best result's rectangle list is given as input to UTS and LGFi. Each set is run twice. Our experimental result is shown in Table 4.1, Table 4.2 and Table 4.3 .

Table 4.1: Result of population-generation for instance number 245 (Part 1)

| Population | Generation | GA Algo(s) | GA Res(s) | UTS | LGFi |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 10 | 10 | LGFof | 29 | 30 | 29 |
|  | 20 | LGFof / FFF | 29/5 | 30.5 | 29 |
|  | 40 | LGFof / FFF | 29.5 | 30 | 29 |
|  | 60 | BFDH / LGFof | 29 | 31.5 | 29 |
|  | 80 | LGFof | 28.5 | 30 | 29 |
|  | 100 | LGFof / BFDH | 28.5 | 30 | 29 |
|  | 200 | LGFof / BFDH | 28 | 30.5 | 29 |
|  | 300 | LGFof | 28 | 30 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 30.5 | 29 |
| 20 | 10 | LGFof | 29 | 30 | 29 |
|  | 20 | LGFof | 29 | 30 | 29 |
|  | 40 | LGFof / BFDH | 29 | 30.5 | 29 |
|  | 60 | LGFof | 28.5 | 30 | 29 |
|  | 80 | LGFof | 28.5 | 29.5 | 29 |
|  | 100 | LGFof | 28 | 29.5 | 29 |
|  | 200 | LGFof | 28 | 30.5 | 29 |
|  | 300 | LGFof | 28 | 29.5 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29.5 | 29 |
| 40 | 10 | LGFof | 29 | 30.5 | 29 |
|  | 20 | LGFof | 29.5 | 30 | 29 |
|  | 40 | LGFof | 28 | 30 | 29 |
|  | 60 | LGFof | 28.5 | 30 | 29 |
|  | 80 | LGFof | 28.5 | 30 | 29 |
|  | 100 | LGFof | 28 | 30 | 29 |
|  | 200 | LGFof | 28 | 30 | 29 |
|  | 300 | LGFof | 28 | 30 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 30 | 29 |
| 60 | 10 | LGFof | 29 | 30 | 29 |
|  | 20 | LGFof | 28.5 | 29.5 | 29 |
|  | 40 | LGFof | 28 | 30 | 29 |
|  | 60 | LGFof | 28 | 30 | 29 |
|  | 80 | LGFof | 28 | 30 | 29 |
|  | 100 | LGFof | 28 | 30 | 29 |
|  | 200 | LGFof | 28 | 30.5 | 29 |
|  | 300 | LGFof | 28 | 30.5 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29.5 | 29 |

Table 4.2: Result of population-generation for instance number 245 (Part 2)

| Population | Generation | GA Algo(s) | GA Res(s) | UTS | LGFi |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 80 | 10 | LGFof | 29 | 30 | 29 |
|  | 20 | LGFof | 29 | 30 | 29 |
|  | 40 | LGFof | 28 | 29.5 | 29 |
|  | 60 | LGFof | 28 | 30 | 29 |
|  | 80 | LGFof | 28 | 29 | 29 |
|  | 100 | LGFof | 28 | 30 | 29 |
|  | 200 | LGFof | 28 | 29.5 | 29 |
|  | 300 | LGFof | 28 | 29.5 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29 | 29 |
| 100 | 10 | LGFof | 29 | 30 | 29 |
|  | 20 | LGFof | 28.5 | 30 | 29 |
|  | 40 | LGFof | 28.5 | 30 | 29 |
|  | 60 | LGFof | 28 | 30 | 29 |
|  | 80 | LGFof | 28 | 29 | 29 |
|  | 100 | LGFof | 28 | 30 | 29 |
|  | 200 | LGFof | 28 | 29.5 | 29 |
|  | 300 | LGFof | 28 | 29.5 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29 | 29 |
| 200 | 10 | LGFof | 29 | 30 | 29 |
|  | 20 | LGFof | 28.5 | 30.5 | 29 |
|  | 40 | LGFof | 28 | 30 | 29 |
|  | 60 | LGFof | 28 | 29.5 | 29 |
|  | 80 | LGFof | 28 | 30 | 29 |
|  | 100 | LGFof | 28 | 29 | 29 |
|  | 200 | LGFof | 28 | 30 | 29 |
|  | 300 | LGFof | 28 | 29 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29.5 | 29 |
| 300 | 10 | LGFof | 29 | 29.5 | 29 |
|  | 20 | LGFof | 28.5 | 30 | 29 |
|  | 40 | LGFof | 28 | 30 | 29 |
|  | 60 | LGFof | 28 | 29.5 | 29 |
|  | 80 | LGFof | 28 | 29 | 29 |
|  | 100 | LGFof | 28 | 29.5 | 29 |
|  | 200 | LGFof | 28 | 30 | 29 |
|  | 300 | LGFof | 28 | 29.5 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29.5 | 29 |

Table 4.3: Result of population-generation for instance number 245 (Part 3)

| Population | Generation | GA Algo(s) | GA Res(s) | UTS | LGFi |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 400 | 10 | LGFof | 29 | 31 | 29 |
|  | 20 | LGFof | 29 | 30 | 29 |
|  | 40 | LGFof | 28 | 29.5 | 29 |
|  | 60 | LGFof | 28 | 29.5 | 29 |
|  | 80 | LGFof | 28 | 29.5 | 29 |
|  | 100 | LGFof | 28 | 29.5 | 29 |
|  | 200 | LGFof | 28 | 29.5 | 29 |
|  | 300 | LGFof | 28 | 29 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 30 | 29 |
| 500 | 10 | LGFof | 29 | 30.5 | 29 |
|  | 20 | LGFof | 28 | 29.5 | 29 |
|  | 40 | LGFof | 28 | 30 | 29 |
|  | 60 | LGFof | 28 | 29 | 29 |
|  | 80 | LGFof | 28 | 29.5 | 29 |
|  | 100 | LGFof | 28 | 29.5 | 29 |
|  | 200 | LGFof | 28 | 30 | 29 |
|  | 300 | LGFof | 28 | 29 | 29 |
|  | 400 | LGFof | 28 | 29.5 | 29 |
|  | 500 | LGFof | 28 | 29.5 | 29 |

As a result of this experiment, we chose our population size as 60 and our generation number as 40 .

### 4.2 A Tool for Visual Analysis of 2DBPPs

In this study, in order to analyze the results of packing heuristics and proposed algorithms we developed an application named as A Tool for Visual Analysis of 2D Bin Packing which is written in Java. We write the results of processed problem into a text file in a predefined format. Then, the result text file is given as input to the application. At the end, the application draws the position and shape of rectangles in the bins and prints them into graphical canvas.

### 4.3 MHO-SOMA

In order to solve single objective O-2DBPP, we applied MHO-SOMA. Experimental setup is composed of :

- population size is 60
- number of generation is 40
- swap mutation is used
- for each population generation step, swap mutation is applied for ten times
- swap mutation is applied for once in each offspring generation to keep diversity The general results of experiment are listed in Table 4.4 and Table 4.5. The results are compared with best known solution lower bound of University of Bologna D.E.I.S. Operations Research [7].

Table 4.4: Result of MHO-SOMA for Berkey-Wang instances

| Class | Num of Rect | lb | Result |
| :---: | :---: | :---: | :---: |
| 1 | 20 | 71 | 71 |
|  | 40 | 134 | 136 |
|  | 60 | 197 | 202 |
|  | 80 | 274 | 277 |
|  | 100 | 317 | 323 |
|  | Av. | 198.6 | 201.8 |
| 2 | 20 | 10 | 10 |
|  | 40 | 19 | 20 |
|  | 60 | 25 | 26 |
|  | 80 | 31 | 32 |
|  | 100 | 39 | 40 |
|  | Av. | 24.8 | 25.6 |
| 3 | 20 | 51 | 54 |
|  | 40 | 92 | 97 |
|  | 60 | 136 | 144 |
|  | 80 | 187 | 198 |
|  | 100 | 221 | 233 |
|  | Av. | 137.4 | 145.2 |
| 4 | 20 | 10 | 10 |
|  | 40 | 19 | 19 |
|  | 60 | 23 | 27 |
|  | 80 | 30 | 34 |
|  | 100 | 37 | 39 |
|  | Av. | 23.8 | 25.8 |
| 5 | 20 | 65 | 66 |
|  | 40 | 119 | 124 |
|  | 60 | 179 | 186 |
|  | 80 | 241 | 253 |
|  | 100 | 279 | 294 |
|  | Av. | 176.6 | 184.6 |
| 6 | 20 | 10 | 10 |
|  | 40 | 15 | 19 |
|  | 60 | 21 | 23 |
|  | 80 | 30 | 30 |
|  | 100 | 32 | 35 |
|  | Av. | 21.6 | 23.4 |

Table 4.5: Result of MHO-SOMA for Martello-Vigo instances

| Class | Num of Rect | lb | Pr. Algo |
| :---: | :---: | :---: | :---: |
| 7 | 20 | 55 | 57 |
|  | 40 | 109 | 115 |
|  | 60 | 156 | 161 |
|  | 80 | 224 | 232 |
|  | 100 | 269 | 276 |
|  | Av. | 162.6 | 168.2 |
| 8 | 20 | 58 | 59 |
|  | 40 | 112 | 115 |
|  | 60 | 159 | 166 |
|  | 80 | 223 | 227 |
|  | 100 | 274 | 284 |
|  | Av. | 165.2 | 170.2 |
| 9 | 20 | 143 | 143 |
|  | 40 | 278 | 278 |
|  | 60 | 437 | 437 |
|  | 80 | 577 | 577 |
|  | 100 | 695 | 695 |
|  | Av. | 426 | 426 |
| 10 | 20 | 42 | 43 |
|  | 40 | 74 | 76 |
|  | 60 | 98 | 106 |
|  | 80 | 123 | 134 |
|  | 100 | 153 | 164 |
|  | Av. | 98 | 104.6 |

### 4.4 MHNO-SOMA

In order to solve single objective NO-2DBPP, we applied MHNO-SOMA. We have two different setups.

First experimental setup is composed of :

- population size is 60
- number of generation is 40
- rotation mutation is used
- for each population generation step, rotation mutation is applied for ten times
- rotation mutation is applied for once in each offspring generation to keep diversity

Second experimental setup is composed of :

- population size is 60
- number of generation is 40
- swap rotate mutation is used
- for each population generation step, swap rotate mutation is applied for ten times
- swap rotate mutation is applied for once in each offspring generation to keep diversity

The general results of two experiments are listed in Table 4.6 and Table 4.7. We used simplest continous lower bound which is described in Equation 4.1.

$$
\begin{equation*}
L B=\left\lceil\frac{\sum_{j=1}^{\# \text { Rect }} w_{j} h_{j}}{W H}\right\rceil \tag{4.1}
\end{equation*}
$$

Table 4.6: Result of MHNO-SOMA for Berkey-Wang instances

| Class | Num of Rect | clb | Pr. Algo(r) | Pr. Algo(sr) |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 20 | 64 | 66 | 66 |
|  | 40 | 120 | 131 | 129 |
|  | 60 | 185 | 196 | 195 |
|  | 80 | 253 | 270 | 270 |
|  | 100 | 305 | 314 | 313 |
|  | Av. | 185.4 | 195.4 | 194.6 |
| 2 | 20 | 10 | 10 | 10 |
|  | 40 | 19 | 20 | 19 |
|  | 60 | 25 | 25 | 25 |
|  | 80 | 31 | 31 | 31 |
|  | 100 | 39 | 39 | 39 |
|  | Av. | 24.8 | 25 | 24.8 |
| 3 | 20 | 44 | 48 | 48 |
|  | 40 | 82 | 95 | 95 |
|  | 60 | 125 | 137 | 137 |
|  | 80 | 173 | 186 | 187 |
|  | 100 | 205 | 225 | 225 |
|  | Av. | 125.8 | 138.2 | 138.4 |
| 4 | 20 | 10 | 10 | 10 |
|  | 40 | 19 | 19 | 19 |
|  | 60 | 23 | 25 | 25 |
|  | 80 | 30 | 32 | 33 |
|  | 100 | 37 | 38 | 38 |
|  | Av. | 23.8 | 24.8 | 25 |
| 5 | 20 | 54 | 59 | 59 |
|  | 40 | 101 | 116 | 115 |
|  | 60 | 157 | 175 | 176 |
|  | 80 | 215 | 240 | 241 |
|  | 100 | 259 | 284 | 284 |
|  | Av. | 157.2 | 174.8 | 175 |
| 6 | 20 | 10 | 10 | 10 |
|  | 40 | 15 | 18 | 17 |
|  | 60 | 21 | 22 | 22 |
|  | 80 | 30 | 30 | 30 |
|  | 100 | 32 | 34 | 34 |
|  | Av. | 21.6 | 22.8 | 22.6 |

Table 4.7: Result of MHNO-SOMA for Martello-Vigo instances

| Class | Num of Rect | clb | Pr. Algo(r) | Pr. Algo(sr) |
| :---: | :---: | :---: | :---: | :---: |
| 7 | 20 | 47 | 52 | 52 |
|  | 40 | 97 | 106 | 107 |
|  | 60 | 140 | 152 | 153 |
|  | 80 | 197 | 216 | 217 |
|  | 100 | 238 | 260 | 259 |
|  | Av. | 143.8 | 157.2 | 157.6 |
| 8 | 20 | 48 | 53 | 53 |
|  | 40 | 96 | 106 | 105 |
|  | 60 | 141 | 155 | 154 |
|  | 80 | 195 | 213 | 214 |
|  | 100 | 241 | 261 | 262 |
|  | Av. | 144.2 | 157.6 | 157.6 |
| 9 | 20 | 94 | 143 | 143 |
|  | 40 | 180 | 275 | 275 |
|  | 60 | 276 | 435 | 435 |
|  | 80 | 371 | 573 | 573 |
|  | 100 | 450 | 693 | 693 |
|  | Av. | 274.2 | 423.8 | 423.8 |
| 10 | 20 | 38 | 41 | 41 |
|  | 40 | 69 | 73 | 73 |
|  | 60 | 94 | 101 | 101 |
|  | 80 | 122 | 129 | 130 |
|  | 100 | 153 | 161 | 162 |
|  | Av. | 95.2 | 101 | 101.4 |

### 4.5 MHO-MOGA

In order to solve multiple objective O-2DBPP, we applied MHO-MOGA. Experimental setup is composed of :

- population size is 60
- number of generation is 40
- swap mutation is used
- for each population generation step, swap mutation is applied for ten times
- swap mutation is applied for once in each offspring generation to keep diversity The general results of experiment are listed in Table 4.8 and Table 4.9. We used LGFof algorithm result in order to compare results.

Table 4.8: Result of MHO-MOGA for Berkey-Wang instances

| Class | Num of Rect | LGFof b | LGFof cg | Pr. Algo b | Pr. Algo cg |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 20 | 74 | 268.1 | 71 | 117 |
|  | 40 | 137 | 469.1 | 136 | 189.4 |
|  | 60 | 202 | 748.1 | 202 | 408.6 |
|  | 80 | 278 | 1114.6 | 277 | 545.7 |
|  | 100 | 323 | 1188.5 | 323 | 616.5 |
|  | Av. | 202.8 | 757.7 | 201.8 | 375.4 |
| 2 | 20 | 10 | 376.5 | 10 | 11.1 |
|  | 40 | 20 | 646.4 | 20 | 6.8 |
|  | 60 | 27 | 1024.2 | 26 | 150 |
|  | 80 | 32 | 1533.7 | 32 | 178.6 |
|  | 100 | 40 | 1854.6 | 40 | 152.2 |
|  | Av. | 25.8 | 1087.1 | 25.6 | 99.7 |
| 3 | 20 | 56 | 1079 | 54 | 305.6 |
|  | 40 | 102 | 2012.6 | 98 | 776.3 |
|  | 60 | 146 | 2987.3 | 145 | 1144.5 |
|  | 80 | 199 | 4375 | 198 | 1687.2 |
|  | 100 | 236 | 5217.8 | 233 | 2353.5 |
|  | Av. | 147.8 | 3134.3 | 145.6 | 1253.4 |
| 4 | 20 | 11 | 1542.3 | 10 | 115.2 |
|  | 40 | 19 | 3469.3 | 19 | 432.3 |
|  | 60 | 27 | 4758.2 | 27 | 277.5 |
|  | 80 | 34 | 6017 | 33 | 794.9 |
|  | 100 | 41 | 8005.5 | 40 | 741.7 |
|  | Av. | 26.4 | 4758.5 | 25.8 | 472.3 |
| 5 | 20 | 66 | 2699.8 | 66 | 1126.2 |
|  | 40 | 124 | 5891.1 | 123 | 2850.6 |
|  | 60 | 187 | 7584.2 | 183 | 4045.6 |
|  | 80 | 254 | 11502.6 | 251 | 5713.5 |
|  | 100 | 296 | 12875 | 292 | 7311.6 |
|  | Av. | 185.4 | 8110.5 | 183 | 4209.5 |
| 6 | 20 | 10 | 4790 | 10 | 879.1 |
|  | 40 | 19 | 9357.9 | 19 | 228.9 |
|  | 60 | 23 | 11447.4 | 23 | 612.2 |
|  | 80 | 30 | 16052.8 | 30 | 1202.3 |
|  | 100 | 36 | 20841 | 36 | 1875.9 |
|  | Av. | 23.6 | 12497.8 | 23.6 | 959.7 |

Table 4.9: Result of MHO-MOGA for Martello-Vigo instances

| Class | Num of Rect | LGFof b | LGFof cg | Pr. Algo b | Pr. Algo cg |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 20 | 58 | 4086.4 | 57 | 1200.2 |
|  | 40 | 115 | 9133.5 | 115 | 2064.1 |
|  | 60 | 161 | 12805.9 | 162 | 3594.5 |
|  | 80 | 235 | 18095.9 | 234 | 4984.6 |
|  | 100 | 276 | 22237.1 | 277 | 6297.1 |
|  | Av. | 169 | 13271.8 | 169 | 3628.1 |
| 8 | 20 | 61 | 2150.7 | 60 | 935.8 |
|  | 40 | 119 | 4482.7 | 115 | 2921.7 |
|  | 60 | 168 | 5912.5 | 166 | 3827.4 |
|  | 80 | 231 | 7719.8 | 230 | 6235.7 |
|  | 100 | 286 | 9266.1 | 285 | 6835.6 |
|  | Av. | 173 | 5906.4 | 171.2 | 4151.2 |
| 9 | 20 | 144 | 3718.7 | 143 | 2120.7 |
|  | 40 | 278 | 7808.1 | 278 | 4280 |
|  | 60 | 438 | 11027.8 | 437 | 7597.2 |
|  | 80 | 577 | 15529.9 | 577 | 9894.4 |
|  | 100 | 696 | 18976 | 695 | 12836.7 |
|  | Av. | 426.6 | 11412.1 | 426 | 7345.8 |
| 10 | 20 | 44 | 3162.7 | 45 | 531.2 |
|  | 40 | 77 | 5445.5 | 77 | 1721.7 |
|  | 60 | 108 | 9310.4 | 105 | 2670.5 |
|  | 80 | 134 | 10527.7 | 137 | 3331.5 |
|  | 100 | 165 | 13949.1 | 167 | 5850.5 |
|  | Av. | 105.6 | 8479.1 | 106.2 | 2821.1 |

### 4.6 MHNO-MOGA

In order to solve multi-objective NO-2DBPP, we applied MHNO-MOGA. We have two different setups.

First experimental setup is composed of :

- population size is 60
- number of generation is 40
- rotation mutation is used
- for each population generation step, rotation mutation is applied for ten times
- rotation mutation is applied for once in each offspring generation to keep diversity

The general results of first experiment are listed in Table 4.10 and Table 4.11. We used LGFi algorithm result in order to compare results.

Table 4.10: Result of MHNO-MOGA Ro for Berkey-Wang instances

| Class | Num of Rect | LGFi b | LGFi cg | Pr. Algo b | Pr. Algo cg |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 20 | 67 | 230.9 | 68 | 67.5 |
|  | 40 | 131 | 392.6 | 131 | 218.7 |
|  | 60 | 199 | 620.5 | 197 | 359.6 |
|  | 80 | 271 | 877.5 | 270 | 584.6 |
|  | 100 | 317 | 1123.7 | 320 | 676.6 |
|  | Av. | 197 | 649 | 197.2 | 381.4 |
| 2 | 20 | 10 | 256.8 | 10 | 1.7 |
|  | 40 | 20 | 530.6 | 20 | 6.4 |
|  | 60 | 25 | 993.6 | 25 | 160.7 |
|  | 80 | 32 | 1107.9 | 31 | 134.1 |
|  | 100 | 39 | 1078.7 | 39 | 195.7 |
|  | Av. | 25.2 | 793.5 | 25 | 99.7 |
| 3 | 20 | 50 | 918.7 | 49 | 241.3 |
|  | 40 | 97 | 1706.4 | 95 | 643.1 |
|  | 60 | 139 | 2785.7 | 139 | 1073 |
|  | 80 | 189 | 3805.2 | 192 | 1709 |
|  | 100 | 227 | 4422.1 | 229 | 2211.8 |
|  | Av. | 140.4 | 2727.6 | 140.8 | 1175.6 |
| 4 | 20 | 10 | 1794.8 | 10 | 0.8 |
|  | 40 | 19 | 3940.2 | 19 | 8 |
|  | 60 | 26 | 4109.7 | 25 | 174.8 |
|  | 80 | 33 | 5558.4 | 33 | 162.9 |
|  | 100 | 39 | 7418.4 | 40 | 360.7 |
|  | Av. | 25.4 | 4564.3 | 25.4 | 141.4 |
| 5 | 20 | 62 | 2097.4 | 60 | 555.3 |
|  | 40 | 117 | 5039.1 | 117 | 2167.1 |
|  | 60 | 180 | 7031 | 177 | 3292 |
|  | 80 | 246 | 9068.5 | 242 | 4477 |
|  | 100 | 290 | 11640 | 286 | 5573.6 |
|  | Av. | 179 | 6975.2 | 176.4 | 3213 |
| 6 | 20 | 10 | 4788.4 | 10 | 8 |
|  | 40 | 19 | 13578.4 | 19 | 14.6 |
|  | 60 | 23 | 15069.6 | 22 | 202.6 |
|  | 80 | 30 | 20176 | 30 | 999.9 |
|  | 100 | 35 | 25989.9 | 34 | 1995.2 |
|  | Av. | 23.4 | 15920.5 | 23 | 644.1 |

Table 4.11: Result of MHNO-MOGA Ro for Martello-Vigo instances

| Class | Num of Rect | LGFi b | LGFi cg | Pr. Algo b | Pr. Algo cg |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 20 | 55 | 1648.6 | 53 | 1145.2 |
|  | 40 | 109 | 3717.4 | 107 | 2968.7 |
|  | 60 | 161 | 6041 | 155 | 3869.6 |
|  | 80 | 223 | 6989.4 | 221 | 6229.4 |
|  | 100 | 271 | 9330.1 | 264 | 7259.2 |
|  | Av. | 163.8 | 5545.3 | 160 | 4294.4 |
| 8 | 20 | 56 | 1858.4 | 55 | 873.5 |
|  | 40 | 111 | 3877.2 | 107 | 2052.2 |
|  | 60 | 163 | 5440.7 | 155 | 3055.2 |
|  | 80 | 221 | 7640.6 | 214 | 4818.7 |
|  | 100 | 269 | 8981.9 | 264 | 6412.7 |
|  | Av. | 164 | 5559.8 | 159 | 3442.5 |
| 9 | 20 | 143 | 3012.3 | 143 | 2060.1 |
|  | 40 | 275 | 6081.2 | 275 | 4432.6 |
|  | 60 | 435 | 10005.9 | 435 | 7661.3 |
|  | 80 | 573 | 13963.7 | 573 | 9919.7 |
|  | 100 | 693 | 16765.3 | 693 | 12324 |
|  | Av. | 423.8 | 9965.7 | 423.8 | 7279.5 |
| 10 | 20 | 41 | 3800.9 | 43 | 396.2 |
|  | 40 | 75 | 6710.7 | 75 | 732.4 |
|  | 60 | 104 | 9447.3 | 105 | 1885.1 |
|  | 80 | 133 | 12136.1 | 134 | 2937.9 |
|  | 100 | 163 | 15835.2 | 165 | 4615.5 |
|  | Av. | 103.2 | 9586 | 104.4 | 2113.4 |

Second experimental setup is composed of :

- population size is 60
- number of generation is 40
- swap rotate mutation is used
- for each population generation step, swap rotate mutation is applied for ten times
- swap rotate mutation is applied for once in each offspring generation to keep diversity

The general results of second experiment are listed in Table 4.12 and Table 4.13. We used LGFi algorithm result in order to compare results.

Table 4.12: Result of MHNO-MOGA SwRo for Berkey-Wang instances

| Class | Num of Rect | LGFi b | LGFi cg | Pr. Algo b | Pr. Algo cg |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 20 | 67 | 230.9 | 66 | 76.9 |
|  | 40 | 131 | 392.6 | 125 | 202.8 |
|  | 60 | 199 | 620.5 | 200 | 359.2 |
|  | 80 | 271 | 877.5 | 264 | 630.7 |
|  | 100 | 317 | 1123.7 | 320 | 711.7 |
|  | Av. | 197 | 649 | 195 | 396.3 |
| 2 | 20 | 10 | 256.8 | 10 | 1 |
|  | 40 | 20 | 530.6 | 20 | 8.7 |
|  | 60 | 25 | 993.6 | 25 | 133 |
|  | 80 | 32 | 1107.9 | 31 | 131.8 |
|  | 100 | 39 | 1078.7 | 39 | 245.5 |
|  | Av. | 25.2 | 793.5 | 25 | 104 |
| 3 | 20 | 50 | 918.7 | 49 | 260.3 |
|  | 40 | 97 | 1706.4 | 97 | 581.1 |
|  | 60 | 139 | 2785.7 | 141 | 1165.3 |
|  | 80 | 189 | 3805.2 | 190 | 1775.7 |
|  | 100 | 227 | 4422.1 | 228 | 2335.5 |
|  | Av. | 140.4 | 2727.6 | 141 | 1223.6 |
| 4 | 20 | 10 | 1794.8 | 10 | 0.5 |
|  | 40 | 19 | 3940.2 | 19 | 22.6 |
|  | 60 | 26 | 4109.7 | 25 | 97.1 |
|  | 80 | 33 | 5558.4 | 33 | 346.4 |
|  | 100 | 39 | 7418.4 | 39 | 633.6 |
|  | Av. | 25.4 | 4564.3 | 25.2 | 220 |
| 5 | 20 | 62 | 2097.4 | 60 | 719.4 |
|  | 40 | 117 | 5039.1 | 116 | 2027 |
|  | 60 | 180 | 7031 | 177 | 3444.2 |
|  | 80 | 246 | 9068.5 | 245 | 5043.8 |
|  | 100 | 290 | 11640 | 286 | 5154.2 |
|  | Av. | 179 | 6975.2 | 176.8 | 3277.7 |
| 6 | 20 | 10 | 4788.4 | 10 | 4.9 |
|  | 40 | 19 | 13578.4 | 19 | 51.3 |
|  | 60 | 23 | 15069.6 | 22 | 112.3 |
|  | 80 | 30 | 20176 | 30 | 502 |
|  | 100 | 35 | 25989.9 | 34 | 1356.8 |
|  | Av. | 23.4 | 15920.5 | 23 | 405.5 |

Table 4.13: Result of MHNO-MOGA SwRo for Martello-Vigo instances

| Class | Num of Rect | LGFi b | LGFi cg | Pr. Algo b | Pr. Algo cg |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | 20 | 55 | 1648.6 | 53 | 1015.6 |
|  | 40 | 109 | 3717.4 | 108 | 2612.5 |
|  | 60 | 161 | 6041 | 156 | 3880.6 |
|  | 80 | 223 | 6989.4 | 220 | 6126.4 |
|  | 100 | 271 | 9330.1 | 267 | 6899.4 |
|  | Av. | 163.8 | 5545.3 | 160.8 | 4106.9 |
| 8 | 20 | 56 | 1858.4 | 53 | 785.3 |
|  | 40 | 111 | 3877.2 | 107 | 2170.4 |
|  | 60 | 163 | 5440.7 | 155 | 3044 |
|  | 80 | 221 | 7640.6 | 214 | 4521.9 |
|  | 100 | 269 | 8981.9 | 265 | 6237.6 |
|  | Av. | 164 | 5559.8 | 158.8 | 3351.8 |
| 9 | 20 | 143 | 3012.3 | 143 | 1885.9 |
|  | 40 | 275 | 6081.2 | 275 | 4241.6 |
|  | 60 | 435 | 10005.9 | 435 | 6762.7 |
|  | 80 | 573 | 13963.7 | 573 | 9355.1 |
|  | 100 | 693 | 16765.3 | 693 | 11598.4 |
|  | Av. | 423.8 | 9965.7 | 423.8 | 6768.7 |
| 10 | 20 | 41 | 3800.9 | 42 | 263.2 |
|  | 40 | 75 | 6710.7 | 75 | 971 |
|  | 60 | 104 | 9447.3 | 102 | 2616.8 |
|  | 80 | 133 | 12136.1 | 133 | 3837.2 |
|  | 100 | 163 | 15835.2 | 166 | 4662.4 |
|  | Av. | 103.2 | 9586 | 103.6 | 2470.1 |

### 4.7 Analysis of Algorithms

In order to analyze runtime and efficiency of algorithms, we randomly picked five different item size $(20,40,60,80,100)$ problems. Each test is runned for five times. Best values of FNF, FFF, BFDH, LGFof and LGFi are used as results. For proposed algorithms and UTS, we used average values of tests.

In each subsection, comparisons of algorithms according to 500 instance test setup are also explained in detail.

### 4.7.1 Oriented Single Objective Problems

Runtime analysis of FNF, FFF, BFDH, UTS, LGFof and MHO-SOMA for oriented single objective random picked tests are shown in Table 4.14.

Table 4.14: Runtime of algorithms for orientated single objective problems in msec

| Rect | FNF | FFF | BFDH | UTS | LGFof | MHO-SOMA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 4.0 | 4.5 | 4.6 | 81.7 | 64.0 | 21748 |
| 40 | 4.4 | 4.6 | 4.9 | 317.4 | 29.7 | 47196 |
| 60 | 4.8 | 6.0 | 7.7 | 27761 | 1529.5 | 754233 |
| 80 | 7.4 | 8.3 | 9.1 | 3050 | 334.5 | 25486 |
| 100 | 10.1 | 12.1 | 17.2 | 3746 | 291.0 | 582892 |
| Av. | 6.1 | 7.1 | 8.7 | 6991.2 | 449.7 | 286311 |

Result (bin) analysis of FNF, FFF, BFDH, UTS, LGFof and MHO-SOMA for oriented single objective random picked tests are shown in Table 4.15:

Table 4.15: Results (bin) of algorithms for orientated single objective problems

| Rect | FNF | FFF | BFDH | UTS | LGFof | MHO-SOMA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 11 | 9 | 9 | 9 | 9 | 8 |
| 40 | 17 | 13 | 13 | 13 | 13 | 12 |
| 60 | 53 | 47 | 47 | 47 | 47 | 46 |
| 80 | 30 | 24 | 23 | 24 | 24 | 23 |
| 100 | 44 | 31 | 31 | 30 | 30 | 30 |
| Av. | 31 | 24.8 | 24.6 | 24.6 | 24.6 | 23.8 |

Extra bin usage of FNF, FFF, BFDH, UTS, LGFof and MHO-SOMA for oriented single objective 500 problem set to reach results Bologna University DEIS Operations Research are shown in Table 4.16:

Table 4.16: Comparisons of heuristics and MHO-SOMA for orientated single objective 500 problem set

| Bin Sol. | FNF | FFF | BFDH | UTS | LGFof | Pr. Algo. |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| No ex. | $16.8 \%$ | $42.4 \%$ | $45.4 \%$ | $46.2 \%$ | $57.2 \%$ | $63.4 \%$ |
| 1 ex. | $20 \%$ | $36.2 \%$ | $38.6 \%$ | $38.6 \%$ | $34.8 \%$ | $32 \%$ |
| 2 ex. | $6 \%$ | $16 \%$ | $15 \%$ | $14.2 \%$ | $7.6 \%$ | $4.4 \%$ |
| 3 ex. | $9.2 \%$ | $5.2 \%$ | $1 \%$ | $1 \%$ | $0.6 \%$ | $0.2 \%$ |
| 3+ ex. | $48 \%$ | $0.2 \%$ | $0 \%$ | $0 \%$ | $0 \%$ | $0 \%$ |

### 4.7.2 Non-Oriented Single Objective Problems

Runtime analysis of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-SOMA for non-oriented single objective random picked tests are shown in Table 4.17;

Table 4.17: Runtime of algorithms for non-orientated single objective problems in msec

| Rect | FNF | FFF | BFDH | UTS | LGFof | LGFi | MHNO-SOMA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 4.0 | 4.5 | 4.6 | 81.7 | 64.0 | 65.0 | 101612 |
| 40 | 4.4 | 4.6 | 4.9 | 317.4 | 29.7 | 30.0 | 46640 |
| 60 | 4.8 | 6.0 | 7.7 | 27761 | 1529.5 | 2416.4 | 873217 |
| 80 | 7.4 | 8.3 | 9.1 | 3050 | 334.5 | 257.2 | 411683 |
| 100 | 10.1 | 12.1 | 17.2 | 3746 | 291.0 | 324.1 | 645343 |
| Av. | 6.1 | 7.1 | 8.7 | 6991.2 | 449.7 | 618.5 | 415699 |

Result (bin) analysis of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-SOMA for non-oriented single objective random picked tests are shown in Table 4.18

Table 4.18: Results (bin) of algorithms for non-orientated single objective problems

| Rect | FNF | FFF | BFDH | UTS | LGFof | LGFi | MHNO-SOMA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 11 | 9 | 9 | 9 | 9 | 8 | 8 |
| 40 | 17 | 13 | 13 | 13 | 13 | 12 | 11 |
| 60 | 53 | 47 | 47 | 47 | 47 | 46 | 46 |
| 80 | 30 | 24 | 23 | 24 | 24 | 22 | 21 |
| 100 | 44 | 31 | 31 | 30 | 30 | 29 | 28 |
| Av. | 31 | 24.8 | 24.6 | 24.6 | 24.6 | 23.4 | 22.8 |

Extra bin usage of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-SOMA (r) and MHNO-SOMA (sr) for oriented single objective 500 problem set ro reach results of continuous lower bound are shown in Table 4.19 and Table 4.20

Table 4.19: Comparisons of heuristics and MHNO-SOMA (r) for non-orientated single objective 500 problem set

| Bin | FNF | FFF | BFDH | UTS | LGFof | LGFi | Pr. Algo.(r) |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| No ex. | $16 \%$ | $24.6 \%$ | $25.6 \%$ | $21.4 \%$ | $27.4 \%$ | $31.8 \%$ | $41.4 \%$ |
| 1 ex. | $15.8 \%$ | $19.2 \%$ | $20.2 \%$ | $23.6 \%$ | $23 \%$ | $27 \%$ | $29 \%$ |
| 2 ex. | $5 \%$ | $17 \%$ | $17.4 \%$ | $19.6 \%$ | $17.8 \%$ | $15.4 \%$ | $14.4 \%$ |
| 3 ex. | $4.6 \%$ | $12.2 \%$ | $12.8 \%$ | $13.6 \%$ | $11.8 \%$ | $12.2 \%$ | $3.8 \%$ |
| 3+ ex. | $58.6 \%$ | $27 \%$ | $24 \%$ | $21.6 \%$ | $20.4 \%$ | $13.6 \%$ | $11.4 \%$ |

Table 4.20: Comparisons of heuristics and MHNO-SOMA (sr) for non-orientated single objective 500 problem set

| Bin | FNF | FFF | BFDH | UTS | LGFof | LGFi | Pr. Algo.(sr) |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| No ex. | $16 \%$ | $24.6 \%$ | $25.6 \%$ | $21.4 \%$ | $27.4 \%$ | $31.8 \%$ | $41.8 \%$ |
| 1 ex. | $15.8 \%$ | $19.2 \%$ | $20.2 \%$ | $23.6 \%$ | $23 \%$ | $27 \%$ | $28.8 \%$ |
| 2 ex. | $5 \%$ | $17 \%$ | $17.4 \%$ | $19.6 \%$ | $17.8 \%$ | $15.4 \%$ | $14 \%$ |
| 3 ex. | $4.6 \%$ | $12.2 \%$ | $12.8 \%$ | $13.6 \%$ | $11.8 \%$ | $12.2 \%$ | $4.2 \%$ |
| 3+ ex. | $58.6 \%$ | $27 \%$ | $24 \%$ | $21.6 \%$ | $20.4 \%$ | $13.6 \%$ | $11.2 \%$ |

### 4.7.3 Oriented Multi-Objective Problems

Runtime analysis of FNF, FFF, BFDH, UTS, LGFof and MHO-MOGA for oriented multi-objective random picked tests are shown in Table 4.21:

Table 4.21: Runtime of algorithms for orientated multi-objective problems in msec

| Rect | FNF | FFF | BFDH | UTS | LGFof | MHO-MOGA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 4.1 | 4.6 | 4.7 | 81.7 | 64.0 | 17515 |
| 40 | 4.7 | 4.8 | 5.0 | 317.8 | 29.9 | 88024 |
| 60 | 5.3 | 6.4 | 7.8 | 27761.4 | 1529.8 | 3034228 |
| 80 | 7.6 | 8.4 | 9.2 | 3050.4 | 334.6 | 4664450 |
| 100 | 10.6 | 12.4 | 17.4 | 3746.9 | 291.1 | 677750 |
| Av. | 6.5 | 7.3 | 8.8 | 6991.6 | 449.9 | 1696393.4 |

Result (bin/cg) analysis of FNF, FFF, BFDH, UTS, LGFof and MHO-MOGA for oriented multi-objective random picked tests are shown in Table 4.22

Table 4.22: Results (bin/cg) of algorithms for orientated multi-objective problems

| Rect | FNF | FFF | BFDH | UTS | LGFof | MHO-MOGA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 11 | 9 | 9 | 9 | 9 | 9 |
|  | 310.4 | 167.3 | 151.3 | 491.1 | 301.4 | 116.6 |
| 40 | 17 | 13 | 13 | 13 | 12 | 12 |
|  | 62.4 | 59.8 | 50 | 64.5 | 33.2 | 19.6 |
| 60 | 53 | 47 | 47 | 47 | 47 | 46 |
|  | 1207.5 | 1024.7 | 1008.5 | 1103.9 | 1021.1 | 813.7 |
| 80 | 30 | 24 | 23 | 24 | 24 | 24 |
|  | 1704.5 | 1850.1 | 1440.7 | 900.9 | 1728.7 | 421 |
| 100 | 44 | 31 | 31 | 30 | 30 | 29 |
|  | 1623.9 | 1391.2 | 1098.2 | 1088.9 | 1331.2 | 769.9 |
| Av. | 31 | 24.8 | 24.6 | 24.6 | 24.6 | 24 |
|  | 981.5 | 898.6 | 749.7 | 729.9 | 883.1 | 428.2 |

Results (bin/cg) of FNF, FFF, BFDH, UTS, LGFof and MHO-MOGA for oriented single objective 500 problem set are shown in Table 4.23

Table 4.23: Results (bin/cg) of heuristics and MHO-MOGA for orientated multiobjective 500 problem set

| Total | FNF | FFF | BFDH | UTS | LGFof | Pr. Alg. |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Bin | 9489 | 7591 | 7514 | 7521 | 7430 | 7389 |
| CG | 340954 | 333247 | 319187 | 297483 | 347076 | 126581 |

Superiority of MHO-MOGA to FNF, FFF, BFDH, UTS and LGFof for oriented multiobjective 500 problem set are shown in Table 4.24;

Table 4.24: MHO-MOGA vs. heuristics for orientated multi-objective 500 problem set

|  | FNF | FFF | BFDH | UTS | LGFof |
| :--- | :--- | :--- | :--- | :--- | :--- |
| MHO-MOGA | $100 \%$ | $100 \%$ | $100 \%$ | $100 \%$ | $97.2 \%$ |

### 4.7.4 Non-Oriented Multi-Objective Problems

Runtime analysis of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-MOGA for non-oriented multi-objective random picked tests are shown in Table 4.25:

Table 4.25: Runtime of algorithms for non-orientated multi-objective problems in msec

| Rect | FNF | FFF | BFDH | UTS | LGFof | LGFi | MHNO-MOGA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 4.1 | 4.6 | 4.7 | 81.7 | 64.0 | 65.0 | 87615 |
| 40 | 4.7 | 4.8 | 5.0 | 317.8 | 29.9 | 30.1 | 39027 |
| 60 | 5.3 | 6.4 | 7.8 | 27761.4 | 1529.8 | 2416.4 | 1545971 |
| 80 | 7.6 | 8.4 | 9.2 | 3050.4 | 334.6 | 257.4 | 645442 |
| 100 | 10.6 | 12.4 | 17.4 | 3746.9 | 291.1 | 324.3 | 257878 |
| Av. | 6.5 | 7.3 | 8.8 | 6991.6 | 449.9 | 618.6 | 515186.6 |

Result (bin/cg) analysis of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-MOGA for non-oriented multi-objective random picked tests are shown in Table 4.26;

Table 4.26: Results (bin/cg) of algorithms for non-orientated multi-objective problems

| Rect | FNF | FFF | BFDH | UTS | LGFof | LGFi | MHNO-MOGA |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 20 | 11 | 9 | 9 | 9 | 9 | 8 | 8 |
|  | 310.4 | 167.3 | 151.3 | 491.1 | 301.4 | 150.9 | 60 |
| 40 | 17 | 13 | 13 | 13 | 12 | 12 | 12 |
|  | 62.4 | 59.8 | 50 | 64.5 | 33.2 | 51.9 | 15 |
| 60 | 53 | 47 | 47 | 47 | 47 | 46 | 46 |
|  | 1207.5 | 1024.7 | 1008.5 | 1103.9 | 1021.1 | 901.4 | 700.2 |
| 80 | 30 | 24 | 23 | 24 | 24 | 22 | 21 |
|  | 1704.5 | 1850.1 | 1440.7 | 900.9 | 1728.7 | 887.4 | 821.3 |
| 100 | 44 | 31 | 31 | 30 | 30 | 29 | 28 |
|  | 1623.9 | 1391.2 | 1098.2 | 1088.9 | 1331.2 | 1131.2 | 575 |
| Av. | 31 | 24.8 | 24.6 | 24.6 | 24.6 | 23.4 | 23 |
|  | 981.5 | 898.6 | 749.7 | 729.9 | 883.1 | 624.6 | 434.3 |

Results (bin/cg) of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-MOGA (r) for non-oriented multi-objective 500 problem set (according to continuous lower bound) are shown in Table 4.27,

Table 4.27: Results (bin/cg) of heuristics and MHNO-MOGA (r) for non-orientated multi-objective 500 problem set

| Total | FNF | FFF | BFDH | UTS | LGFof | LGFi | Pr. Alg. |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Bin | 9489 | 7591 | 7514 | 7521 | 7430 | 7226 | 7175 |
| CG | 340954 | 333247 | 319187 | 297483 | 347076 | 311434 | 113925 |

Superiority of MHO-MOGA (r) vs. FNF, FFF, BFDH, UTS, LGFof and LGFi for non-oriented multi-objective 500 problem set are shown in Table 4.28;

Table 4.28: MHO-MOGA(r) vs. heuristics for non-orientated multi-objective 500 problem set

|  | FNF | FFF | BFDH | UTS | LGFof | LGFi |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| MHNO-MOGA(r) | $100 \%$ | $100 \%$ | $100 \%$ | $100 \%$ | $100 \%$ | $95 \%$ |

Results (bin/cg) of FNF, FFF, BFDH, UTS, LGFof, LGFi and MHNO-MOGA (sr) for non-oriented multi-objective 500 problem set (according to continuous lower bound) are shown in Table 4.27:

Table 4.29: Results (bin/cg) of heuristics and MHNO-MOGA(sr) for non-orientated multi-objective 500 problem set

| Total | FNF | FFF | BFDH | UTS | LGFof | LGFi | Pr. Alg. |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Bin | 9489 | 7591 | 7514 | 7521 | 7430 | 7226 | 7165 |
| CG | 340954 | 333247 | 319187 | 297483 | 347076 | 311434 | 111623 |

Superiority of MHNO-MOGA (sr) vs. FNF, FFF, BFDH, UTS, LGFof and LGFi for non-oriented multi-objective 500 problem set are shown in Table 4.30;

Table 4.30: MHO-MOGA (sr) vs. heuristics for non-orientated multi-objective 500 problem set

|  | FNF | FFF | BFDH | UTS | LGFof | LGFi |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| MHNO-MOGA(sr) | $100 \%$ | $100 \%$ | $100 \%$ | $100 \%$ | $100 \%$ | $96 \%$ |

### 4.8 Visual Analysis of Algorithms

In order to visually analyze proposed algorithms, we randomly picked a problem in our problem set. Our problem (instance) number is 108. It is a Berkey-Wang class 3 problem with item size as 20 . Bin width and bin height of the problem is 40 . Rectangle list of the problem is listed in Table 4.31:

Table 4.31: Rectangle list of instance number 108 with bin width $=40$ and bin height $=40$

| Rectangle | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Width | 30 | 13 | 7 | 29 | 7 | 19 | 13 | 21 | 16 | 17 |
| Height | 15 | 3 | 27 | 12 | 2 | 14 | 12 | 2 | 4 | 14 |
| Rectangle | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
| Width | 19 | 32 | 5 | 2 | 33 | 12 | 5 | 19 | 4 | 27 |
| Height | 4 | 30 | 17 | 8 | 23 | 23 | 15 | 19 | 31 | 23 |

First, we considered the problem as an oriented single objective bin packing problem so rotation of rectangles was not allowed and our objective is minimizing the number of bins. Then, we tested it with FNF, FFF, UTS and MHO-SOMA. Results of test are listed in Table 4.32:

Table 4.32: Results of instance number 108 as single objective BPP

|  | FNF | FFF | UTS | MHO-SOMA |
| :--- | :--- | :--- | :--- | :--- |
| Num of Bin | 7 | 5 | 5 | 4 |

All results of the test is drawn by the use of our application which is described in Section 4.2.

Visual analysis of FNF for the test is shown in Figure 4.1 and Figure 4.1:


Figure 4.1: Visual analysis of FNF (part 1)


Figure 4.2: Visual analysis of FNF (part 2)

Visual analysis of FNF for the test is shown in Figure 4.1 and Figure 4.1:


Figure 4.3: Visual analysis of FFF

Visual analysis of UTS for the test is shown in Figure 4.4.


Figure 4.4: Visual analysis of UTS

Visual analysis of our proposed algorithm MHO-SOMA for the test is shown in Figure 4.5:


Figure 4.5: Visual analysis of MHO-SOMA

## CHAPTER 5

## CONCLUSION AND FUTURE WORK

Bin packing problems are NP-hard combinatorial optimization problems. Engineering, manufacturing, transportation and logistics are dealing with these problems. In this study, we developed four algorithms for two variant (single objective and multiobjective) of 2D offline bin packing problems and tested 500 benchmark problems for each developed algorithm.

First, we tried to minimize number of bins to be packed. Two algorithms are proposed to solve this single objective 2DBPP. Our first algorithm is MHO-SOMA. It tries to minimize the number of bins for orientated 2D offline BPP. MHO-SOMA achieves optimum solutions (according to Bologna University DEIS Operations Research) for $63.4 \%$ of the problems with no extra bin, $32.0 \%$ of the problems with extra one bin and $4.4 \%$ of the problems with extra two bins.

Second single objective algorithm is MHNO-SOMA which solves non-oriented single objective 2D offline BPP. MHNO-SOMA is using two different mutation operation. First mutation operation is rotation and second mutation operation is swapmutation. Each mutation operation is applied for 500 benchmark problems. MHNOSOMA with rotation mutation achieves optimum solutions (according to continuous lower bound) for $41.4 \%$ of the problems with no extra bin, $29.0 \%$ of the problems with extra one bin and $14.4 \%$ of the problems with extra two bins. MHNO-SOMA with swap-rotation mutation achieves optimum solutions (according to continuous lower bound) for $41.8 \%$ of the problems with no extra bin, $28.8 \%$ of the problems with extra one bin and $14 \%$ of the problems with extra two bins. Experimental results show that the combination of global search of well known heuristics with individual
learning (local search of UTS) can achieve optimal results with a reasonable population and number of generations.

For real world applications, load balancing is as important as minimizing number of bins, so we proposed two different algorithms in order to solve multi-objective (bin size-load balancing) 2D offline bin packing problems. First one is MHO-MOGA and second one is MHNO-MOGA. Computational results show that well known heuristics sometimes produce better results about number of bins but not about load balancing. Our proposed MHO-MOGA and MHNO-MOGA give better results about not only number of bins but also load balancing of bins. Analysis of algorithms also shows that proposed algorithms produce better results than heuristics but the runtime of heuristics are shorter than the proposed algorithms. MHO-MOGA achieves better solutions for $97.2 \%$ of the problems than LGFof heuristic. MHNO-MOGA is also using two different mutation operation which are rotation and swap-rotation. MHNOMOGA with rotation mutation achieves better solutions for $95.0 \%$ of the problems than LGFi heuristic and MHNO-MOGA with swap-rotation mutation achieves better solutions for $96.0 \%$ of the problems than LGFi heuristic.

With this study, we have proposed a set of novel metaheuristic algorithms for NPhard combinatorial optimization problem, 2DBPP. Future works of this study will have two categories. First category is integrating other well known heuristics such as Touching Perimeter and Floor-Ceiling into the proposed algorithms. Second one is achieving truly parallelization of MH-SOMA and MH-MOGA.

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